

## ForeRunner<sup>™</sup> ATM Switch Network Configuration Manual

MANU0148-04 - Rev. A - 5/14/98

Software Version 5.2.x

## **FORE Systems, Inc.**

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#### FCC CLASS A NOTICE

**WARNING**: Changes or modifications to this unit not expressly approved by the party responsible for compliance could void this user's authority to operate this equipment.

NOTE: The ASX-200, the ASX-200WG, the ASX-200BX, the ASX-1000, and the *ForeRunnerLE* 155 have been tested and found to comply with the limits for a Class A digital device, pursuant to Part 15 of the FCC Rules. These limits are designed to provide reasonable protection against harmful interference when the equipment is operated in a commercial environment. This equipment generates, uses, and can radiate radio frequency energy and, if not installed and used in accordance with the instruction manual, may cause harmful interference to radio communications. Operation of the equipment in a residential area is likely to cause harmful interference in which case the user will be required to correct the interference at his own expense.

#### DOC CLASS A NOTICE

This digital apparatus does not exceed Class A limits for radio noise emission for a digital device as set out in the Radio Interference Regulations of the Canadian Department of Communications.

Le present appareil numerique n'emet pas de bruits radioelectriques depassant les limites applicables aux appareils numeriques de la class A prescrites dans le reglement sur le brouillage radioelectrique edicte par le ministere des Communications du Canada.

#### VCCI CLASS 1 NOTICE

この装置は、第一種情報処理装置(商工業地域において使用されるべき情報処理装置)で商工業地域での電波障害防止を目的とした情報処理装置等電波障害自主規制協議会(VCCI)基準に適合しております。

従って、住宅地域またはその隣接した地域で使用すると、ラジオ、テレビジョン受信機等に受信障害を与えることがあります。

取扱説明書に従って正しい取り扱いをして下さい。

This equipment is in the Class 1 category (Information Technology Equipment to be used in commercial and/or industrial areas) and conforms to the standards set by the Voluntary Control Council For Interference by Information Technology Equipment aimed at preventing radio interference in commercial and/or industrial areas. Consequently, when used in a residential area or in an adjacent area thereto, radio interference may be caused to radios and TV receivers, etc. Read the instructions for correct handling.

### FCC REQUIREMENTS (Notice to Users of DS1 Service)

The following instructions are provided to ensure compliance with the Federal Communications Commission (FCC) Rules, Part 68.

- This device must only be connected to the DS1 network connected behind an FCC Part 68 registered channel service unit. Direct connection is not allowed.
- (2) Before connecting your unit, you must inform the telephone company of the following information:

Port ID	REN/SOC	FIC	USOC
NM-6/DS1C	6.0N	04DU9-BN,	RJ48C
NM-2/DS1C		04DU9-DN,	
NM-8/DS1D		04DU9-1ZN, and	
NM-4/DS1D		04DU9-1SN	

- (3) If the unit appears to be malfunctioning, it should be disconnected from the telephone lines until you learn if your equipment or the telephone line is the source of the trouble. If your equipment needs repair, it should not be reconnected until it is repaired.
- (4) If the telephone company finds that this equipment is exceeding tolerable parameters, the telephone company can temporarily disconnect service, although they will attempt to give you advance notice if possible.
- (5) Under the FCC Rules, no customer is authorized to repair this equipment. This restriction applies regardless of whether the equipment is in or out of warranty.
- (6) If the telephone company alters their equipment in a manner that will affect use of this device, they must give you advance warning so as to give you the opportunity for uninterrupted service. You will be advised of your right to file a complaint with the FCC.

#### CANADIAN IC CS-03 COMPLIANCE STATEMENT

<u>NOTICE</u>: The Industry Canada label identifies certified equipment. This certification means that the equipment meets certain telecommunications network protective, operational and safety requirements. The Industry Canada label does not guarantee the equipment will operate to the user's satisfaction.

Before installing this equipment, users should ensure that it is permissible to be connected to the facilities of the local telecommunications company. The equipment must also be installed using an acceptable method of connection. In some cases, the company's inside wiring associated with a single line individual service may be extended by means of a certified connector assembly (telephone extension cord). The customer should be aware that compliance with the above conditions may not prevent degradation of service in some situations.

Repairs to certified equipment should be made by an authorized Canadian maintenance facility designated by the supplier. Any repairs or alterations made by the user to this equipment, or equipment malfunctions, may give the telecommunications company cause to request the user to disconnect the equipment.

Users should ensure for their own protection that the electrical ground connections of the power utility, telephone lines and internal metallic water pipe system, if present, are connected together. This precaution may be particularly important in rural areas.

<u>Caution</u>: Users should not attempt to make such connections themselves, but should contact the appropriate electric inspection authority, or electrician, as appropriate.

#### F1 AND E3 NOTICE

The E1 (NM-6/E1C, NM-2/E1C, NM-8/E1D, and NM-4/E1D) and E3 (NM-4/E3C, NM-2/E3C, NM-4/E3D, and NM-2/E3D) network modules that are described in this manual are approved for use in FORE Systems' host systems providing that the instructions below are strictly observed. Failure to follow these instructions invalidates the approval.

#### Pan European Approval - CE Marking

Pan European approval of the E1 network module was issued by BABT following assessment against CTR12. This means that it can be connected to ONP and unstructured PTO-provided private circuits with 120  $\Omega$  interfaces in all European countries, according to Telecommunications Terminal Equipment (TTE) Directive 91/263/EEC. Thus, the following CE mark applies:

## **C€168.X**

The E1 and E3 network modules conform to safety standard EN60950: 1992 following the provisions of Low Voltage Product Safety Directive 73/23/EEC and CE Marking Directive 93/68/EEC, and can be marked accordingly with the CE symbol.

The E1 and E3 network modules conform to EN55022: 1994 and EN50082-1: 1992 following the provisions of the EMC Directive 89/336/EEC, and can be marked accordingly with the CE symbol.

#### **National Approvals**

#### UK

Network Module	Connects to	Approval Number
E1	PTO-provided private circuits with 75 Ω interfaces	AA60953
E3	PTO-provided private circuits with 75 $\Omega$ interfaces	NS/4387/1/T/605954
CEM E1	PTO-provided private circuits with 75 $\Omega$ or 120 $\Omega$ unstructured interfaces	AA607478

#### Required User Guide Statements - UK Installation

The network modules are designed for use only with FORE Systems ATM Switches. Use of the network modules in any product not listed in this manual may result in a hazard and will invalidate the regulatory approval. The network modules must be installed in accordance with the installation instructions provided.

The following table shows the available ports and their safety status:

Ports	Safety Status
E1 and E3 Ports	TNV operating at SELV
Bus Connector	SELV

#### **CE** NOTICE

Marking by the symbol **CE** indicates compliance of this system to the EMC (Electromagnetic Compatibility) directive of the European Community and compliance to the Low Voltage (Safety) Directive. Such marking is indicative that this system meets or exceeds the following technical standards:

- EN 55022 "Limits and Methods of Measurement of Radio Interference Characteristics of Information Technology Equipment."
- EN 50082-1 "Electromagnetic compatibility Generic immunity standard Part 1: Residential, commercial, and light industry."
- IEC 1000-4-2 "Electromagnetic compatibility for industrial-process measurement and control equipment Part 2: Electrostatic discharge requirements."
- IEC 1000-4-3 "Electromagnetic compatibility for industrial-process measurement and control equipment Part 3: Radiate electromagnetic field requirements."
- IEC 1000-4-4 "Electromagnetic compatibility for industrial-process measurement and control equipment Part 4: Electrical fast transient/burst requirements."

#### SAFETY CERTIFICATIONS

ETL certified to meet Information Technology Equipment safety standards UL 1950, CSA 22.2 No. 950, and EN 60950.

## **List of Figures**

#### **List of Tables**

Pr	ef	а	C	е
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Chapter	r Summar	ies		i
Technic	al Suppor	t		ii
Typogra	phical Sty	/les		iii
Importa	nt Informa	ation Indicators		iv
Laser R	adiation N	Notice		V
Safety F	Precaution	าร		v
	Modificat	ions to Equipment		V
		nt of a FORE Systems Product		
	Power Co	ord Connection		vii
CHAPT	ER 1 C	Configuring PVCs		
1.1	General 0	Concepts		. 1 - 1
1.2	Virtual Pa	aths		. 1 - 3
	1.2.1	Through Paths		. 1 - 4
	1.2.2	Originating and Terminating Paths		
1.3	Listing Vi	irtual Paths		. 1 - 7
	1.3.1	Listing Through Paths		. 1 - 7
	1.3.2	Listing Originating and Terminating Paths		. 1 - 9
1.4	Virtual Ch	hannels		1 - 11
	1.4.1	Smart Permanent Virtual Circuits		1 - 14
	1.4.2	Listing Virtual Channels		1 - 15
1.5	Creating	PVCs and SPVCs		1 - 17
	1.5.1	Creating a Through Path		1 - 18
	1.5.2	Creating an Originating or Terminating Path	'	1 - 21
	1	.5.2.1 Shaping Multiple Originating Paths on a Single Port	'	1 - 24
	1	.5.2.2 Terminating a PVC at a Switch		
	1	.5.2.3 Creating ATM ARP Entries		
	-	.5.2.4 Listing ATM ARP Entries		
	1.5.3	Creating a Virtual Channel		
	1.5.4	Creating a SPANS SPVC		
	1.5.5	Displaying SPANS SPVC Information		1 - 33

1.6	Traffic Typ	pes1	- 35
1.7	Traffic Po	licing (Usage Parameter Control)	- 36
	1.7.1	Leaky Bucket Algorithm	- 36
	1.7.2	Non-conforming Cells: Tagging vs. Dropping	
	1.7.3	UPC Traffic Contract Parameters	
	1.7.4	AMI UPC Commands	- 39
СНАРТ	ER 2 C	Configuring Classical IP	
2.1	Introducti	on	2 - 1
	2.1.1	Logical IP Subnets	
	2.1.2	Classical IP Interfaces	
	2.1.3	SPANS Interface	2 - 3
2.2	Address I	Registration and ILMI	2 - 4
	2.2.1	NSAP Addresses	
	2.2.2	Operating with ILMI Support	
	2.2.3	Operating without ILMI Support	
	2.2.4	Configuration	
2.3		ARP Servers	
	2.3.1	Theory	
	2.3.2	Configuring a FORE Switch to be an ARP Server	
	2.3.3 2.3.4	Classical IP Operation	
2.4		IP PVCs	
2.4			
	2.4.1	Theory and Configuration	
2.5		ng the Network	
2.5	Ū	Third-Party Host with No ILMI and No RFC-1577 Support	
	2.5.1 2.5.2	Third-Party Switch with ILMI and No RFC-1577 Support	
	2.5.2	Third-Party Switch with RFC-1577 and No ILMI Support	
CHAPT		Configuring an Emulated LAN	
3.1	Introducti	on	
	3.1.1	Ethernet ELANs	
	3.1.2	Token Ring ELANs	
3.2		mponents	
	3.2.1	LAN Emulation Client (LEC)	
	3.2.2	LAN Emulation Configuration Server (LECS)	
	3.2.3	LAN Emulation Server (LES)	
	3.2.4	Broadcast and Unknown Server (BUS)	
3.3		LAN Operation	
	3.3.1	Initialization	3 - 6

	3.3.2 3.3.3	Registration and Address Resolution	
3.4		uted LAN Emulation	
3.4			
	3.4.1	Single Server LANE Services Model	
		3.4.1.1 Using a Single Server	3 - 10
	3.4.2		
	3.4.2	Distributed LAN Emulation Model	
		3.4.2.2 Advantages of DLE	
		3.4.2.2.1 Load Sharing	
		3.4.2.2.2 Improved Performance for Remote LECs	3 - 13 3 <sub>-</sub> 15
		3.4.2.2.3 Fault Tolerance	
		3.4.2.2.3.1 Single Server ELAN	
		3.4.2.2.3.2 DLE ELAN	
3.5	EL ANL	Access Control	
3.6		uring an ELAN	
3.0	U	•	
	3.6.1	Configuring an LECS Configuration Database File	
		3.6.1.1 Before You Begin	3 - 23
		3.6.1.2 LECS Configuration File Syntax	
		3.6.1.3 Defining an ELAN	3 - 30
		3.6.1.4 Defining a Client	
		3.6.1.5 LECS Control Parameters	
	3.6.2	3.6.1.6 LECS MPOA Parameters	
	3.6.2 3.6.3	Sample LECS Configuration File	
	3.0.3	3.6.3.1 Starting the LECS	
		3.6.3.2 Starting the DLE LES/BUS Peer Servers	
	3.6.4	Starting the LEC(s) and Joining an ELAN	
3.7			
3.7		ding an ELAN to Use DLE	
	3.7.1	Edit the LECS.CFG File	
	3.7.2	Delete the LES and BUS	
	3.7.3	Upgrade the Switches Running Services	
	3.7.4 3.7.5	Create the DLE Peer Servers	
	3.7.5 3.7.6		
	3.7.6	Restart the LECS	
	3.7.7 3.7.8	Create the Last DLE Peer	
	3.7.6 3.7.9	Add the Last DLE Peer to Each Peer List	
	3.7.10	Update the LECS.CFG File	
	3.7.10	Transfer the Final LECS.CFG File	2 - 52 3 - 52
	3.7.11	Restart the LECS	
	0.7.12	Notice the LLOC	55

3.8	Upgrad	ling an ELAN without Using DLE
	3.8.1	Deleting the Non Co-located Services
		3.8.1.1 Administer Down the Services
		3.8.1.2 Delete the Non Co-located LES and BUS
		3.8.1.2.1 Edit the LECS.CFG File
	3.8.2	Upgrade the Switches Running Services
	3.8.3	Recreate the LES and BUS Together 3 - 56
	3.8.4	Administer the Services Up
CHAP	ΓER 4	MPOA
4.1	Overvi	ew of LANE/MPOA
4.2	LANE F	Primer
	4.2.1	LANE Components
	4.2.2	An Example LANE Configuration
		4.2.2.1 The Initialization Process
		4.2.2.2 The Connection Process
		4.2.2.3 Multicast and Broadcast Packets
		4.2.2.4 Accessing Fast Ethernet and FDDI Networks
		4.2.2.5 Multiple ELANs 4 - 6
		4.2.2.6 Distributed LAN Emulation
		4.2.2.7 Automatic ELAN Selection
		4.2.2.8 Intelligent BUS
4.3	An Intro	oduction to Multi-Protocol Over ATM4 - 7
	4.3.1	LANE Without MPOA 4 - 7
	4.3.2	Why MPOA?
	4.3.3	MPOA Components
	4.3.4	MPOA Example 4 - 11
		4.3.4.1 MPS Configuration
		4.3.4.2 Initialization
		4.3.4.3 Flow Analysis
		4.3.4.4 Making a Shortcut
		4.3.4.5 Shortcut Teardown
CHAP	ΓER 5	ForeThought PNNI
5.1	FT-PNN	NI Routing
	5.1.1	Hello Protocol
	5.1.2	Topology Database Exchange 5 - 2
	5.1.3	Flooding
	5.1.4	Hierarchical Routing 5 - 3

	5.1.4.1 Hierarchical Addressing	5 - 3
	5.1.4.1.1 Switch Prefix	5 - 3
	5.1.4.1.2 Switch Summary Prefix	5 - 4
	5.1.4.1.3 Peer Group ID	5 - 4
	5.1.4.2 Path Computation	5 - 4
5.2	The Physical Network	5 - 5
	5.2.1 Peer Groups	5 - 8
	5.2.2 Peer Group Topology	5 - 8
	5.2.3 Border Switches	5 - 8
	5.2.4 Peer Group Summary Node (PGSN)	
	5.2.5 Backbone Topology	
	5.2.6 Single Switch Perspective	5 - 9
CHAP	PTER 6 ATM Forum PNNI	
6.1	PNNI Routing Protocol	6 - 1
	6.1.1 Hello Protocol	6 - 1
	6.1.2 Database Exchange Protocol	
	6.1.3 Flooding Protocol	6 - 2
	6.1.4 Path Computation	6 - 2
	6.1.5 Hierarchical Routing	6 - 3
6.2	PNNI Signalling Protocol	6 - 4
	6.2.1 Source Routing	6 - 4
	6.2.2 Crankback	6 - 4
6.3	Internetworking between PNNI and FT-PNNI	6 - 5
	6.3.1 Gateway Switches and Split Switches	6 - 5
	6.3.2 Dynamic Leaking of Reachability Information	6 - 6
	6.3.2.1 Areas	6 - 6
	6.3.2.1.1 Peer Groups in Areas	
	6.3.2.1.2 Area IDs	
	6.3.2.1.3 Levels	-
	6.3.2.2 Domains	
	6.3.2.2.1 Configuring Domains	
	6.3.2.3 Propagation of Reachability Information	
	6.3.2.3.1 Policies	
	6.3.2.3.2 Scope	i - 10
	6.3.2.3.3 The Process for Leaking Reachability	
	Information	
	6.3.2.4 VP Trunk QoS Extension	) <b>-</b> 11

CHAP1	ΓER 7	Signalling			
7.1	VCI AII	ocation Range			7 - 1
	7.1.1	Determining the V	'CI Allocat	tion Range with ILMI Down	7-2
	7.1.2	Determining the V	CI Allocat	tion Range with ILMI Up	7 - 4
7.2	Signall	ng Scope			7-7
	7.2.1	VC-Space			7-7
	7.2.2	Dynamic Paths			7 - 8
7.3	Signall	ng Channel Auto Cor	nfiguration	Procedures	7 - 9
	7.3.1	Overview of Signa	ılling Char	nnel Auto Configuration	7 - 9
	7.3.2			el Auto Configuration	
		7.3.2.1 Specifying		and Interface Version	
		7.3.2.1.1		s of Valid Configurations	
		7.3.2.1.2	•	s of Invalid Configurations	
				be and Mode	
		7.3.2.2.1		s of Valid Configurations	
<b>-</b> 4	A.I.	7.3.2.2.2		s of Invalid Configurations	
7.4				meters	
	7.4.1				
				tion of Traffic December	
	7.4.2			tion of Traffic Parameters	
	7.4.2	UNI S.A			7 - 19
CHAPT	ΓER 8	Security			
8.1	Config	uring Userids			8 - 1
	8.1.1	Login Authenticati	on Metho	d	8-3
				1	
		8.1.1.2 SecurID A		tion	
		8.1.1.2.1	SecurID	Protection on Switches	8 - 3
		8.1.1.2.2		Passcode	
			1.2.2.1	PIN Number	
			1.2.2.2	SecurID Tokens	
		8.1.1.2.3		Server	
			1.2.3.1	Slave Server	
			1.2.3.2 1.2.3.3	Server Database  Data Encryption between the Server	8-4
		0.1.	1.2.3.3	and Switches	Q <sub>-</sub> 5
		8.1.1.2.4	SecurID	AMI Commands	
		8.1.1.2.5		SecurID on a Switch	
			1.2.5.1	Installing the Server Software	
			1.2.5.2	Transferring the Configuration File	
			1.2.5.3	Editing the Server Configuration File	
		8.1.	1.2.5.4	An Example Login Using SecurID	

	8.1.2	AMI Command Privileges	. 8 - 9
		8.1.2.1 Admin Privileges	. 8 - 9
		8.1.2.2 User Privileges	. 8 - 9
	8.1.3	AMI Access Levels	
		8.1.3.1 Serial Access	
		8.1.3.2 Network Access	
		8.1.3.3 All Access	
		8.1.3.4 No Access	
	8.1.4	Userid Password	
	8.1.5	Privilege Level for Unlisted Users	
8.2	IP Filter	ring	8 - 11
	8.2.1	Authorized IP Address Table	
	8.2.2	IP Filtering Flags	
		8.2.2.1 Strict Source Routing Flag	
		8.2.2.2 Loose Source Routing Flag	
		8.2.2.3 All Flag	
	8.2.3	IP Access Statistics	
8.3	NSAP F	Filtering	8 - 13
	8.3.1	Filters and Templates	8 - 13
	8.3.2	NSAP Filtering Lookup	
	8.3.3	NSAP Filtering Statistics	8 - 14
СНАІ	PTER 9	Configuring Timing	
9.1	Overvie	ew	. 9 - 1
9.2		Modes	
9.3	J	slock	
	9.3.1	Failover of the Switchclock	
9.4		vel Timing	
9. <del>4</del> 9.5		•	
9.5	_	Configuration Examples	
	9.5.1	Configuring Timing on an ASX-200BX or ASX-200WG	
	9.5.2	Configuring Timing on an ASX-1000 (Single Timing Domain)	
	9.5.3	Configuring Timing on an ASX-1000 (Multiple Timing Domains)	
	9.5.4	Configuring Timing on a ForeRunnerLE 155	. 9 - 5
APPE	ENDIX A	Configuring SNMP	
A.1	SNMP	Indexing	. A - 1
A.2	SNMP :	Traps	.A - 3
	A.2.1	Adding SNMP Trap Destinations	A - 21
	A.2.2	Displaying SNMP Trap Destinations	
	A 2 3	Removing SNMP Trap Destinations	

B.1       Configuring CES Connections       B - 2         B.1.1       Creating a New CES Connection       B - 2         B.1.2       Displaying CES Connections       B - 5         APPENDIX C Converting from FT-PNNI to PNNI         C.1       ASX-1000 Routing Configuration Issues       C - 2         C.1.1       ASX-1000s in FT-PNNI Peer Groups       C - 2         C.1.2       ASX-1000s in PNNI Areas       C - 3         C.1.3       Multiple Gateways in an ASX-1000       C - 4         C.1.4       Migrating from FT-PNNI to PNNI Routing       C - 4         C.2       Migration of a Non-Hierarchical FT-PNNI Network       C - 5         C.2.1       Migration Overview       C - 5         C.2.2       Detailed Migration Example       C - 6         C.3.1       Migration of a Hierarchical FT-PNNI Network       C - 10         C.3.1       Migration of a Hierarchical FT-PNNI Network with a       C - 10         C.3.1.1       Migration Starting with the Backbone       C - 10         C.3.1.1.1       Upgrade the Switches       C - 11         C.3.1.2.1       C - 10       C - 12         C.3.1.2.1.1       Upgrade the Switches       C - 15         C.3.1.2.1.2       C convert Peer Group A       C - 26
B.1.2   Displaying CES Connections.   B - 5
APPENDIX C Converting from FT-PNNI to PNNI C.1 ASX-1000 Routing Configuration Issues
C.1         ASX-1000 Routing Configuration Issues         C - 2           C.1.1         ASX-1000s in FT-PNNI Peer Groups         C - 2           C.1.2         ASX-1000s in PNNI Areas         C - 3           C.1.3         Multiple Gateways in an ASX-1000         C - 4           C.1.4         Migrating from FT-PNNI to PNNI Routing         C - 4           C.2         Migration of a Non-Hierarchical FT-PNNI Network         C - 5           C.2.1         Migration Overview         C - 6           C.2.2         Detailed Migration Example         C - 6           C.3         Migration of a Hierarchical FT-PNNI Network with a         C - 10           C.3.1         Migration Starting with the Backbone         C - 10           C.3.1.1         Migration Starting with the Backbone         C - 11           C.3.1.1.1.1         Upgrade the Switches         C - 11           C.3.1.1.1.2         Convert the Backbone         C - 12           C.3.1.2.1         Overview of the Migration         C - 12           C.3.1.2.1.1         Upgrade the Switches         C - 13           C.3.1.2.1.2         Convert Peer Groups         C - 15           C.3.1.2.1.3         Convert Peer Group C         C - 15           C.3.1.2.1.4         Convert Peer Group A         C - 26<
C.1.1 ASX-1000s in FT-PNNI Peer Groups
C.1.2         ASX-1000s in PNNI Areas         C - 3           C.1.3         Multiple Gateways in an ASX-1000         C - 4           C.1.4         Migrating from FT-PNNI to PNNI Routing         C - 4           C.2         Migration of a Non-Hierarchical FT-PNNI Network         C - 5           C.2.1         Migration Overview         C - 6           C.2.2         Detailed Migration Example         C - 6           C.3.1         Migration of a Hierarchical FT-PNNI Network         C - 10           C.3.1         Migration of a Hierarchical FT-PNNI Network with a Contiguous Backbone         C - 10           C.3.1.1         Migration Starting with the Backbone         C - 10           C.3.1.1.1         Migration Overview         C - 10           C.3.1.1.1.1         Upgrade the Switches         C - 10           C.3.1.2.1.2         Convert the Individual Peer Groups         C - 12           C.3.1.2.1.1         Upgrade the Switches         C - 13           C.3.1.2.1.2         Convert Peer Group C         C - 15           C.3.1.2.1.2         Convert Peer Group A         C - 23           C.3.1.2.1.4         Convert Peer Group B         C - 26
C.1.3 Multiple Gateways in an ASX-1000
C.1.4         Migrating from FT-PNNI to PNNI Routing         C - 4           C.2         Migration of a Non-Hierarchical FT-PNNI Network         C - 5           C.2.1         Migration Overview         C - 5           C.2.2         Detailed Migration Example         C - 6           C.3         Migration of a Hierarchical FT-PNNI Network         C - 10           C.3.1         Migration of a Hierarchical FT-PNNI Network with a Contiguous Backbone         C - 10           C.3.1.1         Migration Starting with the Backbone         C - 10           C.3.1.1.1         Migration Overview         C - 11           C.3.1.1.1.1         Upgrade the Switches         C - 12           C.3.1.1.1.2         Convert the Individual Peer Groups         C - 12           C.3.1.2.1         Overview of the Migration         C - 12           C.3.1.2.1.1         Upgrade the Switches         C - 12           C.3.1.2.1.2         Convert Peer Group C         C - 13           C.3.1.2.1.3         Convert Peer Group A         C - 23           C.3.1.2.1.4         Convert Peer Group B         C - 26
C.2 Migration of a Non-Hierarchical FT-PNNI Network
C.2.1         Migration Overview         C - 6           C.2.2         Detailed Migration Example         C - 6           C.3         Migration of a Hierarchical FT-PNNI Network         C - 10           C.3.1         Migration of a Hierarchical FT-PNNI Network with a Contiguous Backbone         C - 10           C.3.1.1         Migration Starting with the Backbone         C - 10           C.3.1.1.1         Upgrade the Switches         C - 11           C.3.1.1.1.2         Convert the Backbone         C - 11           C.3.1.1.1.3         Convert the Individual Peer Groups         C - 12           C.3.1.2.1         Overview of the Migration         C - 19           C.3.1.2.1.1         Upgrade the Switches         C - 19           C.3.1.2.1.2         Convert Peer Group C         C - 19           C.3.1.2.1.3         Convert Peer Group A         C - 23           C.3.1.2.1.4         Convert the Backbone         C - 23           C.3.1.2.1.5         Convert Peer Group B         C - 26
C.2.2         Detailed Migration Example
C.3       Migration of a Hierarchical FT-PNNI Network       C - 10         C.3.1       Migration of a Hierarchical FT-PNNI Network with a       C - 10         C.3.1.1       Migration Starting with the Backbone       C - 10         C.3.1.1.1       Migration Overview       C - 11         C.3.1.1.1.1       Upgrade the Switches       C - 11         C.3.1.1.1.2       Convert the Backbone       C - 12         C.3.1.2.1       Overview of the Individual Peer Groups       C - 19         C.3.1.2.1       Overview of the Migration       C - 19         C.3.1.2.1.1       Upgrade the Switches       C - 19         C.3.1.2.1.2       Convert Peer Group C       C - 19         C.3.1.2.1.3       Convert Peer Group A       C - 23         C.3.1.2.1.4       Convert the Backbone       C - 23         C.3.1.2.1.5       Convert Peer Group B       C - 26
C.3.1 Migration of a Hierarchical FT-PNNI Network with a  Contiguous Backbone
Contiguous Backbone         C - 10           C.3.1.1 Migration Starting with the Backbone         C - 10           C.3.1.1.1 Migration Overview         C - 11           C.3.1.1.1.1 Upgrade the Switches         C - 11           C.3.1.1.1.2 Convert the Backbone         C - 11           C.3.1.1.1.3 Convert the Individual Peer Groups         C - 14           C.3.1.2 Migration Starting with the Peer Groups         C - 19           C.3.1.2.1 Overview of the Migration         C - 19           C.3.1.2.1.1 Upgrade the Switches         C - 19           C.3.1.2.1.2 Convert Peer Group C         C - 19           C.3.1.2.1.3 Convert Peer Group A         C - 23           C.3.1.2.1.4 Convert the Backbone         C - 23           C.3.1.2.1.5 Convert Peer Group B         C - 26
C.3.1.1       Migration Starting with the Backbone       C - 10         C.3.1.1.1       Migration Overview       C - 11         C.3.1.1.1.1       Upgrade the Switches       C - 11         C.3.1.1.1.2       Convert the Backbone       C - 14         C.3.1.1.1.3       Convert the Individual Peer Groups       C - 19         C.3.1.2.1       Overview of the Migration       C - 19         C.3.1.2.1.1       Upgrade the Switches       C - 19         C.3.1.2.1.2       Convert Peer Group C       C - 19         C.3.1.2.1.3       Convert Peer Group A       C - 23         C.3.1.2.1.4       Convert the Backbone       C - 23         C.3.1.2.1.5       Convert Peer Group B       C - 26
C.3.1.1.1 Migration Overview
C.3.1.1.1.1       Upgrade the Switches.       C - 11         C.3.1.1.1.2       Convert the Backbone       C - 11         C.3.1.1.1.3       Convert the Individual Peer Groups.       C - 12         C.3.1.2       Migration Starting with the Peer Groups.       C - 19         C.3.1.2.1       Overview of the Migration       C - 19         C.3.1.2.1.1       Upgrade the Switches.       C - 19         C.3.1.2.1.2       Convert Peer Group C       C - 19         C.3.1.2.1.3       Convert Peer Group A       C - 23         C.3.1.2.1.4       Convert the Backbone       C - 23         C.3.1.2.1.5       Convert Peer Group B       C - 26
C.3.1.1.1.2       Convert the Backbone       C - 14         C.3.1.2.1       Convert the Individual Peer Groups       C - 14         C.3.1.2.1       Overview of the Migration       C - 19         C.3.1.2.1.1       Upgrade the Switches       C - 19         C.3.1.2.1.2       Convert Peer Group C       C - 19         C.3.1.2.1.3       Convert Peer Group A       C - 23         C.3.1.2.1.4       Convert the Backbone       C - 23         C.3.1.2.1.5       Convert Peer Group B       C - 26
C.3.1.1.1.3 Convert the Individual Peer Groups C - 14 C.3.1.2 Migration Starting with the Peer Groups C - 19 C.3.1.2.1 Overview of the Migration
C.3.1.2       Migration Starting with the Peer Groups.       C - 19         C.3.1.2.1       Overview of the Migration       C - 19         C.3.1.2.1.1       Upgrade the Switches.       C - 19         C.3.1.2.1.2       Convert Peer Group C       C - 19         C.3.1.2.1.3       Convert Peer Group A       C - 23         C.3.1.2.1.4       Convert the Backbone       C - 23         C.3.1.2.1.5       Convert Peer Group B       C - 26
C.3.1.2.1 Overview of the Migration
C.3.1.2.1.1       Upgrade the Switches.       C - 19         C.3.1.2.1.2       Convert Peer Group C       C - 19         C.3.1.2.1.3       Convert Peer Group A       C - 23         C.3.1.2.1.4       Convert the Backbone       C - 23         C.3.1.2.1.5       Convert Peer Group B       C - 26
C.3.1.2.1.2       Convert Peer Group C
C.3.1.2.1.3 Convert Peer Group A
C.3.1.2.1.4 Convert the Backbone
C.3.1.2.1.5 Convert Peer Group B
Non-Contiguous Backbone
APPENDIX D Configuring FramePlus Modules
D.1 Frame Relay Overview
D.1.1 Interworking Function (IWF)
D.1.1.1 Translation Mode
D.1.1.2 Transparent Mode
D.2 Configuring the Module Level
D.2.1 Dividing the Buffer Space
D.2.2 Setting the Thresholds
D.2.2.1 Noting the CLP0PPD Threshold
D.2.2.2 Configuring the CLP1EPD Threshold

	D.2.2.3	Configuring the CLP0EPD Threshold	D - 7
	D.2.2.4	Configuring the CLP1PPD Threshold	D - 8
D.3	Profiles		D <b>-</b> 9
	D.3.1 EPD/F	PPD Profile	D - 9
		Profile	
	D.3.3 Frame	Relay Rate Profile	D - 10
	D.3.4 LMI P	rofile	D - 11
		e Profile	
		Profile	
D.4	Services		D - 11
D.5	Configuring Fran	ne Relay	D - 12
	D.5.1 Choos	sing Frame Relay Profiles	D - 12
		ng the Services for Frame Relay	
		ng Frame Relay PVCs	
		juring Frame Relay SPVCs	
		Creating a Frame Relay SPANS SPVC	
		Creating a Frame Relay PNNI SPVC	
D.6	Configuring FUNI		
		ging the Application Key	
		ng the Profiles for FUNI	
		D - 19	
		D - 20	
	D.6.5 Config	juring FUNI SPVCs	D - 21
		Creating a FUNI SPANS SPVC	
		Creating a FUNI PNNI SPVC	
D.7	Upgrading the FramePlus Network Module Software		

Acronyms

Glossary

Index

## **List of Figures**

CHAPTER 1	Configuring PVCs
Figure 1.1	The Cell
Figure 1.2	Virtual Channels in a Virtual Path1 - 3
Figure 1.3	An Example of a Virtual Path
Figure 1.4	Composition of a Virtual Path
Figure 1.5	An Example of a Through Path1 - 5
Figure 1.6	Through Paths are Unidirectional
Figure 1.7	Using Originating and Terminating Paths for Bandwidth Allocation
Figure 1.8	An Example of a Virtual Channel
Figure 1.9	Example of a Virtual Channel1 - 12
Figure 1.10	Virtual Channels are Unidirectional1 - 12
Figure 1.11	Virtual Channels Created on Terminating Path C3 3 1 - 13
Figure 1.12	Virtual Channels Created on Originating Path C2 21 - 13
Figure 1.13	The Path of a Cell Via SPVCs1 - 14
Figure 1.14	PVPs Looped through Port 1A2 and Output on Port 1A1 to WAN
Figure 1.15	PVPs Coming in Port 1A1 from WAN and Looped through Port 1A2
CHAPTER 2	Configuring Classical IP
Figure 2.1	Configuring a Third-Party Host with No ILMI and No RFC-1577 Support
Figure 2.2	Configuring a Third-Party Switch with ILMI Support and No RFC-1577
Figure 2.3	Configuring a Third-Party Switch with RFC-1577 and No ILMI Support2 - 14
CHAPTER 3	Configuring an Emulated LAN
Figure 3.1	Basic Emulated LAN Interconnections
Figure 3.2	ELAN Operation
Figure 3.3	Single Server LANE Services Model
Figure 3.4	Broadcast IP-ARP Request
Figure 3.5	IP ARP Response Handling

	Figure 3.6	Distributed LAN Emulation Model
	Figure 3.7	IP ARP Broadcast from LEC 1 to LEC 9 3 - 13
	Figure 3.8	Re-distributing the Broadcast across DLE Peer Servers 3 - 13
	Figure 3.9	LE-ARP for Unknown Host Sent to Proxies (not shown) and
		DLE Peer Servers
	Figure 3.10	LE-ARP Query Answered by One DLE Peer Server and
	=	Re-distributed by Another
	Figure 3.11	LE-ARP Response Delivered and LEC 9 Contacts LEC 1 3 - 15
	Figure 3.12	ELAN with a Single Server and Multiple Switches
	Figure 2.42	Connecting to Services
	Figure 3.13	ELAN with Single Server and Remote Connection to Server 3 - 17
	Figure 3.14	ELAN with Single Server in Operation
	Figure 3.15	Registrations on an ELAN with Multiple Servers
	Figure 3.16	ELAN with Multiple Servers in Operation
	Figure 3.17	Failure of One ELAN Server and the Recovery Process 3 - 19
	Figure 3.18	ELAN Re-established Using the Second Server
	Figure 3.19	Sample LECS Configuration File (Part One of Two)
	Figure 3.20	Sample LECS Configuration File (Part Two of Two) 3 - 37
CH	IAPTER 4	MPOA
Cŀ	IAPTER 4 I Figure 4.1	<b>MPOA</b> An Example of an ELAN
Cŀ		
CH	Figure 4.1	An Example of an ELAN 4 - 4
CH	Figure 4.1 Figure 4.2	An Example of an ELAN
	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4	An Example of an ELAN
	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5	An Example of an ELAN       4 - 4         LANE without MPOA       4 - 7         LANE with MPOA       4 - 9         MPOA Example Network       4 - 11    ForeThought PNNI
	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1	An Example of an ELAN       4 - 4         LANE without MPOA       4 - 7         LANE with MPOA       4 - 9         MPOA Example Network       4 - 11         ForeThought PNNI         Example of a 13-byte Switch Prefix       5 - 3
	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5	An Example of an ELAN
	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1 Figure 5.2	An Example of an ELAN
	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1	An Example of an ELAN
	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1 Figure 5.2	An Example of an ELAN
CH	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1 Figure 5.2 Figure 5.3 Figure 5.4	An Example of an ELAN
CH	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1 Figure 5.2 Figure 5.3 Figure 5.4	An Example of an ELAN
CH	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1 Figure 5.2 Figure 5.3 Figure 5.4  IAPTER 6	An Example of an ELAN
CH	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1 Figure 5.2 Figure 5.3 Figure 5.4  IAPTER 6 Figure 6.1	An Example of an ELAN
CH	Figure 4.1 Figure 4.2 Figure 4.3 Figure 4.4  IAPTER 5 Figure 5.1 Figure 5.2 Figure 5.3 Figure 5.4  IAPTER 6 Figure 6.1 Figure 6.2	An Example of an ELAN

APPENDIX C	Converting from FT-PNNI to PNNI
Figure C.1	Invalid Configuration of ASX-1000 Split between Two FT-PNNI Peer Groups
Figure C.2	Invalid Configuration of ASX-1000 Split between Two
9	PNNI Peer Groups
Figure C.3	Multiple Fabrics of an ASX-1000 Incorrectly Configured
J	as Gateways
Figure C.4	A Non-Hierarchical FT-PNNI Network
Figure C.5	S5 Changed to a Gateway Switch
Figure C.6	S3 Changed to a Gateway Switch
Figure C.7	S4 Changed to a Gateway Switch and S5 to PNNI
Figure C.8	A Completely Converted PNNI Network
Figure C.9	Hierarchical FT-PNNI Network with 3 Peer Groups and a
	Contiguous Backbone
Figure C.10	C.1 and A.1 as Gateway Switches
Figure C.11	Peer Group Severed from the Rest of the FT-PNNI Area C - 13
Figure C.12	Peer Group C before the Conversion to PNNI
Figure C.13	C.1 and C.2 Not Part of Peer Group C
Figure C.14	A Completely Converted PNNI Network
Figure C.15	C.6 as a Gateway
Figure C.16	Peer Group C Fully Migrated to PNNI
Figure C.17	Peer Group B Disconnected from Peer Group A
Figure C.18	A Migrated Backbone
Figure C.19	Hierarchical FT-PNNI Network with 3 Peer Groups and a
	Non-contiguous Backbone
Figure C.20	Hierarchical PNNI Network after Migration
APPENDIX D	Configuring <i>FramePlus</i> Modules
Figure D.1	Buffer Sizes Configured Using the setmem Command
Figure D.2	CLP0PPD Automatically Calculated
Figure D.3	Calculated CLP1EPD
Figure D.4	Calculated CLP0EPD
Figure D.5	Calculated CLP1PPD D - 8

## List of Figures

## **List of Tables**

CHAPTER 1	Configuring PVCs
Table 1.1	Summary of Traffic Contract Variables and Policing Actions1 - 37
CHAPTER 3	Configuring an Emulated LAN
Table 3.1	LECS Configuration File Parameters
CHAPTER 7	Signalling
Table 7.1	Action Taken Based on Both Switches' Signalling Channel
	Configurations
Table 7.2	Action Taken Based on the Peer's Supported MIB Variable 7 - 11
Table 7.3	Valid Type and Version Combinations
Table 7.4	Invalid Type and Version Combinations
Table 7.5	Valid Scope and Mode Combinations7 - 15
Table 7.6	Invalid Scope and Mode Combinations7 - 16
Table 7.7	UNI 3.1 Allowable Combination of Traffic Parameters in ForeThought 5.2.x
CHAPTER 8	Security
Table 8.1	Possible Login Scenarios
APPENDIX A	Configuring SNMP
Table A.1	ASX-200WG/ASX-200BX Port NumberingA - 2
Table A.2	SNMP Traps Supported on the ForeRunner Switches
Table A.3	Message Type Encodings for Trap 2003
Table A.4	Error Codes for Trap 2003
APPENDIX D	Configuring <i>FramePlus</i> Modules
Table D.1	Buffer Models

#### List of Tables

## **Preface**

This manual provides the technical information needed to configure the *ForeRunner*<sup>TM</sup> family of ATM Switches, the *ForeRunner* LAN and WAN network modules, and the accompanying *ForeThought*<sup>TM</sup> software. This document also provides general ATM information and general product information. This document was created for users with various levels of experience. If you have any questions or problems, please contact FORE Systems' Technical Support.

## **Chapter Summaries**

**Chapter 1 - Configuring PVCs** - Describes how to create PVCs on a switch through the ATM Management Interface (AMI).

**Chapter 2 - Configuring Classical IP** - Describes how to design, configure, and maintain a Classical IP ATM network.

**Chapter 3 - Configuring an Emulated LAN** - Provides an overview of LAN Emulation and gives an example of how to configure an Emulated LAN.

**Chapter 4 - MPOA** - Contains an overview of LAN Emulation (LANE) and Multi-Protocol Over ATM (MPOA).

**Chapter 5 - ForeThought PNNI** - Provides an overview of *ForeThought* PNNI and shows how this scalable routing and signalling protocol can be used to simplify large network topologies.

**Chapter 6 - ATM Forum PNNI** - Provides an overview of ATM Forum PNNI and shows how this scalable routing and signalling protocol can be used to simplify large network topologies.

**Chapter 7 - Signalling** - Describes signalling protocol information.

Chapter 8 - Security - Describes the various forms of security that can be used on the switch.

**Chapter 9 - Configuring Timing** - Describes how to set up timing on a switch.

**Appendix A - Configuring SNMP** - Describes the remote SNMP configuration of a switch.

**Appendix B - Configuring Circuit Emulation Services** - Contains information for configuring Circuit Emulation Services (CES) network modules.

**Appendix C - Converting from FT-PNNI to PNNI** - Discusses the conversion of both non-hierarchical and hierarchical FT-PNNI networks to ATM Forum PNNI networks.

**Appendix D - Configuring FramePlus Modules** - Contains information for configuring *FramePlus* network modules.

## **Technical Support**

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1. Select the "Support" link from FORE's World Wide Web page:

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- A support contract ID number
- The serial number of each product in question
- All relevant information describing the problem or question

## **Typographical Styles**

Throughout this manual, all specific commands meant to be entered by the user appear on a separate line in bold typeface. In addition, use of the Enter or Return key is represented as **ENTER>**. The following example demonstrates this convention:

#### cd /usr <ENTER>

File names that appear within the text of this manual are represented in the following style: "...the fore\_install program installs this distribution."

Command names that appear within the text of this manual are represented in the following style: "...using the flush-cache command clears the bridge cache."

Subsystem names that appear within the text of this manual are represented in the following style: "...to access the bridge subsystem..."

Parameter names that appear within the text of this manual are represented in the following style: "...using  $\langle seg-list \rangle$  allows you to specify the segments for which you want to display the specified bridge statistics."

Any messages that appear on the screen during software installation and network interface administration are shown in Courier font to distinguish them from the rest of the text as follows:

.... Are all four conditions true?

## **Important Information Indicators**

To call your attention to safety and otherwise important information that must be reviewed to ensure correct and complete installation, as well as to avoid damage to the FORE Systems product or to your system, FORE Systems utilizes the following *WARNING/CAUTION/NOTE* indicators.

**WARNING** statements contain information that is critical to the safety of the operator and/or the system. Do not proceed beyond a **WARNING** statement until the indicated conditions are fully understood or met. This information could prevent serious injury to the operator, damage to the FORE Systems product, the system, or currently loaded software, and is indicated as follows:

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Hazardous voltages are present. To reduce the risk of electrical shock and danger to personal health, follow the instructions carefully.

**CAUTION** statements contain information that is important for proper installation/operation. Compliance with **CAUTION** statements can prevent possible equipment damage and/or loss of data and are indicated as follows:

#### **CAUTION**



You risk damaging your equipment and/or software if you do not follow these instructions.

**NOTE** statements contain information that has been found important enough to be called to the special attention of the operator and is set off from the text as follows:



If you change the value of the LECS control parameters while the LECS process is running, the new values do not take effect until the LECS process is stopped, and then restarted.

## **Laser Radiation Notice**

Class 1 Laser Product: This product conforms to applicable requirements of 21 CFR 1040 at the date of manufacture.

Class 1 lasers are defined as products which do not permit human access to laser radiation in excess of the accessible limits of Class 1 for applicable wavelengths and durations. These lasers are safe under reasonably foreseeable conditions of operation. Do not view beam with optical instruments.

Single mode fiber optic network modules contain Class 1 lasers.



This Laser Notice section only applies to products or components containing Class 1 lasers.

## **Safety Precautions**

For your protection, observe the following safety precautions when setting up equipment:

- Follow all warnings and instructions marked on the equipment.
- Ensure that the voltage and frequency of your power source matches the voltage and frequency inscribed on the equipment's electrical rating label.
- Never push objects of any kind through openings in the equipment. Dangerous
  voltages may be present. Conductive foreign objects could produce a short circuit
  that could cause fire, electric shock, or damage to your equipment.

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### Preface

# CHAPTER 1

## **Configuring PVCs**

To establish a permanent communication link between two sites, it is necessary to establish permanent virtual circuits (PVCs) at each switch along the communications path. This chapter discusses the creation of PVCs through the ATM Management Interface (AMI), a command-line user interface to *ForeRunner* ATM switches.

## 1.1 General Concepts

Each ATM cell contains a virtual path identifier (VPI) and a virtual channel identifier (VCI) as part of its five-byte ATM header. The VPI and VCI are used to route the cell through the ATM network. When a switch fabric receives a cell, it examines the ATM header to determine the correct output port, VPI, and VCI for the cell. For example, an ATM switch fabric can be configured such that any cell received on port A1 with VPI |VCI| = 0 32 is switched to port B2 with VPI |VCI| = 0 35. The translation from input port, VPI, and VCI to output port, VPI, and VCI is achieved via a mapping table in the switch fabric's memory.

The VCI value of cells does not change as the cell is switched through the ATM network via a virtual path. In a single switch environment, a cell's VPI and VCI are translated only once, but in a multiple switch environment a cell's VPI and VCI are translated many times. It is important to remember that a cell's VPI and VCI are of local significance only (i.e., link-by-link). It is also important to note that virtual connections are unidirectional; that is, they are valid in one direction only. The VPI and VCI may change as the cell is switched through the network.

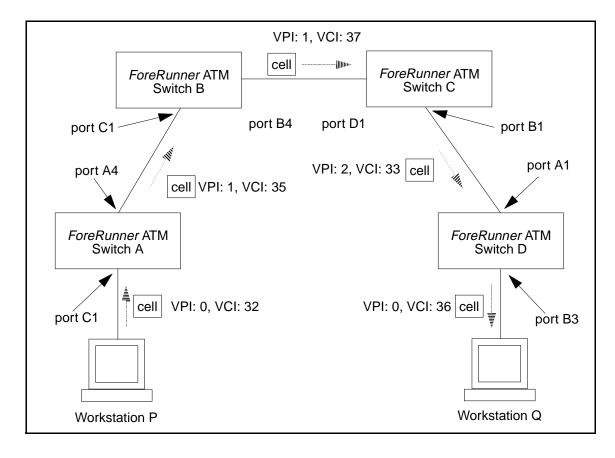


Figure 1.1 - The Cell

The mappings in an ATM network used to route cells from a source to a destination are generally referred to as virtual channels and virtual paths. The following sections is to explain how to create the necessary mappings to establish these virtual paths and virtual channels in a network of *ForeRunner* ATM switches.

### 1.2 Virtual Paths

Virtual paths, which are carried within a physical transit medium (e.g., DS1, E1, DS3, E3, OC3c, or OC12c link), are used to establish connections between two nodes in an ATM network. Many virtual paths can be transmitted within a single physical link. Two types of virtual paths exist: virtual path connections (VPCs), also known as through paths, and originating/terminating paths, also known as virtual path terminators (VPTs). VPCs allow virtual paths to be cross-connected at a switch node while VPTs allow virtual channels (VCCs) to be cross-connected or switched at a switch node.

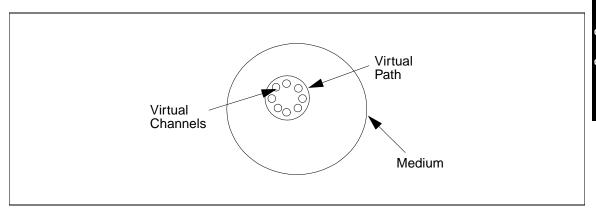


Figure 1.2 - Virtual Channels in a Virtual Path

A single virtual path can be used to route many virtual channels through the ATM network. Because a virtual path simply routes virtual channels through the network, a cell is guaranteed to have the same VCI when it exits the virtual path as it had when it entered the virtual path.

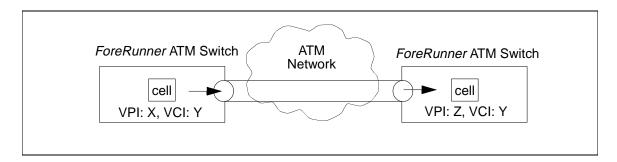


Figure 1.3 - An Example of a Virtual Path

The VCI value of cells does not change as the cell is switched through the ATM network via a virtual path. Each virtual path must originate at a switch fabric, pass through zero or more switch fabrics and terminate at another switch fabric. The origination and termination points are referred to as originating and terminating paths. Virtual paths are switched through switch fabrics via through paths. Virtual paths are made up of an originating path, zero or more through paths, and a terminating path.

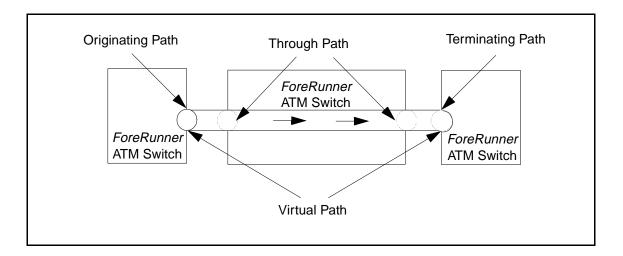


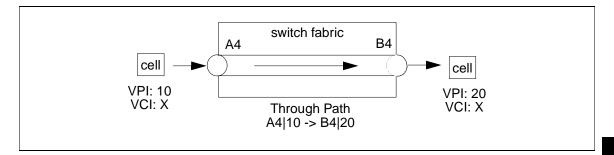
Figure 1.4 - Composition of a Virtual Path

## 1.2.1 Through Paths

Through paths route an entire virtual path through an ATM switch fabric. When a cell is received by a switch fabric on a through path, the VPI is examined to determine the output port and VPI. The VCI component of the ATM header remains unchanged and can have any value. So, all of the channels within the through path are switched correctly without altering the VCI value of cells on these channels.

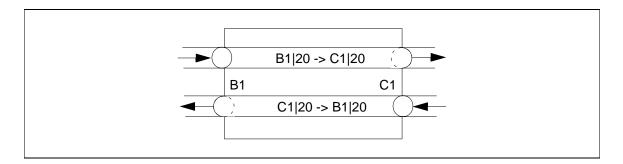
Four parameters are needed to define a through path on a *ForeRunner* switch fabric: input port, input VPI, output port, and output VPI. Through paths are represented as follows:

The VCI value remains unchanged when cells are switched via a through path. For example, the through path A4  $\mid$  10 -> B4  $\mid$  20 maps cells received on port A4 with VPI: 10 and any VCI to port B4 with VPI: 20 and the same VCI.



**Figure 1.5 -** An Example of a Through Path

By definition, through paths only switch cells in one direction; they are unidirectional. For example, switch fabric X is configured with the through path B1 | 20 -> C1 | 20. If a cell is received on port C1 with VPI: 20, it is not transmitted on port B1 with a new VPI: 20. In order for this to happen, the through path C1 | 20 -> B1 | 20 must exist as well. Since through paths are unidirectional, two through paths are necessary for bidirectional communication.



**Figure 1.6 -** Through Paths are Unidirectional

## 1.2.2 Originating and Terminating Paths

As previously noted, originating and terminating paths (also called virtual path terminators) are points at which a virtual path originates and terminates. For example, if a virtual path exists from switch fabric A to switch fabric B, then there must be an originating path on switch fabric A and a terminating path on switch fabric B.

An originating path is defined by two parameters: output VPI and output port. Similarly, a terminating path is defined by the parameters: input VPI and input port. Because originating and terminating paths do not define the way cells are switched through an ATM switch fabric, virtual channels must exist to switch cells from a terminating path to an originating path. (See the section about virtual channels for more information.) Originating and terminating paths are the endpoints of virtual paths and are used primarily for bandwidth allocation.

The bandwidth allocated to originating and terminating paths is used to control the amount of virtual channel (VCC) bandwidth entering or leaving a virtual path. The total guaranteed bandwidth used by virtual channels on an originating path or a terminating path cannot exceed the amount of bandwidth allocated to that path. For example, as illustrated in Figure 1.7, if each of the four virtual channels shown is using 10 Mbps of bandwidth, then the originating and terminating paths must have at least 40 Mbps of bandwidth allocated.



UBR traffic bandwidth, which is a "best effort" service class, is not limited by the VP's allocated bandwidth since its bandwidth is not guaranteed. Actual UBR VCC traffic transmitted within a VP may exceed the VP's allocated bandwidth.

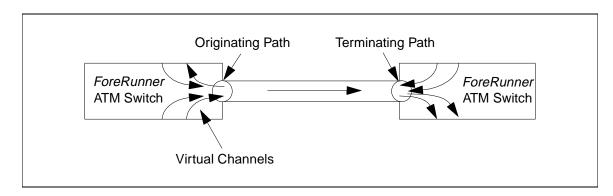


Figure 1.7 - Using Originating and Terminating Paths for Bandwidth Allocation

# 1.3 Listing Virtual Paths

# 1.3.1 Listing Through Paths

By logging in to AMI (see Chapter 1 of the *ATM Management Interface (AMI) Manual* for information about logging into AMI), it is possible to display either all of the existing through paths on an individual switch fabric or all of the existing through paths on a specified port. To list all of the existing through paths on an individual switch fabric, enter the following:

#### configuration vpc show

Input		Output					
	Port	VPI	Port	VPI	UPC	Prot	Name
	3B1	40	3B4	40	0	pvc	customer_a
	3B1	75	3B5	75	0	pvc	customer_b
	3B2	95	3B3	95	0	pvc	customer_e
	3B6	62	3B2	62	0	pvc	customer_c
	3B6	68	3B3	68	0	pvc	customer_d

The fields in this display are defined as follows:

Field	Description		
Input Port	The incoming port number of the through path.		
Input VPI	The incoming virtual path number.		
Output Port	Port The outgoing port number of the through path.		
Output VPI	The outgoing virtual path number.		
UPC	The integer index that refers to a specific UPC traffic contract assigned to this through path. UPC contracts can be displayed using conf upc show		
Prot	The type of protocol running on this channel.		
Name	The user-assigned name which helps to identify this through path uniquely.		

# Configuring PVCs

To list advanced options about all of the existing virtual (through) paths, enter the following parameters:

### configuration vpc show advanced

Input		Outpu	.t		
Port	VPI	Port	VPI	Shape	ConType
3B1	40	3B4	40		N/A
3B1	75	3B5	75		N/A
3B2	95	3B3	95		tran-tran-pmp
3B6	62	3B2	62		tran-tran-pp
3B6	68	3B3	68		N/A

The fields in the advanced display are defined as follows:

Field	Description
Input Port	The incoming port number of the through path.
Input VPI	The incoming virtual path number.
Output Port	The outgoing port number of the through path.
Output VPI	The outgoing virtual path number.
Shape	Indicates whether or not traffic shaping has been enabled for this path. This field only applies to the Series C network modules.
СопТуре	The connection type for the endpoints of this path with respect to a particular network. Orig (originating) means that the ingress/egress endpoint of the path is connected to the source node which is outside the network, tran (transit) means that the ingress/egress endpoint of the path is connected to a node within the network, and term (terminating) means that the ingress/egress endpoint of the path is connected to the destination node which is outside the network. pp means this is labelled as a point-to-point path, pmp means this is labelled as a multipoint-to-point path. mpmp means this is labelled as a multipoint-to-multipoint path.

# 1.3.2 Listing Originating and Terminating Paths

By logging in to AMI, it is possible to display either all of the existing originating and terminating paths on an individual switch fabric or on a specified port. To list all of the originating and terminating paths on an individual switch fabric, enter the following parameters:

#### configuration vpt show

t Output		t						
VPI	Port	VPI	ResBW	CurBW	MinVCI	MaxVCI	VCs	Protocol
0	termi	nate	N/A	0.8K	1	511	6	pvc
1	terminate		1.0M	0.8K	1	511	2	pvc
0	terminate		N/A	0.8K	1	511	6	pvc
0	termi	nate	N/A	0.8K	1	511	6	pvc
0	terminate		N/A	0.8K	1	511	6	pvc
0	termi	nate	N/A	7.6K	1	1023	24	pvc
nate	1C1	0	N/A	0.8K	1	511	6	pvc
nate	1C1	1	1.0M	0.8K	1	511	2	pvc
nate	1C2	0	N/A	0.8K	1	511	6	pvc
nate	1C3	0	N/A	0.8K	1	511	6	pvc
nate	1C4	0	N/A	0.8K	1	511	6	pvc
nate	1CTL	0	N/A	7.6K	1	1023	30	pvc
	VPI 0 1 0 0 0 0 nate nate nate	VPI Port 0 termi 1 termi 0 termi 0 termi 0 termi 1 ter	VPI Port VPI 0 terminate 1 terminate 0 terminate 0 terminate 0 terminate 1 ter	VPI Port VPI ResBW  0 terminate N/A  1 terminate 1.0M  0 terminate N/A  0 terminate N/A  0 terminate N/A  0 terminate N/A  nate 1C1 0 N/A  nate 1C1 1 1.0M  nate 1C2 0 N/A  nate 1C3 0 N/A  nate 1C4 0 N/A	VPI Port VPI ResBW CurBW 0 terminate N/A 0.8K 1 terminate 1.0M 0.8K 0 terminate N/A 0.8K nate 1C1 0 N/A 0.8K nate 1C1 1 1.0M 0.8K nate 1C2 0 N/A 0.8K nate 1C3 0 N/A 0.8K nate 1C4 0 N/A 0.8K	VPI Port VPI ResBW CurBW MinVCI 0 terminate N/A 0.8K 1 1 terminate 1.0M 0.8K 1 0 terminate N/A 0.8K 1 nate 1C1 0 N/A 0.8K 1 nate 1C2 0 N/A 0.8K 1 nate 1C3 0 N/A 0.8K 1 nate 1C4 0 N/A 0.8K 1 nate 1C4 0 N/A 0.8K 1	VPI         Port         VPI         ResBW         CurBW         MinVCI         MaxVCI           0         terminate         N/A         0.8K         1         511           1         terminate         1.0M         0.8K         1         511           0         terminate         N/A         0.8K         1         511           nate         1C1         0         N/A         0.8K         1         511           nate         1C2         0         N/A         0.8K         1         511           nate         1C3         0         N/A         0.8K         1         511           nate         1C4         0         N/A         0.8K         1         511	VPI         Port         VPI         ResBW         CurBW         MinVCI         MaxVCI         VCs           0         terminate         N/A         0.8K         1         511         6           1         terminate         1.0M         0.8K         1         511         2           0         terminate         N/A         0.8K         1         511         6           0         terminate         N/A         0.8K         1         511         6           0         terminate         N/A         7.6K         1         1023         24           nate         1C1         0         N/A         0.8K         1         511         6           nate         1C1         1         1.0M         0.8K         1         511         2           nate         1C2         0         N/A         0.8K         1         511         6           nate         1C3         0         N/A         0.8K         1         511         6           nate         1C4         0         N/A         0.8K         1         511         6

The fields in this display are defined as follows:

Field	Description
Input Port	The incoming port number of the vpt. Shows originate if it is an originating path.
Input VPI	The incoming virtual path number.
Output Port	The outgoing port number of the vpt. Shows the number of the output port of the vpt. Shows terminate if it is a terminating path.
Output VPI	The outgoing virtual path number.
ResBW	The maximum amount of bandwidth, in Kbps, that is reserved for the virtual channels using this vpt. A value of $N/A$ indicates that this path is an elastic path. Elastic paths allocate and de-allocate bandwidth for their channels from the link.
CurBW	The amount of bandwidth, in Kbps, being used by the virtual channels using this vpt.
MinVCI	The bottom number for the range of VCIs that are reserved for VCCs on this virtual path terminator. The default is 1.
MaxVCI	The top number for the range of VCIs that are reserved for VCCs on this virtual path terminator. The default is 511.
VCs	The number of virtual channels that are currently using this vpt.
Protocol	The type of protocol running on this channel.

To list all of the advanced options about all of the existing virtual path terminators, enter the following parameters:

#### configuration vpt show advanced

Input		Output					
Port	VPI	Port	VPI	Shape	VBROB	BuffOB	
1C1	0	termi	nate	N/A	N/A	N/A	
1C1	1	termi	nate	N/A	N/A	N/A	
1C2	0	termi	nate	N/A	N/A	N/A	
1C3	0	termi	nate	N/A	N/A	N/A	
1C4	0	terminate		N/A	N/A	N/A	
1CTL	0	terminate		N/A	N/A	N/A	
origi	nate	1C1 0			port	port	
origi	nate	1C1	1		100	100	
originate		1C2	0	) port		port	
origi	nate	1C3 0		port		port	
origi	nate	1C4 0			port		
origi	nate	1CTL	0		N/A	N/A	

The fields in the advanced display are defined as follows:

Field	Description				
Input Port	The incoming port number of the vpt. Shows originate if it is an originating path.				
Input VPI	The incoming virtual path number.				
Output Port	The outgoing port number of the vpt. Shows terminate if it is a terminating path.				
Output VPI	The outgoing virtual path number.				
Shape	Indicates whether or not traffic shaping has been enabled for this originating vpt. field only applies to the Series C network modules.				
VBROB	The bandwidth overbooking level assigned to this vpt, specified as a percentage. The default is 100, which means that no overbooking has been defined. Values less than 100 cause underbooking. Values greater than 100 denote overbooking. port means this is an elastic path. Since elastic paths derive their overbooking factors from their parent ports, use conf port show to display the overbooking value.				
BuffOB	The buffer overbooking level assigned to this vpt, specified as a percentage. The default is 100, which means that no overbooking has been defined. Values less than 100 cause underbooking. Values greater than 100 denote overbooking. port means this is an elastic path. Since elastic paths derive their overbooking factors from their parent ports, use conf port show to display the overbooking value.				



For more information about setting overbooking parameters (VBROB and BuffOB), see Section 1.5.2.

## 1.4 Virtual Channels

Virtual channels "ride" inside of virtual paths. The combination of the two specifies a virtual connection. On a switch fabric, each virtual channel switches cells with a specific VPI and VCI received on a specific port to another port with a new VPI and a new VCI. Unlike through paths, which carry one or more VCCs, virtual channels describe a single virtual connection between two endpoints.

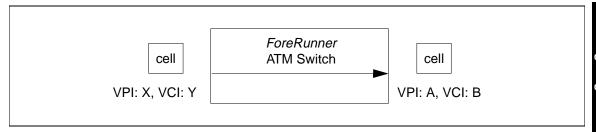


Figure 1.8 - An Example of a Virtual Channel

Six parameters are needed to define a virtual channel: input port, input VPI, input VCI, output port, output VPI, and output VCI. Virtual channels are represented by the following notation:

```
<iport> <ivpi> <ivci> <oport> <ovpi> <ovci>
```

Virtual channels switch cells using both the VPI and VCI values. Both the VPI and VCI values may change when a cell is switched via a virtual channel. For example, the virtual channel  $C2 \mid 1 \mid 20 \rightarrow D2 \mid 9 \mid 25$  switches cells received on port C2 with VPI: 1 and VCI: 20 such that they are transmitted out port D2 with VPI: 9 and VCI: 25.

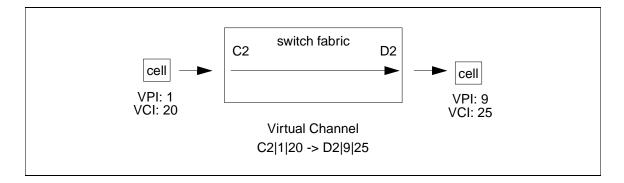


Figure 1.9 - Example of a Virtual Channel

In order to establish two-way communications between two ports on a switch fabric, two virtual channels are necessary because virtual channels are unidirectional. For example, switch fabric A is configured with the virtual channel C3 | 7 | 12 -> D1 | 8 | 2. If a cell is received on port D1 with VPI: 8 and VCI: 2, it is not transmitted out port C3 with VPI: 7 and VCI: 12. An additional channel, namely D1 | 8 | 2 -> C3 | 7 | 12, would have to exist.

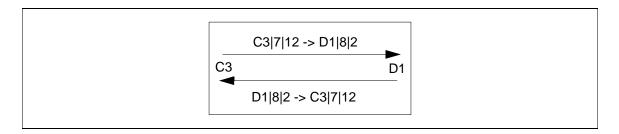


Figure 1.10 - Virtual Channels are Unidirectional

Before a virtual channel can be created, the corresponding terminating and originating paths must exist. For example, before the channels shown on the switch fabric in Figure 1.11 can be created, the terminating path C3  $\mid$  3 must exist.

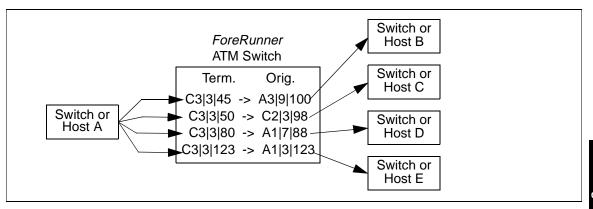
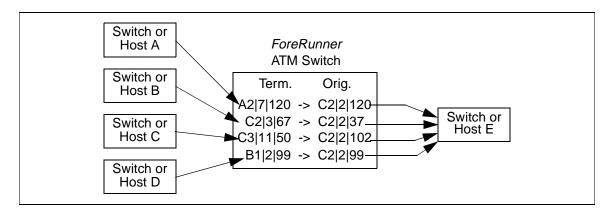


Figure 1.11 - Virtual Channels Created on Terminating Path C3 | 3

Similarly, before the virtual channels shown in Figure 1.12 can be created, the originating path  $C2 \mid 2$  must exist.



**Figure 1.12 -** Virtual Channels Created on Originating Path  $C2 \mid 2$ 

Furthermore, in these examples, the terminating path  $C3 \mid 3$  and originating path  $C2 \mid 2$  must have enough bandwidth allocated to support the total bandwidth used by the virtual channels (see Figure 1.7).

### 1.4.1 Smart Permanent Virtual Circuits

Smart Permanent Virtual Circuits (SPVCs) are connections that go across multiple switch fabrics. An SPVC looks like a PVC at the local and remote endpoints with an SVC (Switched Virtual Circuit) in the middle. SVCs are channels established on demand by network signalling. Similar to a dialed telephone call, SVCs transport information between two locations and last only for the duration of the transfer.

SPVCs are more robust than PVCs. If a link carrying a PVC goes down, then the PVC goes down. If a link carrying a SPVC goes down and there is an alternate route, then the end switch fabrics of the SPVC automatically reroute the SPVC around the failed link.

As shown in Figure 1.13, interswitch links exist between the *ForeRunner* switches. The endpoints exist at switch A and switch E. If a link goes down between switch A and switch B, an SPVC can reroute the cell (via an SVC) through switch C.

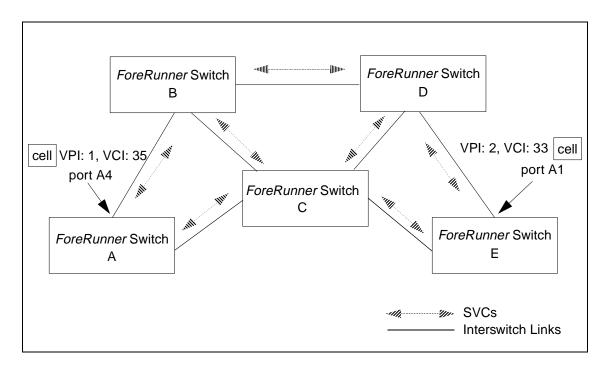


Figure 1.13 - The Path of a Cell Via SPVCs

# 1.4.2 Listing Virtual Channels

By logging in to AMI, you can display either all of the existing virtual channels on an individual switch fabric or on a specified port. To list all of the virtual channels on an individual switch fabric, enter the following parameters:

#### configuration vcc show

Input		Ou	tput					
Port	VPI	VCI	Port	VPI	VCI	UPC	Protocol	Name
3B1	0	5	3CTL	0	49	0	uni	N/A
3B1	0	14	3CTL	0	48	0	spans	N/A
3B1	0	15	3CTL	0	47		spans	N/A
3B1	0	16	3CTL	0	50		uni	N/A
3B1	0	100	3B4	0	100	0	pvc	N/A
3B2	0	5	3CTL	0	53	0	uni	N/A
3B2	0	14	3CTL	0	52	0	spans	N/A
3B2	0	15	3CTL	0	51		spans	N/A

Press return for more, q to quit: q

### The fields in this display are defined as follows:

Field	Description
Input Port	The incoming port number of the virtual channel.
Input VPI	The incoming virtual path number.
Input VCI	The incoming virtual channel number.
Output Port	The outgoing port number of the virtual channel.
Output VPI	The outgoing virtual path number.
Output VCI	The outgoing virtual channel number.
UPC	The integer index that refers to the specific UPC traffic contract assigned to this VCI.
Protocol	Indicates what type of channel this is. Can be spans, pvc, uni, spvc, or rcc. rcc is the routing control channel (0, 18) on PNNI links over which PNNI exchanges routing information.
Name	The unique, user-assigned name for this channel. If no name is assigned, shows N/A.

# Configuring PVCs

To list advanced information about all of the existing permanent virtual channels on a switch board, enter the following parameters:

#### configuration vcc show advanced

Input 0		Ou	tput				
Port	VPI	VCI	Port	VPI	VCI	Protocol	ConType
3B1	0	5	3CTL	0	49	uni	N/A
3B1	0	14	3CTL	0	48	spans	N/A
3B1	0	15	3CTL	0	47	spans	N/A
3B1	0	16	3CTL	0	50	uni	N/A
3B1	0	100	3B4	0	100	pvc	tran-tran-pp
3B2	0	5	3CTL	0	53	uni	N/A
3B2	0	14	3CTL	0	52	spans	N/A
3B2	0	15	3CTL	0	51	spans	N/A
3B2	0	16	3CTL	0	54	uni	N/A
Press	retur	n for	more,	q to	quit	: q	

The fields in the advanced display are defined as follows:

Field	Description			
Input Port	The incoming port number of the virtual channel.			
Input VPI	The incoming virtual path number.			
Input VCI	The incoming virtual channel number.			
Output Port	The outgoing port number of the virtual channel.			
Output VPI	The outgoing virtual path number.			
Output VCI	The outgoing virtual channel number.			
Protocol	Indicates what type of channel this is. Can be spans, pvc, uni, spvc, or rcc. rcc is the routing control channel (0, 18) on PNNI links over which PNNI exchanges routing information.			
СопТуре	The connection type for the endpoints of this channel with respect to a particular network. Orig (originating) means that the ingress/egress endpoint of the channel is connected to the source node which is outside the network, tran (transit) means that the ingress/egress endpoint of the channel is connected to a node within the network, and term (terminating) means that the ingress/egress endpoint of the channel is connected to the destination node which is outside the network. pp means this is labelled as a point-to-point channel, pmp means this is labelled as a multipoint-to-point channel. mpmp means this is labelled as a multipoint-to-multipoint channel.			

# 1.5 Creating PVCs and SPVCs

FORE's ATM network modules provide ATM transmission connectivity, while its intelligent network modules, such as the Circuit Emulation Services (CES) network module, provide adaptation for ports carrying one transmission format (e.g., TDM) to ATM cells. This section describes how to create the following from one ATM port to another:

- · a permanent virtual path (through path)
- a permanent virtual path terminator (originating or terminating path)
- · a permanent virtual channel through the network
- · a smart permanent virtual channel through the network

This section assumes that the physical port parameters of the switches have already been configured and that the ATM network module traffic models have been set appropriately. For more information about these configurations, see the ATM Management Interface (AMI) Manual.



When these paths and channels are created, a command is entered automatically into the current configuration database (CDB), meaning that these paths and channels are created every time the switch control processor (SCP) is restarted. The CDB should be backed up frequently.

# 1.5.1 Creating a Through Path

To create a new through path, log in to AMI and enter the following parameters:

The optional parameters for call records and VP shaping (using Series C network modules) are as follows:

These parameters are defined as follows:

Parameter	Description			
iport	The incoming port number.			
ivpi	The incoming virtual path number.			
oport	The outgoing port number.			
ovpi	The outgoing virtual path number.			
-upc <index></index>	The integer index that refers to a specific UPC traffic contract. If no index is specified, then no traffic policing will take place on this VPI. It is assigned a UPC index of 0, and all traffic on this VPI is treated as UBR traffic. This is the default.			
-name <name></name>	The name you want to assign to this through path to help identify it uniquely. It is most useful for billing purposes so you can identify which paths are being used by which customers. Can be up to 32 ASCII characters long.			
-inctype (orig   tran   term)	The path connection type for the incoming path. For billing purposes, it denotes on which switch this path is arriving. Orig (originating) means that the ingress endpoint of the path is connected to the source node which is outside the network, tran (transit) means that the ingress endpoint of the path is connected to a node within the network, and term (terminating) means that the ingress endpoint of the path is connected to the destination node which is outside the network.			
-outctype(orig   tran   term)	The path connection type for the outgoing path. For billing purposes, it denotes on which switch this path is leaving. Orig (originating) means that the egress endpoint of the path is connected to the source node which is outside the network, tran (transit) means that the egress endpoint of the path is connected to a node within the network, and term (terminating) means that the egress endpoint of the path is connected to the destination node which is outside the network.			
pmp <sup>1</sup>	Indicates this is a point-to-multipoint path.			
трр	Indicates this is a multipoint-to-point path.			
тртр	Indicates this is a multipoint-to-multipoint path.			

Parameter	Description
-shapeivpi <vpi><sup>2</sup></vpi>	The incoming VPI for this through path. When the traffic shaping port is not the port connected to the WAN, a through path must be created from the WAN port to the traffic shaping port. Cells arrive from the network at the traffic shaping port with this value equal to the VPI of the terminating path at the traffic shaping port. This parameter only applies to the Series C network modules.

<sup>1.</sup> By indicating pmp, mpp, or mpmp, you are only assigning a label for record keeping purposes. The switch does not necessarily create the type of path you have specified. If you assign a connection type, but do not assign a pmp, mpp, or mpmp label, the switch assigns a label of pp (point-to-point).

<sup>2.</sup> If you want to shape traffic on more than two ports on a given Series C network module, it is recommended that you set the traffic memory model to model number 5 for that network module using conf module traffic c setmodel.



Terminating and originating paths cannot be created across the intra-fabric ports on an ASX-1000; only through paths can be created across the intra-fabric ports as shown in the third example.

The following is an example of how to create a virtual path which specifies a name:

```
myswitch::configuration vpc> new 3b1 75 3b5 75 -name customer_b
```

The following is an example of how to create a virtual path which specifies a name and a connection type:

```
myswitch::configuration vpc> new 3b6 62 3b2 62 -name customer_c -inctype tran
-outctype tran
```

The following is an example of how to create a virtual path on an ASX-1000. To create a through path going in port 2A1, VPI 1 on the switch board installed in slot 2 and going out port 4B1, VPI 1 on the switch board installed in slot 4, enter the following:

```
myswitch::configuration vpc> new 2a1 1 2e4 1
myswitch::configuration vpc> new 2e4 1 2a1 1
myswitch::configuration vpc> new 4b1 1 4e2 1
myswitch::configuration vpc> new 4e2 1 4b1 1
```

#### Configuring PVCs

In the first line in the first pair, notice that the output port is 2E4. This is the intra-fabric port. The 2 means the connection is coming out of the switch board in slot 2 through the intra-fabric port. The E represents the intra-fabric port. The 4 means the connection is destined for switch board in slot 4. 2E4 then becomes the input port in the second line.

In the first line in the second pair, notice that the output port is 4E2. This is the intra-fabric port. The 4 means the connection is coming out of the switch board in slot 4 through the intra-fabric port. The E represents the intra-fabric port. The 2 means the connection is destined for switch board in slot 2. 4E2 then becomes the input port in the second line.

At the same time, a command is entered automatically into the current configuration database, which means that this virtual path is created every time the SCP is restarted.

# 1.5.2 Creating an Originating or Terminating Path

To create an originating or terminating path (virtual path terminator), log in to AMI and enter the following parameters:

The **-reserved**, **-minvci**, and **-maxvci** parameters are optional for creating originating or terminating paths. The advanced traffic management options for creating originating paths are as follows:

```
[-shapeovpi <vpi>] [-loopvpi <vpi>]
[-vbrob <percent>] [-vbrbuffob <percent>]
```

The advanced QoS options for creating originating or terminating paths are as follows:

```
[-cbr (none | default | <qosindex>)]
[-rtvbr (none | default | <qosindex>)]
[-nrtvbr (none | default | <qosindex>)]
[-ubr (none | default | <qosindex>)]
[-abr (none | default | <qosindex>)]
```



The <qosindex> must exist (conf qosext new) before it can be applied to the originating/terminating path.

These parameters are defined as follows:

Parameter	Description			
port	The port number for this vpt.			
vpi	The path number for this vpt.			
term	Specifies that the vpt to be created is a terminating path.			
orig	Specifies that the vpt to be created is an originating path.			
reserved	The amount of bandwidth, in Kbps, that you want to reserve on this vpt. If this option is not used, an elastic path is created. Elastic paths allocate and de-allocate bandwidth for their channels from the link.			
minvci	The bottom number for the range of VCIs to be reserved for VCCs on this vpt. The default is 1.			
maxvci	The top number for the range of VCIs to be reserved for VCCs on this vpt. The default is 511.			

Parameter	Description			
shapeovpi	The output path on a traffic shaping originating vpt. Setting this value configures traffic shaping on the originating path. Cells bound for the network leave the traffic shaping port with this VPI. When the traffic shaping port is the WAN port, this value equals the input VPI of the originating path. If the traffic shaping port is not the WAN port, this value equals the input VPI of the through path from the shaping port to the WAN port. This parameter only applies to the Series C network modules. See Section 1.5.2.1.			
loopvpi	The originating vpi will be shaped by a through path going to a Series D network module. You should enter the input vpi of the through path that goes from the looping port to the WAN port. This option is also used when creating the through path that connects from the WAN port to the looping port. The through path loopvpi should be the same vpi as the terminating path on the looping port. See Section 1.5.2.1 for more information.			
vbrob	The bandwidth overbooking level for this vpt, specified as a percentage. Valid values are integers from 1 to 32,767. 100 means that no overbooking has been defined. Values less than 100 cause underbooking. Values greater than 100 cause overbooking. Overbooking cannot be specified on an elastic path. Therefore, you can only specify an overbooking factor for an originating path when you also have reserved bandwidth for the path (i.e., specified the -reserved <kbs> parameter).</kbs>			
vbrbuffob	The buffer overbooking level for this vpt, specified as a percentage. Valid values are integers greater than or equal to 1. 100 means that no overbooking has been defined. Values less than 100 cause underbooking. Values greater than 100 cause overbooking. Overbooking cannot be specified on an elastic path. Therefore, you can only specify an overbooking factor for an originating path when you also have reserved bandwidth for the path (i.e., specified the -reserved <kbs> parameter).</kbs>			
none	The specified class of service (CBR, real-time VBR, non real-time VBR, UBR, ABR) is not supported.			
default	The default parameters of 0 CTD, 0 CDV, and 0 CLR are to be used for the CBR class of service.			
qosindex	The index of the set of QoS extension parameters. See conf qosext show for this number.			



Bandwidth and/or buffer overbooking are used when the number of VBR VCCs configured across an originating or terminating VPT exceeds the guaranteed capacity of the VP. While overbooking removes service guarantees (in the case when all users transmit at full contract rate simultaneously), a network can be engineered in a more cost-effective manner if it is assumed that all VCC sources do not transmit simultaneously.

**Guideline:** An initial buffer overbooking value of 500-700% is recommended until live traffic patterns and VPC utilization can be measured. If bandwidth is still under-utilized, a higher VBROB setting can be used.

The following is an example of how to create a terminating path:

```
myswitch::configuration vpt> new 3b3 99 term
Would you like to create the originating side also [y]? y
```

The following is an example of how to create a originating path:

```
myswitch::configuration vpt> new 3b4 88 orig
Would you like to create the terminating side also [y]? y
```

The following is an example of how to delete a terminating path:

```
myswitch::configuration vpt> del 3b4 88 term Would you like to delete the originating side also [y]? y
```

The following is an example of how to delete an originating path:

```
myswitch::configuration vpt> del 3b3 99 orig Would you like to delete the terminating side also [y]? y
```

If you do not specify term or orig, the switch automatically deletes both sides of the path:

```
myswitch::configuration vpt> del 3b4 88
```



Before deleting a virtual path, you must first delete all VCCs which use that path.

### 1.5.2.1 Shaping Multiple Originating Paths on a Single Port

This feature allows you to shape several originating paths to be output on a single port. This feature is useful if you have several remote sites interconnected by a PVP mesh. If you only need to shape one originating path, you can simply use the conf port traffic d ratelimit command on a Series D network module.

This feature requires the use of two ports. One port, which can be on a Series C, Series LC, or Series D network module, is used for originating and terminating path(s). This port is put in either physical or diagnostic loopback and is configured with the new option -loopvpi under conf vpt new. The other port provides the actual output on a shaped through path, so this port must be on a Series D network module.

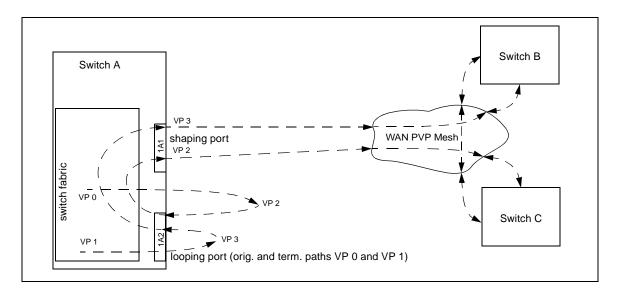


Figure 1.14 - PVPs Looped through Port 1A2 and Output on Port 1A1 to WAN

As shown in the example in Figure 1.14, the traffic that is going out originating path 0 on port 1A2 gets looped so that it is output to the WAN on through path 2 on shaping port 1A1. Similarly, the traffic that is going out originating path 1 on port 1A2 gets looped so that it is output to the WAN on through path 3 on shaping port 1A1.

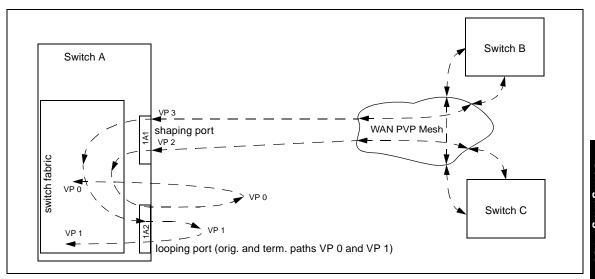


Figure 1.15 - PVPs Coming in Port 1A1 from WAN and Looped through Port 1A2

Then, as shown in the example in Figure 1.15, the traffic that is coming in through path 2 on port 1A1 from the WAN gets looped back to terminating path 0 on on port 1A2. Similarly, the traffic that is coming in through path 3 on port 1A1 from the WAN gets looped back to terminating path 1 on port 1A2.

In your own network, you will repeat this process on these same two ports for as many paths as you want to shape.



Before you configure the paths, you should configure the network module that will contain the looping port to maximize the number of multicast connections. On a Series C network module, use model number 5 under conf module traffic c setmodel. On a Series LC network module, use model number 6, 7, or 8 under conf module traffic lc setmodel.

To configure the shaped originating paths as shown in the example in Figure 1.14 and Figure 1.15, perform the following steps:

1. First, create the UPC contract to be used on the through paths that are going to be output on the shaping port on the Series D network module to the WAN.

myswitch::configuration upc> new 2 cbr 42452 noGCRA -scheduling smoothed



You must use the -nogCRA option to disable policing. This allow bursts without dropping cells. The shaped rate is equal to PCR for CBR traffic and equal to SCR for VBR traffic. You must also use the smoothed scheduling option so that a rate group is created when a PVC or PVP is configured with this UPC.

2. Next, delete any existing path(s) on the looping port as follows:

```
myswitch::configuration upc> conf vpt
myswitch::configuration vpt> conf vpt del 1a2 0
```



By default, path 0 is the only path that exists. If path 1 already exists, you must delete it using the command conf vpt del 1a2 1.

3. Recreate the paths on the looping port using the **-loopvpi** option to ensure that cells from the WAN port get looped back into the terminating path.

```
myswitch::configuration vpt> new 1a2 0 -loopvpi 2 myswitch::configuration vpt> new 1a2 1 -loopvpi 3
```

4. Create a PVP to the shaping port on a Series D network module and apply the UPC contract that you created in step 1 so that it shapes the cells.

```
myswitch::configuration vpc> conf vpc
myswitch::configuration vpc> new la2 2 la1 2 -upc 2
myswitch::configuration vpc> new la2 3 la1 3 -upc 2
```

5. Create a PVP in the opposite direction and use **-loopvpi** to ensure that the cells on each through path get looped back the terminating paths on the looping port (1a2).

```
myswitch::configuration vpc> new 1a1 2 1a2 2 -loopvpi 0 myswitch::configuration vpc> new 1a1 3 1a2 3 -loopvpi 1
```

6. Now use a loopback command to loop the transmit side to the receive side on the looping port (1a2).

```
myswitch::configuration vpt> conf port admin 1a2 down
myswitch::configuration vpt> conf port sonet loop 1a2 diag
```



Any required signalling paths must be recreated at this point.

7. Additional shaped originating paths can be added later on these same two ports as needed by repeating the steps in this section.

## 1.5.2.2 Terminating a PVC at a Switch

Sometimes it is necessary to create a PVC between a host and a switch fabric that is at a remote location. In this case, the PVC should be created from the host to the control port (CTL) of the switch and vice versa.

Some additional configuration is necessary for communication to be established between the host and the switch fabric. The switch needs an entry in its ATM ARP cache in order to send cells destined for the host with the correct VPI and VCI and to pass received cells with a specific VPI and VCI to IP. This configuration can be done using AMI as shown in the following subsections. (See the atmarp (8c) man page for more information.)

### 1.5.2.3 Creating ATM ARP Entries

To create a FORE IP PVC ARP entry, log in to AMI. Data on this PVC is encapsulated using null encapsulation (also known as VC-based multiplexing) as specified in RFC-1483. Enter the following parameters:

configuration atmarp newforeip <host> <vpi> <vci> (4 | 5) [ <interface>]

These parameters are defined as follows:

Parameter	Description			
host	The IP address of the remote host.			
vpi	The virtual path number of the FORE IP PVC. Must be 0.			
vci	The virtual channel number of the FORE IP PVC.			
4   5	The connection's ATM Adaptation Layer (AAL) type. The default is 4.			
interface	The FORE IP interface to be used for this connection. The default is asx0.			

Once the parameters are entered, the entry is created instantly by the SCP. At the same time, a command is entered automatically into the current CDB which creates this ATM ARP entry each time the SCP is restarted.

To create a new Classical IP PVC ARP entry, log in to AMI. All data is sent LLC/SNAP encapsulated. Enter the following parameters:

configuration atmarp newclassicalip <host> <vpi> <vci> [<interface>]

These parameters are defined as follows:

Parameter	Description			
host	The host IP address of the remote IP endstation.			
vpi	the virtual path number of the Classical IP PVC.			
vci	The virtual channel number of the Classical IP PVC.			
interface	The Classical IP interface to be used for this connection: qaa0, qaa1, qaa2, or qaa3. The default is qaa0.			

Once the parameters are entered, the Classical IP PVC ARP entry is created instantly by the SCP. At the same time, a command is entered automatically into the current CDB which creates this ATM ARP entry each time that the SCP is restarted.

### 1.5.2.4 Listing ATM ARP Entries

To verify that the ARP entries exist correctly for the outgoing PVC connection from the SCP to the host, display the ATM ARP cache by logging in to AMI and entering the following parameters:

#### configuration atmarp show

IPaddress	If	VPI	VCI	AAL	Type	Direction
198.29.22.9	asx0	0	63	aal5	foreIpSVC	pending
198.29.22.15	asx0	0	231	aal5	foreIpSVC	pending
198.29.22.37	asx0	0	65	aal34	foreIpSVC	pending
IPaddress	If	NSAP 2	Addres	s		
198.29.17.3	qaa0	0x47.	0005.8	0.ffe10	0.0000.f21b.0138.0	02048102754.00
198.29.17.10	qaa0	0x47.	0005.8	0.ffe10	0.0000.f21b.0137.0	02048100be6.00
198.29.17.15	qaa0	0x47.	0005.8	0.ffe10	0.0000.f21b.0137.0	0204810048d.00
198.29.17.52	gaa0	0×47.	0005.8	0.ffe10	0.0000.f21b.0138.0	020481b0138.00

### The fields in this display are defined as follows:

Field	Description				
IPaddress	he IP address for this connection.				
If	ne name of the IP interface for this connection.				
VPI	The virtual path number.				
VCI	The virtual channel number.				
AAL	The AAL type of the given connection.				
Туре	Shows what kind of connection this is. Can be foreIpPVC, foreIpSVC, classicalIp-PVC, or classicalIpSVC.				
Direction	Outgoing means this is an outgoing connection. Incoming means this is an incoming connection. Pending means that a connection has not (yet) been established. Incomplete means that the IP-to-ATM address mapping is not yet known for the given IP address.				
NSAP Address	The NSAP address for this connection.				

# 1.5.3 Creating a Virtual Channel

To create a new virtual channel, log in to AMI and enter the following parameters:

```
configuration vcc new <iport> <ivpi> <ivci> <oport> <ovpi> <ovci>
[-upc <index>] [-name <name>]
```

The advanced options for call records are as follows:

These parameters are defined as follows:

Parameter	Description				
iport	The incoming port number.				
ivpi	The incoming virtual path number.				
ivci	The incoming virtual channel number.				
oport	The outgoing port number.				
ovpi	The outgoing virtual path number.				
ovci	The outgoing virtual channel number.				
-upc <index></index>	The integer index that refers to a specific UPC traffic contract. If no index is specified, then no traffic policing will take place on this VCI. It is assigned a UPC index of 0, and all traffic on this VCI is treated as UBR traffic. This is the default.				
name	The name you want to assign to this channel to identify it uniquely. It is useful for billing purposes so you can identify which channels are being used by which customers. Can be up to 32 ASCII characters long.				
inctype	The channel connection type for the incoming channel. For billing purposes, it denotes which switch this channel is arriving. Orig (originating) means that the ingress endpo of the channel is connected to the source node which is outside the network, tran (tr sit) means that the ingress endpoint of the channel is connected to a node within the r work, and term (terminating) means that the ingress endpoint of the channel is connected to the destination node which is outside the network.				
outctype	The channel connection type for the outgoing channel. For billing purposes, it denotes on which switch this channel is leaving. Orig (originating) means that the egress endpoint of the channel is connected to the source node which is outside the network, tran (transit) means that the egress endpoint of the channel is connected to a node within the network, and term (terminating) means that the egress endpoint of the channel is connected to the destination node which is outside the network.				
pmp <sup>1</sup>	Indicates this is a point-to-multipoint channel.				
mpp	Indicates this is a multipoint-to-point channel.				

Parameter	Description	
mpmp Indicates this is a multipoint-to-multipoint channel.		

<sup>1.</sup> By indicating pmp, mpp, or mpmp, you are only assigning a label for record keeping purposes. The switch does not necessarily create the type of channel you have specified. If you assign a connection type, but do not assign a pmp, mpp, or mpmp label, the switch assigns a label of pp (point-to-point).

The following is an example of how to create a uni-directional virtual channel which specifies the call record connection type as transit:

```
localhost::config vcc> new 3b1 0 100 3b4 0 100 -inctype tran -outctype tran
```

The following is an example of how to create a uni-directional virtual channel which has the name "customer\_a" assigned to it:

```
myswitch::configuration vcc> new 3b2 0 145 3b3 0 145 -name customer_a
```

The following is an example of how to create a bi-directional virtual channel on an ASX-1000. To create a VCC going in port 2A1, VPI 0, VCI 100 on the switch board installed in slot 2 and going out port 4B1, VPI 0, VCI 100 on the switch board installed in slot 4, enter the following:

```
myswitch::configuration vcc> new 2a1 0 100 2e4 0 100 myswitch::configuration vcc> new 2e4 0 100 2a1 0 100 myswitch::configuration vcc> new 4b1 0 100 4e2 0 100 myswitch::configuration vcc> new 4e2 0 100 4b1 0 100
```

In the first line in the first pair, notice that the output port is 2E4. This is the intra-fabric port. The 2 means the connection is coming out of the switch board in slot 2 through the intra-fabric port. The E represents the intra-fabric port. The 4 means the connection is destined for switch board in slot 4. 2E4 then becomes the input port in the second line.

In the first line in the second pair, notice that the output port is 4E2. This is the intra-fabric port. The 4 means the connection is coming out of the switch board in slot 4 through the intra-fabric port. The E represents the intra-fabric port. The 2 means the connection is destined for switch board in slot 2. 4E2 then becomes the input port in the second line.

Once the parameters are entered, the virtual channel is created instantly by the SCP. At the same time, a command is entered automatically into the current configuration database, which means that this virtual channel is created every time the SCP is restarted.

## 1.5.4 Creating a SPANS SPVC

To create a SPANS SPVC, you must configure both ends of the connection concurrently on the two switch fabrics. This means you must have an AMI session open on both the local switch fabric and the destination switch fabric. To create a new SPANS SPVC, log in to AMI and enter the following parameters:

```
myswitch::configuration spvc spans> new <port> <vpi> <vci> <dest-session> <dest-port> <dest-vpi> <dest-vci>\
[-peak <Kb/sec>] [(source | destination | bidirectional)]
```

#### These parameters are defined as follows:

Parameter	Description			
port	The port number on the local switch fabric.			
vpi	The virtual path number on the local switch fabric.			
vci	The virtual channel number on the local switch fabric.			
dest-session	The IP address of the remote switch.			
dest-port	The port number on the remote switch fabric.			
dest-vpi	The virtual path number on the remote switch fabric.			
dest-vci	The virtual channel number on the remote switch fabric.			
-peak <kb sec=""></kb>	The amount of peak bandwidth allocated for this SPANS SPVC, specified in kilobits per second. The default is 0.			
source   destination   bidirectional	source means a unidirectional SPANS SPVC going from the local switch fabric to the remote switch fabric will be created. destination means a unidirectional SPANS SPVC going from the remote switch fabric to the local switch fabric will be created. bidirectional means the pair of unidirectional SPANS SPVCs will be created. The default direction, if you do not specify one, is bidirectional.			



To create a bidirectional SPANS SPVC, you must either specify bidirectional, or you must set up two unidirectional SPANS SPVCs with one going in each direction.

To create a SPANS SPVC, you need to configure the two ends concurrently on the two switch fabrics. Therefore, you first need to open an AMI session to the destination switch fabric by using the SCP's IP address, along with the SNMP read-write community string. The following example depicts how to create a bidirectional SPVC from the local switch fabric (myswitch) to a remote switch fabric (198.29.22.46 named fishtank). The asterisk (\*) in front of the prompt indicates that it is a remote session. To return to the local session, you must type localhost (instead of the prompt name).

```
myswitch::> open 198.29.22.46 private

Opening a session for "198.29.22.46", please wait...
Connected to "198.29.22.46" (asx200bx).
*fishtank::> localhost

myswitch::> configuration spvc spans new ?

usage: new <port> <vpi> <vci> <dest-session> <dest-port> <dest-vpi> <dest-vci> \[-peak <Kb/sec>] [(source | destination | bidirectional)]

myswitch::configuration spvc spans> new lcl 0 49 198.29.22.46 lbl 0 50
```

# 1.5.5 Displaying SPANS SPVC Information

This command allows you to display all of the SPANS SPVCs on an individual switch fabric. Enter the following parameters:

The fields in this display are defined as follows:

Field	Description			
Local ID	The unique number that the local switch fabric's SCP assigned to this SPANS SPVC when it was created.			
Local Port	The port number on the local switch fabric.			
Local VPI	The virtual path number on the local switch fabric.			
Local VCI	The virtual channel number on the local switch fabric.			

# Configuring PVCs

Field	Description
Local BW	The amount of peak bandwidth allocated for this SPANS SPVC, specified in Kbps.
Remote ID	The unique number that the remote switch fabric's SCP assigned to this SPANS SPVC when it was created.
Remote Port	The port number on the remote switch fabric.
Remote VPI	The virtual path number on the remote switch fabric.
Remote VCI	The virtual channel number on the remote switch fabric.
Switch	The IP address or name of the remote switch fabric's SCP.

The following is displayed if no SPANS SPVCS have been configured:

myswitch::configuration spvc spans> show
No SPVC information is available

# 1.6 Traffic Types



The following sections contain informal definitions of the concepts presented. For a detailed understanding of these issues, please see the ATM Forum's UNI 3.1 and TM 4.0 Specifications.

Quality of Service (QOS) Management is based on the bandwidth parameters associated with a virtual connection and the class of service and ATM Adaptation Layer (AAL) used for that connection. In order to support voice, video, and data, the ATM Forum has defined four classes of service, or traffic types: Constant Bit Rate (CBR), Variable Bit Rate (VBR), Available Bit Rate (ABR), and Unspecified Bit Rate (UBR).

- At connection set-up time, traffic that uses a CBR parameter, such as a voice signal, makes a request for a dedicated Peak Cell Rate (PCR). Once the PCR is defined, the ATM network must be able to guarantee that amount of bandwidth for the duration of the connection.
- At connection set-up time, traffic that uses a VBR parameter, such as a video and data, makes a request for a dedicated PCR, Sustainable Cell Rate (SCR), and Maximum Burst Size (MBS). Once these cell rates are defined, the ATM network must be able to guarantee these rates for the duration of the connection.
- At connection set-up time, ABR traffic makes a request for a dedicated PCR and Minimum Cell Rate (MCR). Once these cell rates are defined, the ATM network must be able to guarantee the MCR rate for the duration of the connection. ABR traffic sources adjust their transmission rate in response to information they receive describing the status of the network and its capability to successfully deliver data.
- UBR traffic, such as broadcast information and ARP messages, is also known as "best effort" service. UBR provides no bandwidth guarantees.

Because ATM is designed to provide a single network to transport this variety of traffic classes, FORE's traffic policing and Connection Admission Control (CAC) schemes are vital to allowing this mix of traffic to flow smoothly.

# 1.7 Traffic Policing (Usage Parameter Control)

Traffic policing, also known as Usage Parameter Control (UPC), is a method of ensuring fair allocation of network resources and of assessing the cells entering the switch for conformance with pre-established traffic bandwidth contracts. Those cells that exceed the specified contract are "tagged" or "dropped," depending on what is defined in the contract. This ensures that the connections with reserved bandwidth are not exceeding their reservations. FORE Systems' switches use a combination of "leaky bucket," or Generic Cell Rate Algorithm (GCRA) hardware in the switch fabric and user-configurable parameters in AMI to perform these policing functions.

# 1.7.1 Leaky Bucket Algorithm

The first important concept to understand is the leaky bucket algorithm. Leaky buckets are a mechanism by which cells entering the switch fabric are monitored for compliance with UPC traffic contracts that have been negotiated at connection set-up time. Before the leaky buckets are discussed, it is important to understand the parameters that are being measured by the buckets, as shown in Table 1.1. These parameters are informally defined as follows:

- Peak Cell Rate (PCR) the maximum number of cells per second
- Cell Delay Variation Tolerance (CDVT) the tolerance for variation in the interarrival time of these cells, or the amount of jitter that can be accepted by the network
- Sustainable Cell Rate (SCR) the average rate of cell transmission for this connection, taking bursting into account
- Burst Tolerance (BT) the maximum amount of cells that can be transmitted
- Minimum Cell Rate (MCR) the minimum rate that the network has to guarantee for an ABR connection

The leaky bucket algorithm is basically a timer which assesses if cells entering the switch fabric conform to the parameters listed above. As a cell arrives, the timer assesses if the cell is on time, late, or early. If the cell is determined to be on time or late (based on the traffic parameters), the cell is allowed to pass unchanged. If the cell is early (which, in turn, causes the cell stream to exceed the specified parameters), the cell is considered non-conforming and is either dropped or tagged (the CLP bit is set to 1), depending on the specified contract.

The first bucket in this analogy measures the PCR, or the rate at which the bucket drains. It also considers the CDVT, or the depth of the bucket. The second bucket measures the SCR, or the rate at which the bucket drains, and the BT, or the depth of the second bucket.

# 1.7.2 Non-conforming Cells: Tagging vs. Dropping

Second, it is important to understand the concept of tagging and dropping. Each ATM cell has a Cell Loss Priority (CLP) bit which indicates if the network can drop it under congested conditions. When the CLP bit is set to 0 (or CLP=0), the cell is assessed for compliance with traffic parameters associated with the CLP=0 stream. If the traffic parameters dictate that non-compliant cells should be "tagged," the CLP bit is set to 1 (or CLP=1) by the UPC contract, which means that upon experiencing congestion further in the network, these CLP=1 cells are dropped in preference to CLP=0 cells.

Table 1.1 shows the various traffic contracts and user-configurable actions to be executed upon the detection of non-conforming cells.

Traffic Type	PCR	SCR	MBS	Bucket 1	Bucket 2	Policing on Switch Fabric	Action on non-conforming cells	
							Bucket 1	Bucket 2
CBR	PCR01			PCR01, CDVT		Yes	Drop CLP=0+1	
CBR0	PCR0, PCR01			PCR01, CDVT	PCR0, CDVT	Yes	Drop CLP=0+1	If enabled, Tag CLP=0; else, Drop CLP=0
VBR	PCR01	SCR01	MBS01	PCR01, CDVT	SCR01, BT01+CDVT	Yes	Drop CLP=0+1	Drop CLP=0+1
VBR0	PCR01	SCR0	MBS0	PCR01, CDVT	SCR0, BT0+CDVT	Yes	Drop CLP=0+1	If enabled, Tag CLP=0; else, Drop CLP=0

**Table 1.1 - Summary of Traffic Contract Variables and Policing Actions** 

### 1.7.3 UPC Traffic Contract Parameters

The ATM Forum has defined different types of traffic contracts to be used in conjunction with these leaky buckets. The parameters that make up these types of contracts are defined as follows:

- pcr0 PCR for cells with CLP=0
- pcr01 PCR for the aggregate of the CLP=0 cells and the CLP=1 cells (all cells)
- scr0 SCR for cells with CLP=0
- scr01 SCR for the aggregate of the CLP=0 cells and the CLP=1 cells (all cells)
- mbs0 MBS for cells with CLP=0
- mbs01 MBS for the aggregate of the CLP=0 cells and the CLP=1 cells (all cells)
- tag sets CLP bit=1 for CLP=0 cells that fail the PCR0 test for CBR0 contracts or the SCR0/MBS0 test for VBR0 contracts

The specific combinations of these parameters that make up the ATM Forum contracts are defined as follows:

- 1. cbr <pcr01>
- 2. cbr0 < pcr0 > < pcr01 > [tag]
- 3. vbr < pcr01 > < scr01 > < mbs01 >
- 4. vbr0 < pcr01 > < scr0 > < mbs0 > [tag]
- 5. abr <pcr01> <mcr>

The cbr <pcr01> contract is for CBR traffic. It only uses the first leaky bucket to assess the conformance to PCR of the aggregate of the CLP=0 cells and the CLP=1 cells. Cells which fail the PCR CLP=0+1 test are discarded.

The cbr0 <pcr0> <pcr0> <pcr01> [tag] contract is for CBR traffic. It uses the first leaky bucket to assess the conformance to PCR of the CLP=0 cells. It uses the second leaky bucket to assess the conformance to PCR of the aggregate of the CLP=0 and the CLP=1 cells. If the tag option is set, the cells which fail the PCR CLP=0 test are tagged as CLP=1 and passed on to the second leaky bucket to be tested for PCR CLP=0+1 conformance. Cells which fail the PCR test on CLP=0+1 are discarded. If the tag option is not set, cells which fail the PCR CLP=0 test are discarded and cells which fail the PCR test on CLP=0+1 are discarded.

The vbr <pcr01> <scr01> <mbs01> contract is for VBR traffic. The first leaky bucket assesses the conformance to PCR of the aggregate of CLP=0 cells and the CLP=1 cells and the second leaky bucket assesses the conformance to SCR and BT of this same combination. Cells which fail the PCR test are dropped. Cells which pass the PCR test, but which fail SCR and BT test are dropped.

The vbr0 <pcr01> <scr0> <mbs0> [tag] contract is for VBR traffic. It uses the first leaky bucket to assess the conformance to PCR of the aggregate of the CLP=0 and the CLP=1 cells. Cells that fail this test are discarded. It uses the second leaky bucket to assess the conformance to SCR and BT of the CLP=0 cells. If the tag option is set, the cells which fail the SCR and BT CLP=0 test are tagged as non-conforming cells. If the tag option is not set, cells which fail the SCR and BT CLP=0 test are discarded.

The abr cr01> <mcr> contract is for ABR traffic. It only uses the first leaky bucket to assess the conformance to PCR of the aggregate of the CLP=0 cells and the CLP=1 cells. Cells which fail the PCR CLP=0+1 test are discarded. The MCR is the guaranteed minimum rate and the PCR is the maximum required rate. The MCR may be set to 0, which indicates "best effort" service. If MCR is set greater than 0, then the rate is guaranteed, up to MCR. However, between MCR and PCR, cells may be dropped when congestion is experienced.

### 1.7.4 AMI UPC Commands

AMI allows you to create a UPC contract using these combinations of traffic parameters. To create a UPC contract in AMI, enter the following parameters:

Where UPC is one of the following combinations of traffic parameters:

```
cbr <pcr01>
cbr0 <pcr0> <pcr01> [tag]
vbr <pcr01> <scr01> <mbs01>
vbr0 <pcr01> <scr0> <mbs0> [tag]
```

These parameters are defined as follows:

Parameter	Description
index	The integer index that refers to this specific traffic contract. Valid index numbers are from 0 to 32,767.
UPC	One of the types of traffic contracts shown above. The parameters in these contracts are defined as follows:
ubr	Indicates UBR traffic.
$\mathrm{cbr}^1$	Indicates CBR traffic.
cbr0	Indicates CBR0 traffic.
vbr	Indicates VBR traffic.
vbr0	Indicates VBR0 traffic.
pcr0	Indicates the peak cell rate for cells with CLP = 0.
pcr01	Indicates the peak cell rate for all cells.
scr0	Indicates the sustainable cell rate for cells with CLP = 0.
scr01	Indicates the sustainable cell rate for all cells.

Parameter	Description
mbs0	Indicates the maximum burst size for cells with $CLP = 0$ .
mbs01	Indicates the maximum burst size for all cells.
tag	tag means that non-conforming CLP = 0 cells are tagged. Otherwise, they are dropped. The default is that they are dropped. This option only applies to the PCR0 parameter of the CBR0 contract and to the SCR0 and MBS0 parameters of the VBR0 contract.
abr	Indicates ABR traffic. Currently, ABR UPC contracts are supported only on Series D network modules.
mcr	Indicates the minimum cell rate for all cells. ABR connections with an MCR equal to 0 use the roundrobin scheduling discipline. ABR connections with an MCR greater than 0 use the guaranteed scheduling discipline. However, if there are no suitable rate groups in the rate controller, the ABR connections with an MCR greater than 0 are rejected. Currently, ABR UPC contracts are supported only on Series D network modules.
-cdvt us	The Cell Delay Variation Tolerance (CDVT) associated with the peak cell rates, in microseconds. If the CDVT is not specified here, the default CDVT value associated with the port will be used. (See conf port show and conf port cdvt for more information).
noGCRA	nogcra means that GCRA policing is disabled on CBR or VBR (depending on what is configured) connections using this contract. If nogcra is not entered, then GCRA policing is enabled on CBR or VBR (depending on what is configured) connections using this contract. By default, nogcra is not entered (GCRA policing is enabled). You must use the nogcra option when applying a UPC contract to the outbound signalling channel using the -outsigupc <upc-index> variable under conf signalling new to prevent the outbound signalling channel from being policed.</upc-index>
aal5	The connection is using the AAL5 Adaptation Layer.
noPktDisc	This optional parameter can only be used if the connection is AAL5 (i.e., the aal5 parameter is present). This parameter suppresses EPD/PPD (AAL5 packet discard) on the connection. The default is for this parameter not to be present (EPD/PPD is enabled).
ubrTagging	ubrTagging means that all UBR traffic is tagged (set to CLP=1) on this connection. If ubrTagging is not entered, then UBR traffic is not tagged on this connection. This command only applies to UBR traffic. By default, UBR traffic is not tagged.
PPPol <sup>2</sup>	This optional parameter can only be used if the connection is AAL5 (i.e., the aal5 parameter is present). This parameter indicates that Partial Packet Policing is going to be performed on this connection. The default is for this parameter not to be present, which leaves Partial Packet Policing disabled.
AltCLP	This optional parameter only applies to connections on Series D network modules. It indicates that the alternate CLP threshold (configured using conf module traffic d altclpthresh) should be used for all connections created with this UPC contract. The default is for this parameter not to be present, which means the connections will not use the alternate CLP threshold.

Parameter	Description
-scheduling (roundrobin   smoothed   guaranteed) <sup>3</sup>	Indicates the scheduling mode to be used for servicing traffic on the output side of a Series D network module. roundrobin means that all service for these connections comes from one of the round-robin queues in the network module. This is the default mode for both SVCs and PVCs. smoothed means that all service for these connections comes from the network module's rate controller, which ensures that cells for these connections are transmitted into the network at a fixed rate of R cells per second. guaranteed is a combination of the round-robin and smoothed modes. Service for these connections are scheduled with both fixed rate R from the rate controller, and they have an entry in the appropriate round-robin queue.
-name <name></name>	The user-defined name associated with this UPC traffic contract. This helps you remember for what traffic type this specific contract is used. If you do not specify a name, a default name that relates to this type of traffic contract is assigned automatically.
-bc <bits></bits>	The committed burst size of a connection, in bits. Can only be used on a Frame Relay connection.
-be <bits></bits>	The excess burst size of a connection, in bits. Can only be used on a Frame Relay connection.
-cir <kbps></kbps>	The committed information rate of a connection, in kbps. Can only be used on a Frame Relay connection.
-ar <kbps></kbps>	The access rate of a Frame Relay UNI, in kbps. Can only be used on a Frame Relay connection.
-frsize <bytes></bytes>	The average frame size, in bytes. Can only be used on a Frame Relay connection.

<sup>1.</sup> The units for pcr0, pcr01, scr0, scr01, mbs0, and mbs01 are specified either in cells per second or in kilobits per second, depending on what you used for conf system units. To display the current setting, use conf system show. The default is cps (cells per second).

<sup>3.</sup> The -scheduling option has an effect only on connections with outputs on Series D network modules. All other network module platforms only use roundrobin scheduling.



Remember, when you create this UPC contract, it is not actually used until you assign it to a VPC, VCC, or SPANS path. UPC contracts are not assigned to VPTs.



When the advanced options are used, they are converted to ATM UPC parameters and are displayed as ATM parameters.

<sup>2.</sup> The HDCOMP ASIC must be version 1 or greater to support AAL5 partial packet policing. To display the ASIC version, use the conf board show advanced command.

The following is an example of how to create a UPC contract:

myswitch::configuration upc> new 5 vbr0 500 200 250 -cdvt 1000 aal5 PPPol -name vbr0\_upc

This example specifies a contract named "vbr0\_upc", which is a VBR0 contract with an index of 5, a pcr01 of 500 cells/sec (or kbps), an scr0 of 200 cells/sec (or kbps), an mbs0 of 250 cells (or kilobits), a CDVT of 1,000 microseconds, and partial packet policing enabled.



For more information regarding traffic contracts, please refer to Table 5-7 in the ATM Forum UNI 3.0 Specification.



PVCs that use UPC contracts containing any of the [nogCRA], [aal5 [noPktDisc] [PPPol]], and [ubrTagging] options are valid only when the conf port gcrapolicing, conf port aal5packetdiscard, conf port pppolicing, and conf port ubrtagging commands are set to svcOn or svcOff. Use conf port show tm to check these settings.

# CHAPTER 2

# **Configuring Classical IP**

#### 2.1 Introduction

This chapter describes how to design, configure, and maintain a Classical IP ATM network. The term classical indicates that the ATM network has the same properties as existing legacy LANs. That is, even though ATM technology allows for large, globally connected networks, for example, it is only used in the LAN environment as a direct replacement of existing LAN technology. The classical model of LANs connected through IP routers is maintained in ATM networks. RFC-1577 provides the standard for Classical IP over ATM.

Classical IP over ATM is different than IP in legacy LANs in that ATM provides a virtual connection environment through the use of Permanent Virtual Circuits (PVCs) and/or Switched Virtual Circuits (SVCs). SVC management is performed via the ATM Forum UNI 3.0 Specification, which specifies Q.2931. Q.2931 is a broadband signalling protocol designed to establish connections dynamically at the User-Network Interface (UNI). Q.2931 uses Service Specific Connection Oriented Protocol (SSCOP) as a reliable transport protocol, and all signalling occurs over VPI: 0, VCI: 5. Q.2931 connections are bidirectional, with the same VPI/VCI pair used to transmit and receive.

Once a Classical IP connection has been established, IP datagrams are encapsulated using IEEE 802.2 LLC/SNAP and are segmented into ATM cells using ATM Adaptation Layer type 5 (AAL5). Additionally, the default Maximum Transmission Unit (MTU) is 9,180 bytes (the SNAP header adds 8 more bytes) with a maximum packet size of 65,535 bytes. There is currently no support for IP broadcast datagrams or IP multicast datagrams in a Classical IP environment.

## 2.1.1 Logical IP Subnets

An important concept in Classical IP networks is that of a Logical IP Subnet (LIS). An LIS is a group of hosts configured as members of the same IP subnet (that is, they have the same IP network and subnetwork numbers). In this sense, one LIS can be equated to one legacy LAN. It is possible to maintain several overlaid LISs on the same physical ATM network. Therefore, in a Classical IP ATM network, placing a host on a specific subnet is a logical choice rather than a physical one. In this type of environment, communication between hosts in different LISs is only permitted by communicating through an IP router which is a member of both LISs (as per RFC-1577).

The number of LISs, and the division of hosts into each LIS, is purely an administrative issue. Limitations of IP addressing, IP packet filtering, and administrative boundaries may guide a manager into establishing several LISs onto a single ATM network. Keep in mind, though, that communication between LISs must occur through IP routers.

#### 2.1.2 Classical IP Interfaces

In order to support routing between multiple LISs, the switch software allows a switch to be configured as a member of (and a router between) up to four distinct LISs. (The host adapter software allows a host to be configured as a member of (and a router between) up to 16 distinct LISs.) Each LIS membership is through a separate Classical IP network interface. Existing system level IP routing configuration tools are used to control routing through each of the Classical IP interfaces in the same manner as routing among several physical interfaces. Even though each Classical IP interface associated with a given physical interface uses the same physical hardware, they are each configured separately with their own MTU, IP address, and ATM address.

By default, the name of each of the Classical IP interfaces on a switch begins with qa. (On a host, each of the Classical IP interfaces begins with ci and is user-configurable.) On a switch, all of the Classical IP interfaces associated with physical unit zero have a as the next letter. All of the Classical IP interfaces associated with physical unit one have b as the next letter, and so forth. Finally, each Classical IP interface has its interface number as a suffix. As an example of the above naming convention for switches, the name of the third Classical IP interface (unit 2) on physical unit one is qab2.

#### 2.1.3 SPANS Interface

While each of the Classical IP interfaces for a given physical interface is designed to support Classical IP using Q.2931 signalling, a SPANS interface also exists for each physical interface. The SPANS interface is asx0 in switch software and is fa0 in host adapter software.

The SPANS interface supports FORE IP on top of SPANS signalling. FORE IP allows communication using AAL4 or AAL5 with no encapsulation, uses a broadcast ARP for SPANS address resolution, and supports direct communication of all hosts on a physical ATM network without the use of IP routers. Since SPANS and Q.2931 signalling use different VCIs, a host can simultaneously support FORE IP over SPANS as well as Classical IP over Q.2931 on the same physical interface.

As a result of standard IP routing, all traffic sent out a SPANS interface uses FORE IP, while all traffic sent out a Classical IP interface uses Classical IP. The Classical IP interfaces are qaaX (where X is 0 - 3) in switch software and are ci in host adapter software.

Each of the SPANS interfaces should be assigned an IP address on a subnet different than the subnets of any of the Classical IP interfaces. It is permissible to place multiple SPANS interfaces on the same subnet, and the driver load balances connections across these interfaces.

It is only necessary to configure the SPANS and Classical IP interfaces if the specific service provided by that interface is required. A host sending only Classical IP would not need to configure the SPANS interfaces. Likewise, a host sending only FORE IP would not need to configure the Classical IP interfaces. Both the SPANS and Classical IP interfaces may be configured simultaneously, but they must be in separate subnets. Remember that Classical IP specific configuration changes can only be done with the Classical IP devices, while SPANS specific configuration changes can only be done with the SPANS devices.

# 2.2 Address Registration and ILMI

Before a host can establish connections over a physical interface, the host must know the NSAP address for that interface. The primary purpose of Interim Local Management Interface (ILMI) is to discover and register these NSAP addresses dynamically.

#### 2.2.1 NSAP Addresses

For private ATM networks, addresses uniquely identify ATM endpoints. The UNI 3.0 address format is modeled after that of an OSI Network Service Access Point, hence the name NSAP address.

Three address formats have been specified: DCC, ICD, and E.164. FORE implements the ICD ATM format. Per the UNI 3.0 specification, all private networks should accept initial call set-up messages containing ATM addresses with any of the approved formats and forward the calls as necessary.

An NSAP address consists of the following:

- a 13-byte network-side prefix The prefix is the NSAP prefix of the switch to which the host is attached.
- a seven-byte user-side part This consists of the following:
  - a six-byte End System Identifier (ESI) The ESI is the unique IEEE MAC address of the interface.
  - a one-byte selector Although each Classical IP interface for a given physical interface uses the same prefix and ESI, the selector field is the part that indicates the number of the specific Classical IP interface. On a switch, the selector field is 00 for qaa0, 01 for qaa1, 02 for qaa2, and 03 for qaa3.

## 2.2.2 Operating with ILMI Support

FORE Systems switches running *ForeThought* software provide support for ILMI. If ILMI is supported on all of the switches and hosts in a given network, when a switch boots up, ILMI enables the switch to discover all of the hosts attached to it and to send its ATM prefix associated with the port to those hosts dynamically. In return, the host prepends that prefix to its ESI and selector fields, forming a complete ATM address. The host then notifies the switch of its complete ATM address. These registration messages are sent and received over AAL5 using VPI: 0, VCI: 16. Once ILMI registration has been completed, then connection setup can occur.

If a host changes network ports after an ATM address has been registered for its interface, all existing connections are closed. If the new port is on a different switch, a new ATM address (with a different network address prefix) is registered. The host can then begin to establish new connections.

## 2.2.3 Operating without ILMI Support

If ILMI is not supported on a particular switch or host in a given network, then the ATM addresses must be manually configured. If another vendor's switch does not support ILMI, it can not supply an ATM prefix to the hosts. Therefore, the user must assign a unique, valid prefix to the switch. Additionally, the same prefix should be used for all hosts attached to that switch.

On the host, uniconfig is used to configure the ATM address for a specific interface. The switch directly attached to this interface is then informed of this ATM address/port combination through commands in the ATM Management Interface (AMI). Once the host and network have both been informed of this ATM address/port pair, the host may begin signalling.

## 2.2.4 Configuration

The choice to use ILMI for address registration is made at software installation time. Since ILMI uses SNMP as its management protocol, the use of ILMI is tied into snmpd. The choice can be made to run FORE's SNMP agent and use ILMI (snmpd), run FORE's SNMP agent without using ILMI (snmpd -i), or just use ILMI (snmpd -i or ilmid -i).

#### 2.3 ARP and ARP Servers

## 2.3.1 Theory

In order for a host to establish a connection to another host, it must first determine the other host's ATM address. ATM ARP (ATM address resolution protocol) is the procedure used to resolve an IP address into an ATM address. Since the ATM standards do not currently support broadcast on an ATM LAN, address resolution is performed by direct communication with a special ARP server, rather than broadcasting ARP requests as is done in legacy LANs. Each LIS must have only one ARP server configured, but a single ARP server can be the server for several LISs.

Each host in an LIS must be configured with the ATM address of the host providing ARP service for its LIS. On a switch, the ATM address of the ARP server can be obtained by using the AMI command conf atmarp arpserver show. On a host, the ATM address of the ARP server can be obtained by using clipconfig show (remember to use the interface associated with the given LIS).

On a switch, the ATM address of the ARP server can be configured by using the instructions found in Section 2.3.2 of this manual. On a host, the ARP server address is configured into the host at installation time using the configure\_atm script.

Since only one ARP server can be functioning at a time in a given LIS, and since the ARP server's address is configured into each host, it is not possible to use multiple, redundant ARP servers to improve robustness. If an ARP server becomes nonfunctional, a new ARP server must be configured, and then each host within the LIS must be configured to use the new ARP server. To configure a new ARP server address on a switch, use the instructions found in Section 2.3.2 of this manual. To configure a new ARP server address on a host, you must delete the interface and recreate it.

## 2.3.2 Configuring a FORE Switch to be an ARP Server

FORE's ATM switches also have the capability of being an ARP server. (FORE's ATM adapters can not be configured as an ARP server.) To configure a *ForeRunner* ATM switch as an ARP server, perform the following steps on <u>only one</u> of the SCPs:

1. On one of the SCPs, determine the ATM address of that SCP for the relevant interface (qaa0 -> qaa3) using the following AMI command:

configuration atmarp getnsap <interface>

For example:

configuration atmarp getnsap qaa0

gaa0 NSAP address: 47000580ffe1000000f12400de0020481900de00

2. Set the NSAP address of the ARP server to be the ATM address of the interface that you displayed in step 1 using the following AMI command:

conf atmarp arpserver set <NSAPaddress> [<interface>]

For example:

conf atmarp arpserver set 0x47000580ffe1000000f12400de0020481900de00 qaa0

3. Set the ATM address of the ARP server on each of the other switches that will use that switch as the ARP server using the same command found in step 2.

## 2.3.3 Classical IP Operation

Once a host knows its own ATM address and the ATM address of its ARP server it attempts to establish a connection to the ARP server, which is used to send ARP requests and receive ARP replies. When the connection to the ARP server has been established, the ARP server sends an inverse ARP (InARP) request on the new VC to learn the host's IP address. When an InARP reply is received, the ARP server places that host's IP address to ATM address mapping in its ARP cache. Therefore, over time, the ARP server dynamically learns the IP-to-ATM address mappings of all the hosts in its LIS. It can then respond to ARP requests directed toward it for hosts in its LIS.



In order for a host to communicate with an ARP server, it must have learned its own ATM address and have been configured with the ATM address of the ARP server.

A host can not resolve the ATM addresses of other hosts in its LIS unless it can communicate with its ARP server.

Since there is no mechanism for ARP servers to exchange mapping information with each other, it is imperative that each LIS be configured with only one ARP server.

When a host wants to communicate with another host in its LIS, it first sends an ARP request to the ARP server containing the IP address to be resolved. When an ARP reply is received from the ARP server, the host creates an entry in its ARP cache for the given IP address and stores the IP-to-ATM address mapping. This ARP cache entry is marked as complete. To ensure that all of the IP-to-ATM address mappings known by a certain host are up-to-date, hosts are required to age their ARP entries. A host must validate its ARP entries every 15 minutes (20 minutes on an ARP server). Any ARP entries not associated with open connections are immediately removed.

A host validates its SVCs by sending an ARP request to the ARP server. A host validates its PVCs, and an ARP server validates its SVCs, by sending an InARP request on the VC. If a reply is not received, the ARP entry is marked invalid. Once an ARP entry is marked invalid, an attempt is made to revalidate it before transmitting. Transmission proceeds only when validation is successful. If a VC associated with an invalid ARP entry is closed, the entry is removed.

#### 2.3.4 Operational Issues

Certain hosts in an LIS may not support Classical IP. It is still possible to communicate with these hosts (and for these hosts to communicate with one another) by using static ARP entries. If a host does not support Classical IP, its IP-to-ATM address mapping should be placed in its ARP server's cache as a static entry. This allows other hosts that do support Classical IP to contact their ARP server as usual and obtain the correct address mapping. If a host that does not support Classical IP wants to initiate connections, the IP-to-ATM address mappings of the destination hosts should be put in its ARP cache, again as static entries. By using static ARP entries in the above fashion, the ability for all hosts to communicate is maintained.

There are some restrictions on the number of hosts that can be maintained dynamically. They are as follows:

- In the default configuration, a host can only have approximately 250 virtual connections open simultaneously. This means that an ARP server can only serve 250 clients, since each client must maintain a connection with its ARP server. This may be a limitation if the ARP server is servicing multiple LISs.
- It is possible to increase the number of connections using AMI.
- Some hosts may be limited to supporting a maximum of 1,024 connections per adapter.

## 2.4 Classical IP PVCs

## 2.4.1 Theory and Configuration

Normally, ATM connections in a Classical IP environment are established dynamically using UNI 3.0 or UNI 3.1. ARP, ILMI, and UNI 3.0 or UNI 3.1 all work together as described previously to set up an SVC. If a host from another vendor does not support Classical ARP or ILMI, it is still possible to set up an SVC using work-arounds. If a host or a switch in an LIS does not support UNI 3.0 or UNI 3.1, however, it is not possible to establish an SVC. In this case, a Classical IP PVC can be used for communication.

On each of the hosts, cliparp add -pvc is used to establish the PVC. An unused VPI/VCI pair must be chosen for each host. PVCs using the chosen VPI/VCI pairs must also be set up from each of the hosts to their connecting switch, and then on all of the switches between the two connecting switches.



Both the incoming and outgoing connections are set up simultaneously on the host, but they must be set up individually on the switches. The same VPI/VCI pair is used by a host to send on the PVC as well as receive on the PVC. The IP datagrams are sent over the PVC using AAL5 with LLC/SNAP encapsulation.

#### 2.4.2 Revalidation and Removal

Normally, the device driver periodically checks that its PVCs are still established and functioning. A host revalidates a PVC by sending InARP requests over the PVC, if the user specifies that revalidation should occur by using the <code>-reval</code> option to <code>cliparp</code> <code>add</code> <code>-pvc</code> at the time the PVC is created (e.g., <code>-reval</code> 10 means revalidation will occur every 10 minutes). If the equipment attached to the FORE equipment supports revalidation, the user must choose the <code>-reval</code> option. If an InARP reply is not received, the revalidation fails, the PVC is marked invalid (as shown through <code>cliparp</code> <code>show</code>), and communication over the PVC is no longer possible.

Once a PVC is marked invalid, an attempt is made to validate the PVC before transmitting. Transmission proceeds only when validation is successful. It is possible to disable this revalidation feature by not specifying the <code>-reval</code> option. This is often desirable when the remote end of the PVC (such as a video camera) does not support InARP.

A Classical IP PVC is removed on the host side using cliparp delete -pvc. Both the incoming and outgoing connections are removed simultaneously. The PVC must then be removed from each of the network switches involved.

# 2.5 Configuring the Network

In an ATM network, before any connections can be made, the two parties must know each other's ATM address in order to set up that connection.

To allow those connections to work, the ideal scenario is for all hosts and switches in the network to have support for both ILMI and for RFC-1577 (Classical IP over ATM). However, when using other vendors' equipment, this may not be the case. This section describes how to configure a network with the following scenarios:

- Configuring a third-party host that has no ILMI and no RFC-1577 support
- Configuring a third-party switch that has ILMI support, but no RFC-1577 support
- Configuring a third-party switch that has no ILMI support, but has RFC-1577 support

## 2.5.1 Third-Party Host with No ILMI and No RFC-1577 Support

To configure a network with a third-party vendor's host (or an edge device) that supports neither ILMI nor RFC-1577 (as shown in Figure 2.1), perform the following steps:

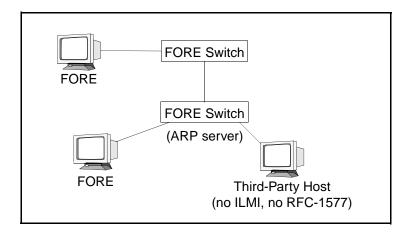


Figure 2.1 - Configuring a Third-Party Host with No ILMI and No RFC-1577 Support

- 1. Before beginning this process, be sure that *ForeThought* software is installed and running on the FORE equipment.
- 2. Using the configuration software of the third-party host, assign that host an NSAP address that has the same prefix as the switch fabric to which it is connected.
- 3. Configure the switch that is the ARP server so that it has a static route to the third-party host using the following AMI command:

conf atmr ftpnni staticroute new <NSAP> <mask> -port <port> -vpi <vpi>

Be sure to use a host mask value of 152.

## 2.5.2 Third-Party Switch with ILMI and No RFC-1577 Support

To configure a network with a third-party vendor's switch that supports ILMI, but not RFC-1577, (as shown in Figure 2.2), perform the following steps:

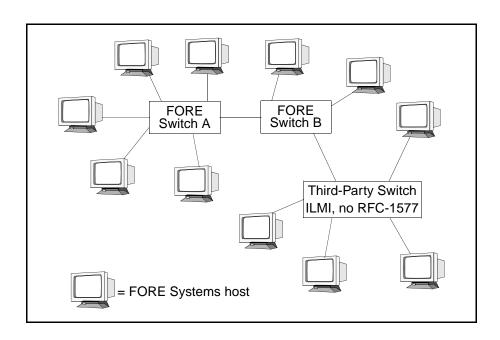


Figure 2.2 - Configuring a Third-Party Switch with ILMI Support and No RFC-1577

- 1. Be sure that *ForeThought* software has been installed on all of the hosts and that ILMI was set in the process. ILMI dynamically performs address registration for all of the hosts.
- 2. Configure a static ATM route to the third-party switch on FORE switch "B" that is physically connected to the third-party switch using the following AMI command:

conf atmr ftpnni staticroute new <NSAP> <mask> -port <port> -vpi <vpi>

Be sure to use a network mask value of 104.

3. Configure two static NSAP routes on the third-party switch, one to each of the FORE switches to which the third-party switch is connected, using the third-party vendor's configuration software. Be sure to use a network mask value of 104.

## 2.5.3 Third-Party Switch with RFC-1577 and No ILMI Support

To configure a network with a third-party vendor's switch that does not support ILMI, but does support RFC-1577 (as shown in Figure 2.3), perform the following steps:

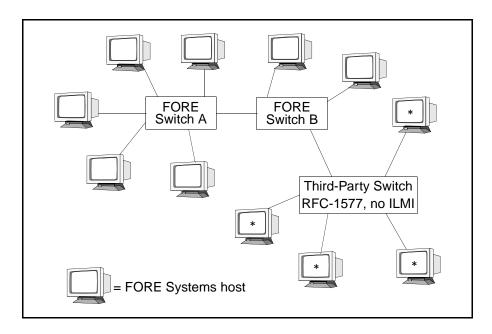


Figure 2.3 - Configuring a Third-Party Switch with RFC-1577 and No ILMI Support

- Be sure that *ForeThought* software has been installed on all of the FORE hosts and that ILMI was set in the process. ILMI dynamically performs address registration for all of the FORE hosts and FORE switches.
- 2. Statically configure the non-FORE (\*) hosts with ATM addresses (edit the firmware download script), using the same switch prefix for all of the hosts.
- 3. Configure a static ATM route to the third-party switch on FORE switch "B" that is physically connected to the third-party switch using the AMI command:

conf atmr ftpnni staticroute new <NSAP> <mask> -port <port> -vpi <vpi> Be sure to use a network mask value of 104. Also, be sure to use the same prefix that was used to configure the hosts.

4. Configure two static ATM routes on the third-party switch, one to each of the FORE switches to which the third-party switch is connected, using the third-party vendor's configuration software. Be sure to use a network mask value of 104.

# CHAPTER 3

# **Configuring an Emulated LAN**

#### 3.1 Introduction

This chapter describes how to design, configure, and maintain an Emulated LAN (ELAN) over an ATM network. An ELAN provides communication of user data frames among all members of the ELAN, similar to a physical LAN. One or more ELANs may run simultaneously (and independently) on the same ATM network.

Each ELAN is composed of a set of LAN Emulation Clients (LECs), a LAN Emulation Configuration Server (LECS), and at least one LAN Emulation Server (LES) and Broadcast and Unknown Server (BUS) pair (also referred to as colocated BUS or an intelligent BUS). In the current software release, the LECS may reside either in a *ForeRunner* ASX-200WG, ASX-200BX, ASX-1000, or LE 155 switch, or in a UNIX workstation running Solaris 2.5, 2.5.1, or 2.6. The LES/BUS pair may reside either in a *PowerHub* 7000, an ASX-200WG, ASX-200BX, ASX-1000, or LE 155 switch, or in a UNIX workstation running Solaris 2.5, 2.5.1, or 2.6. An additional software feature is Distributed LAN Emulation (DLE), which provides load sharing and fault tolerance to the ELAN.

The current software release supports the emulation of both Ethernet (IEEE 802.3) and Token Ring ELANs.

#### 3.1.1 Ethernet ELANs

The current software release supports emulation of Ethernet (IEEE 802.3) ELANs. In the current release, each LEC resides on an ATM host system (PC, Macintosh, UNIX workstation, ATM switch, *PowerHub* 7000, or ES-3810).

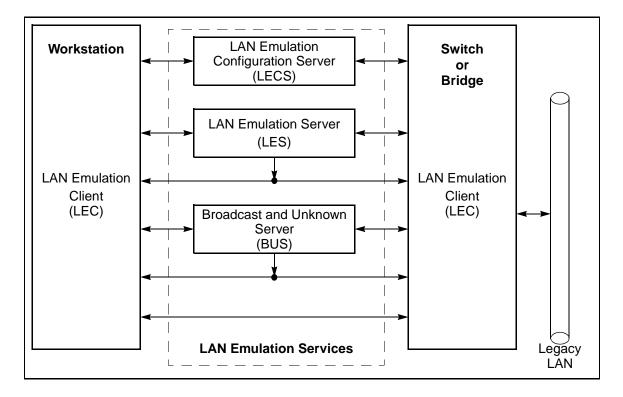
## 3.1.2 Token Ring ELANs

The current software release supports emulation of Token Ring (IEEE 802.5) ELAN services. In the current release, Token Ring LECs can not be created on the switch. Token Ring LECs can only be created on some hosts.

## 3.2 ELAN Components

The components of an ELAN include LECs, and LAN Emulation services consisting of a LECS, a LES, and a BUS. Although the ATM Forum specification allows the LES and BUS to be located on different devices, more intelligent traffic handling is possible when they are located on the same device. *ForeThought* 5.0.x or greater software requires the LES and BUS be co-located (reside on the same device).

The LECS may reside in the same physical system as the LES/BUS or in a separate physical system. For example, the LECS could reside in a switch, while the LES/BUS reside in a work-station. In *ForeThought* 5.2.x software, the LECS is supported only on *ForeRunner* switches and on systems running Solaris. The LES/BUS are supported only on *ForeRunner* switches, on *PowerHub* 7000s, and on systems running Solaris. The functional interconnections of a simple ELAN consisting of two LECs, an LECS, a LES, and a BUS are shown in Figure 3.1.



**Figure 3.1 -** Basic Emulated LAN Interconnections

## 3.2.1 LAN Emulation Client (LEC)

The LEC is the component in an end system that performs data forwarding, address resolution, and other control functions when communicating with other components within the ELAN. It also provides a MAC level emulated Ethernet or Token Ring interface and appears to higher level software as though a physical interface is present. Each LEC must register with both the LES and BUS associated with the ELAN it wishes to join before it may participate in the ELAN. To participate in multiple ELANs, an end system must have multiple LECs. *ForeThought* 5.2.x supports up to 16 LECs on ES-3810 switches and on adapter cards running Solaris or Windows/NT software, and up to 4 LECs on on adapter cards running Windows/95.

## 3.2.2 LAN Emulation Configuration Server (LECS)

The LECS is responsible for the initial configuration of LECs. It provides information about available ELANs that a LEC may join, together with the address of the LES associated with each ELAN. Using DLE in *ForeThought* 5.2, the user may also configure the LECS to associate multiple LES/BUS pairs with a given ELAN. This feature allows LECs to use a single, anycast address to reach one of the other DLE peer servers for their ELAN if their local server goes down. Normal address resolution through *ForeThought* PNNI, ATM Forum PNNI, or IISP will locate the closest, active LES which is using the anycast address.

## 3.2.3 LAN Emulation Server (LES)

The LES implements the control coordination function for the ELAN. The LES provides the service of registering and resolving MAC addresses to ATM addresses. A LEC registers its own address with the LES. A LEC also queries the LES when the client wishes to resolve a MAC address to an ATM address. The LES either responds directly to the client or forwards the query to other clients so they may respond. There may be more than one instance of an active LES per ELAN.

## 3.2.4 Broadcast and Unknown Server (BUS)

Unlike traditional shared-media LAN architectures such as Ethernet or Token Ring, ATM is connection based. Therefore, it has no built-in mechanism for handling connectionless traffic such as broadcasts, multicasts, and unknown unicasts. In an ELAN, the BUS is responsible for servicing these traffic types by accepting broadcast, multicast, and unknown unicast packets from the LECs via dedicated point-to-point connections, and forwarding the packets to all of the members of the ELAN using a single point-to-multipoint connection. (Unknown unicast packets are packets that the sending station broadcasts because it does not yet know the ATM address for the packet's destination MAC address. There may be more than one instance of an active BUS per ELAN. Using *ForeThought* 5.2 each BUS must be a colocated BUS (also referred to as an intelligent BUS or a LES/BUS pair), which allows the BUS to use the LES's registration table to direct unicast traffic.

# 3.3 Emulated LAN Operation

This section describes the operation of an ELAN and its components from the point of view of a LEC. The operation of an ELAN may be divided into three phases:

- 1. Initialization
- 2. Registration and Address Resolution
- Data Transfer

ELAN components communicate with each other using ATM connections. LECs maintain separate connections for traffic control functions and data transfer. The following connection types are used by the LEC when operating in an ELAN:

- *Configuration-Direct Connection*: a temporary bidirectional point-to-point VCC set up by the LEC to the LECS.
- Control-Direct Connection: a bidirectional point-to-point VCC set up by the LEC to the LES. This connection must be maintained for the duration of the LEC's participation in the ELAN.
- *Control-Distribute Connection*: a unidirectional point-to-multipoint VCC set up by the LES to the LEC. This connection must be maintained for the duration of the LEC's participation in the ELAN.
- *Multicast-Send Connection*: a bidirectional point-to-point VCC set up by the LEC to the BUS for sending multicast data to the BUS. The LEC must attempt to maintain this connection while participating in the ELAN.
- *Multicast-Forward Connection*: a unidirectional point-to-multipoint VCC set up from the BUS to LECs participating in the ELAN. The LEC must attempt to maintain this connection while participating in the ELAN.
- Data-Direct Connection: a bidirectional point-to-point VCC set up between LECs that want to exchange unicast data traffic, and torn down after 20 minutes (default) of inactivity. Each LEC normally establishes many Data-Direct Connections.

For the following discussion, please refer to Figure 3.2.

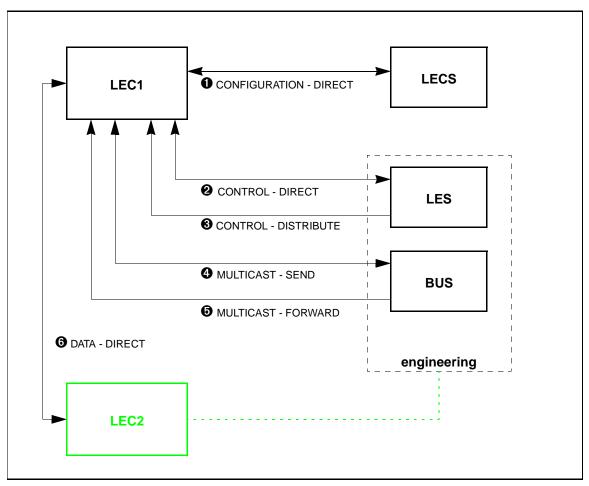


Figure 3.2 - ELAN Operation

#### 3.3.1 Initialization

Once it knows the location of the LECS, LEC1 establishes a configuration-direct connection to the LECS. When connected, the LECS provides LEC1 with the information necessary to connect to the ELAN it wishes to join. This information includes such parameters as: the ATM address of the ELAN's LES, the type of LAN being emulated, the maximum packet size, and the name of the ELAN (engineering, for example). This configuration information is contained in a configuration file that must be built and maintained by the network administrator.



Detailed information about the LECS configuration file may be found in Section 3.6.1.

## 3.3.2 Registration and Address Resolution

After obtaining the address of the LES, LEC1 establishes a control-direct connection **②** to the LES.



When using DLE, this address is a single, anycast address which allows the LEC to reach one of the other DLE peer servers for its ELAN if its local server goes down. This address is routed via PNNI to the nearest active DLE peer server for this ELAN.



If the LES is configured to perform ELAN access control (see Section 3.5), upon receiving a request from a LEC to join the ELAN, the LES sends a message to the LECS to verify that the LEC is allowed to join. If verification is received from the LECS, then the LES gives the LEC permission to join. If verification is not received from the LECS, the LES rejects the join request and the LEC is dropped.

The LES assigns LEC1 a unique identifier, and LEC1 registers its own MAC and ATM addresses with the LES. (The LES maintains a table containing the MAC addresses and corresponding ATM addresses of all members of the ELAN.) At this point, LEC1 has "joined" the ELAN.

The LES then establishes a control-distribute connection ③ back to LEC1. Connections ② and ③ can now be used by LEC1 to send LAN Emulation ARP (LE\_ARP) requests to the LES, and receive replies.

LEC1 now sends an LE\_ARP request to the LES to get the ATM address of the BUS corresponding to the broadcast MAC address (FF-FF-FF-FF-FF). The LEC then establishes a multicast-send connection **4** to the BUS. The BUS responds by setting up a multicast-forward connection **5** to the LEC.

At this point, the LEC is ready to transfer data.

#### 3.3.3 Data Transfer

When LEC1 receives a network-layer packet from a higher layer protocol to transmit to some destination MAC address (for example, LEC2), LEC1 initially does not know the corresponding ATM address of the destination. Consequently, LEC1 transmits an LE\_ARP request to the LES.



The example shown in Figure 3.2 assumes that LEC2 has already registered with the LES, and that connections similar to those described for LEC1 already exist.

While waiting for the LES to respond, LEC1 forwards the packet to the BUS. The BUS broadcasts the packet to all LECs on the ELAN. This is done to avoid data loss, and to minimize connection set-up latency (due to the LE\_ARP process) that may not be acceptable to some network protocols.

If the LE\_ARP response is received, LEC1 establishes a data-direct connection **6** to the destination address of LEC2. This path will be used for subsequent data transfers. Before LEC1 begins to use this connection, it first sends a "flush" packet via the BUS to the destination, LEC2. When LEC2 acknowledges receipt of this packet, signifying that the BUS path is empty, only then does LEC1 begin to use the data-direct connection **6** for data transfer. This process ensures that the network protocol's frames arrive in the proper order.

If no response is received to the LE\_ARP, LEC1 continues to send data via the BUS, while continuing to LE\_ARP until a response is received and a data-direct connection to LEC2 is established.

If LEC1 already has a data-direct connection to a MAC address it wishes to reach, it need not go through the LE\_ARP process again. Instead, it continues to use the current connection. This is possible because each LEC maintains a cache of MAC address to ATM address mappings that it receives in response to the LE\_ARPs it has sent. Entries in this cache are "aged" out over a period of time. Data-direct connections are also cleared if they remain inactive for a period of time.

#### 3.4 Distributed LAN Emulation

Distributed LAN Emulation (DLE) allows the LES and BUS functions that are provided to each ELAN to be distributed among multiple, interconnected server platforms. In this way, DLE provides these ELANs with resiliency and scalability.

To understand DLE operation, it is useful to compare DLE to the current LANE service model, which uses a single LES and BUS for each ELAN. This section first describes a simple example of the single server model and then gives a detailed overview of the DLE model.

## 3.4.1 Single Server LANE Services Model

Figure 3.3 shows the topology of a single server supporting an ELAN. In this example, the LECs are hosts that are using IP, and the LES and BUS are running on the same switch. Three LANE LECs are all registered in the same ELAN called Eng, and each is, therefore, connected to a LES and to a BUS for that ELAN.

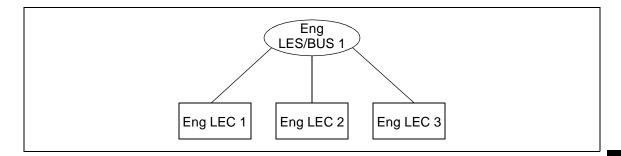


Figure 3.3 - Single Server LANE Services Model

#### 3.4.1.1 Using a Single Server

When LEC 1 wants to contact LEC 3, several messages are exchanged. First, LEC 1 attempts to learn the MAC address of LEC 3 by broadcasting an IP-ARP request with LEC 3's IP address. As Figure 3.4 shows, this ARP request is sent in two steps: ① as a point-to-point message from LEC 1 to the LANE BUS, then ② as a point-to-multipoint message from the BUS to all of the LECs registered in the ELAN.

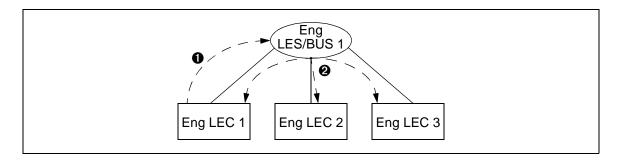


Figure 3.4 - Broadcast IP-ARP Request

When LEC 3 receives the IP ARP request, it recognizes that it is the intended destination, and, therefore, attempts to send an IP ARP response to LEC 1 (whose MAC address was supplied in the ARP request packet).

As shown in Figure 3.5, the delivery of the ARP response is a three-step process: **3** LEC 3 sends an LE-ARP query to the LES, asking for the ATM address that corresponds to LEC 1's MAC address; **4** the LES sends an LE-ARP response to LEC 3; and **5** LEC 3 establishes a circuit to LEC 1's ATM address.

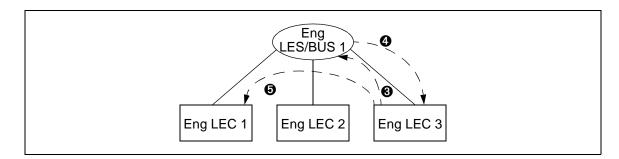


Figure 3.5 - IP ARP Response Handling

#### 3.4.1.2 Limitations of a Single Server

Because the there is only one LES/BUS supporting the ELAN, the following limitations exist:

- The number of LECs in a single ELAN is limited by the number of virtual circuits that the single LES/BUS can establish through their platform's ATM port. This usually limits the ELAN to about 500 LECs.
- Clusters of LECs that are geographically separated from the LES/BUS may have poor throughput, even when connecting to each other, because address queries and broadcasts may traverse slow wide-area links.
- A failure of the LES or BUS brings down the ELAN.

#### 3.4.2 Distributed LAN Emulation Model

To address the limitations of the single server model, DLE distributes the LANE services load among a mesh of LES/BUS DLE peer servers, as shown in Figure 3.6.

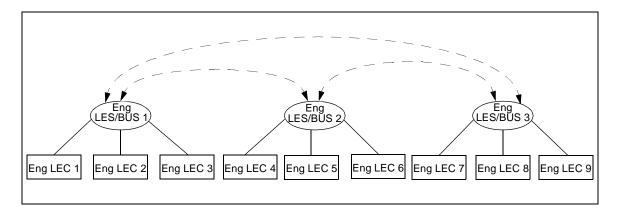


Figure 3.6 - Distributed LAN Emulation Model

Each DLE peer server actually maintains two sets of connections: one is a point-to-multipoint connection to each of its peers for broadcasting multicast data and flooding control information, and the other includes individual point-to-point connections to each peer for directed control traffic.

Each DLE peer server that supports the ELAN is responsible for registering and giving reports about the LECs that are attached to it directly. Each DLE peer server propagates this information to both its locally attached LECs and its peers.



Each device running a DLE peer server must use *ForeThought* 5.0 or greater; however, the DLE peer servers support clients and attached switches using *ForeThought* 4.0 and 4.1, and third-party devices that are ATM Forum LANE 1.0 compliant.

#### 3.4.2.1 Using DLE

Figure 3.7 shows how a connection begins to be established through DLE peer servers. LEC 1 wants to communicate with LEC 9, which is in the same ELAN, but is locally attached to a different DLE peer server. First, **①** LEC 1 sends an IP ARP broadcast request to its local DLE BUS. Then, **②** the BUS broadcasts the packet to both its locally attached LECs and its DLE peer servers.

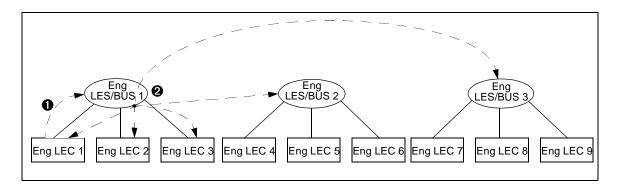
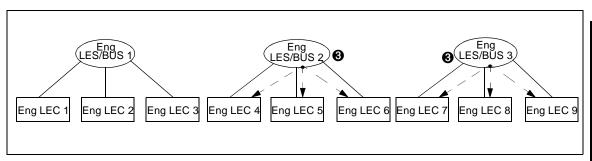


Figure 3.7 - IP ARP Broadcast from LEC 1 to LEC 9

Upon receiving the broadcast from the first DLE peer server, the peers re-distribute the packet to their own locally attached LECs **3**, as shown in Figure 3.8, so the packet arrives its actual destination at LEC 9.



**Figure 3.8 -** Re-distributing the Broadcast across DLE Peer Servers



The peers do <u>not</u> re-distribute the packet to other peers; this would create a loop.

LEC 9 recognizes its IP address, and prepares an IP ARP response. As shown in Figure 3.9, it then sends an LE-ARP request to its local LES **4**, asking for the ATM address that matches LEC 1's MAC address. Since LEC 9's local LES does not have an entry for LEC 1, the local LES passes the query along to all of its locally-attached proxy LECs (none are shown in this figure) and all of its DLE peer servers **5**.

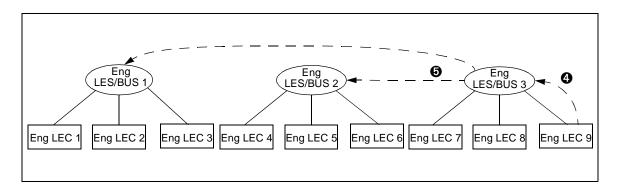


Figure 3.9 - LE-ARP for Unknown Host Sent to Proxies (not shown) and DLE Peer Servers

In Figure 3.10, the second DLE peer server is attached to two proxy LECs (LEC 4 and LEC 5). When the DLE peer server receives the LE-ARP query, it cannot resolve the query, so the DLE peer server re-distributes the query to its proxy LECs **6** (but not to its peer servers again, to avoid a loop). Meanwhile, the first peer server has been able to resolve the LE-ARP for the address of LEC 1 and has sent an LE-ARP response to the third server **6**.

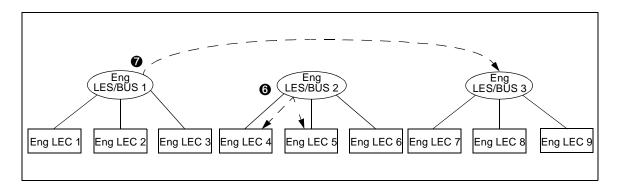


Figure 3.10 - LE-ARP Query Answered by One DLE Peer Server and Re-distributed by Another

When the third DLE peer server receives the LE-ARP response, it passes it directly to LEC 9 (which sent the original query) **3**. The third DLE peer server also caches the registration information for LEC 1 so that other local LECs do not have to go through the entire process again. However, this cache ages out over time. LEC 9 can now open a connection to LEC 1, and send its IP ARP response **9**, as shown in Figure 3.11.

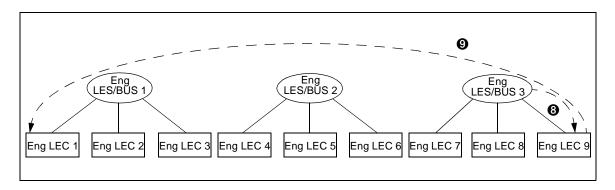


Figure 3.11 - LE-ARP Response Delivered and LEC 9 Contacts LEC 1

#### 3.4.2.2 Advantages of DLE

As mentioned earlier, using DLE provides solutions to the problems of using a single server.

#### 3.4.2.2.1 Load Sharing

DLE peer servers distribute the circuit and processing load. The number of LANE LECs is no longer limited by the number of circuits one LES/BUS platform can maintain, since many platforms can support a single ELAN. Also, more VCs are available for use by other applications.

#### 3.4.2.2.2 Improved Performance for Remote LECs

With DLE, broadcast delivery and LE-ARP resolution across peer servers can take a little longer than if all LECs were connected to a single server, since extra processing steps and transmissions are needed. However, ELANs with groups of LECs in different locations can be designed for higher performance by providing a DLE peer server with each group. Broadcasts and address resolution within each group will improve.

#### 3.4.2.2.3 Fault Tolerance

Perhaps the most important advantage of DLE is fault tolerance. In a single server ELAN, the server can be a single point of failure. If the server fails, endstations in the ELAN are unable to discover each other through broadcast queries and unable to resolve MAC addresses into ATM addresses. Increased network reliability, therefore, requires that ELANs have backups for LES and BUS functions. To illustrate this point, the single server model is again discussed.

#### 3.4.2.2.3.1 Single Server ELAN

Figure 3.12 shows a single server ELAN composed of nine LECs attached to three different switches. The LECS and the LES/BUS are attached to a host connected to a single switch. The process for LEC 1 to connect to the LANE services takes several steps:

- 1. LEC 1 asks the signalling software on its switch to open a connection to the "well-known" LECS address. (Other addressing methods could also be used).
- 2. The signalling software knows that this address is attached to port N (the port on which the host resides on the switch), and opens a circuit between LEC 1 and port N.
- 3. LEC 1 sends a message to the LECS, asking for the address of the LES for LEC 1's ELAN. The LECS responds with the ATM address of the LES, and LEC 1 establishes a circuit to the LES and then the BUS.

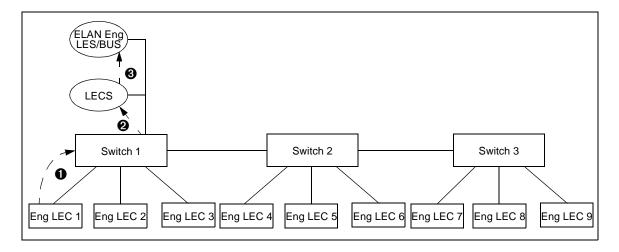


Figure 3.12 - ELAN with a Single Server and Multiple Switches Connecting to Services

In Figure 3.13, when LEC 7 goes through the same process, it is slightly more complicated. When LEC 7 asks its local switch's signalling software to establish a circuit to the LECS **①**, the local switch must use inter-switch link information (IISP or PNNI tables) to establish a cross-switch circuit to the LECS. Once this circuit is established, however, the process is identical.

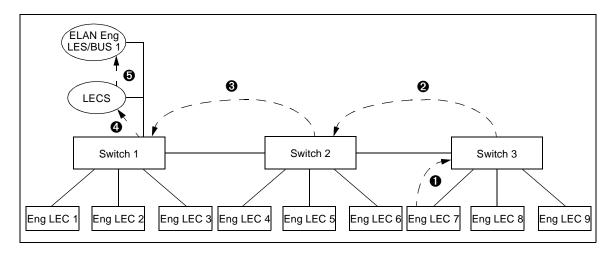


Figure 3.13 - ELAN with Single Server and Remote Connection to Server

Figure 3.14 shows the ELAN in operation after three LECs have gone through the registration process.

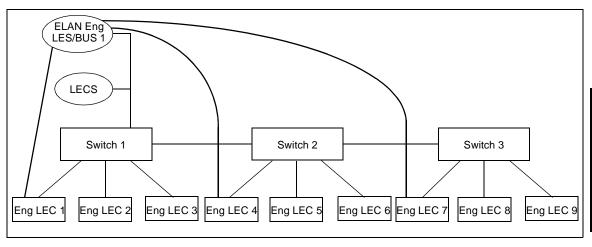


Figure 3.14 - ELAN with Single Server in Operation

If the single server in Figure 3.14 goes down, then the entire ELAN goes down. At this point, the administrator must intervene and reconfigure the ELAN.

#### 3.4.2.2.3.2 DLE ELAN

As noted previously, having a single server supporting an ELAN has a potential problem because the server can be a single point of failure. However, DLE can address this problem. By attaching the ELAN LECs to multiple DLE peer servers which communicate with each other as described earlier, the number of LECs affected by a server failure is reduced, and a backup server is provided for affected LECs to use. Figure 3.15 shows the configuration of such an ELAN as three stations register.

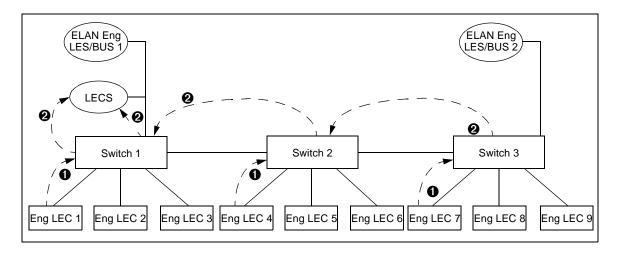


Figure 3.15 - Registrations on an ELAN with Multiple Servers

LECs 1, 4, and 7 are directed by their switches to the LECS. The result is shown in Figure 3.16.



The connection between the two servers carries broadcasts and LE-ARP traffic as described earlier.

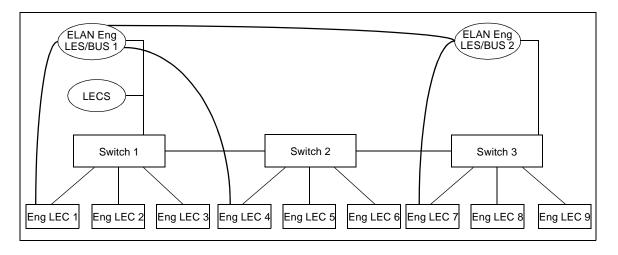


Figure 3.16 - ELAN with Multiple Servers in Operation

This ELAN may experience significant performance improvements for the reasons described earlier. Even if the actual performance is similar to using a single server in a particular network, a great advantage is gained through its fault-tolerance if one of the servers fails as depicted in Figure 3.17.

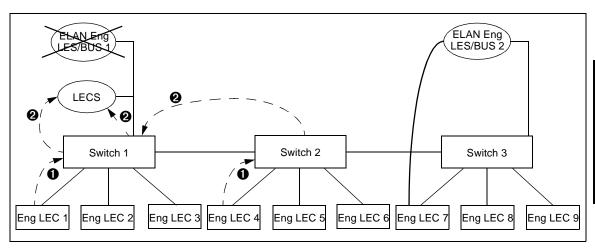


Figure 3.17 - Failure of One ELAN Server and the Recovery Process

The failure and recovery process occurs as follows:

- 1. Eng LES/BUS 1 has lost power. All circuits connected to it are torn down. Low-level signalling traffic (e.g., SSCOP messages) stop, and Switch 1 removes the address of Eng LES/BUS 1 from its link tables.
- 2. LECs 1 and 4 had been connected to Switch 1. They detect that their connections to Eng LES/BUS 1 have been torn down; user intervention is not necessary.
- 3. LECs 1 and 4 follow LANE 1.0 protocols to locate an LECS again to find the address of their ELAN's LES. In this example, they again connect to the LECS.
- 4. The LECS reports to LECs 1 and 4 that their ELAN server is at ATM address N1. This address is used by every peer LES supporting the ELAN; both Eng LES 1 and Eng LES 2 in this example.
- 5. LEC 1 sends a request to Switch 1 to establish a connection to address N1. Switch 1 no longer believes it has a direct connection to N1, and instead uses PNNI to establish a circuit through Switches 2 and 3 to Eng LES 2.
- 6. LEC 4 sends a request to Switch 2 to establish a connection to N1. Switch 2 may have learned from Switch 1 that it no longer offers a connection to N1, or Switch 2 may attempt a route through Switch 1 and be "bounced back" through ATM Forum PNNI crankback. Either way, Switch 2 finally routes the connection through Switch 3 to Eng LES 2.

This recovery process occurs quickly -- clients typically recover at a rate of 100 clients per minute -- and the result is a re-configured ELAN as shown in Figure 3.18.

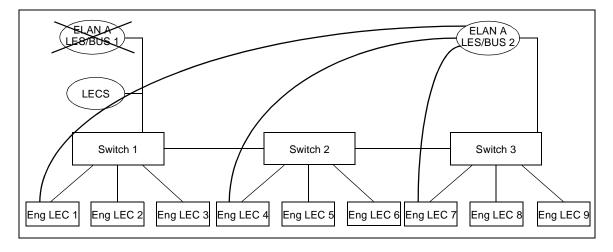


Figure 3.18 - ELAN Re-established Using the Second Server

#### 3.5 ELAN Access Control

Basic ATM Forum LAN Emulation Servers do not guard against unauthorized users learning an ELAN's LES address and then joining the ELAN. However, a method of authorization checking is available in *ForeThought* 5.2.x. After a LEC obtains the address of its LES, the LEC sends a request to the LES to join the ELAN. If the LES has ELAN access control enabled, it sends a message to the LECS to verify that the LEC is allowed to join. If verification is received from the LECS, then the LES gives the LEC permission to join. If verification is not received from the LECS, the LES rejects the join request and the LEC is dropped.

Using this feature, an authorization check is also performed each time the LECS reloads the LECS configuration file. (The LECS periodically checks whether its configuration file has been modified, and, if it has, the file is re-read. The length of this period, in seconds, is defined by the Reload\_Period key.) If the file has changed to disallow some clients that were previously allowed, those clients will be dropped from the ELAN.



ELAN access control also works with a thirdparty LECS. The LES revalidates the client every 600 seconds since the third-party LECS will not contact the LES with configuration changes.

You can enable ELAN access control when you are creating the LES. When you use the **conf** lane les new command, specify the -secure option. This indicates you want to activate a secure LES/BUS pair. (ELAN access control is disabled by default.)

```
myswitch::configuration lane les> new <LES Selector Byte (HEX)> <LES name>\
    [-bus <BUS Selector Byte (HEX)>]\
    [-type (ethernet | token-ring)] \
    [-mtu (1516 | 1580 | 4544 | 9234 | 18190)] \
    [-secure wka | <LECS ATM Address>] \
    [-registertlvs (enable | disable)] \
    [-anycast <LES Anycast ATM Address>]
    [-peers <atm-addr> ...]
```

If you enter wka with the -secure option, the ATM Forum well-known LECS address is used. In this case, you do not have to type the actual well-known address. However, if you are using an LECS address that is different than the well-known address, then you must type the full LECS ATM address to be used.

If you want to disable ELAN access control, or if you want to enable ELAN access control at a later time after the LES has been created, you can use the <code>conflamelessecurity</code> command to do this. See the *AMI Configuration Commands Reference Manual* for more information about this command. By using this command, you do not have to delete and recreate the LES.

# 3.6 Configuring an ELAN

There are different instructions for configuring an ELAN, depending on how your network is currently configured. Please read the following list to determine which set of instructions to use.

- If you had previously configured LANE, you want to upgrade some or all of the clients to *ForeThought* 5.2.x, and you want to upgrade all the equipment that is running services to *ForeThought* 5.2.x using DLE, use the instructions found in Section 3.7.
- If you had previously configured LANE, you want to leave the clients running *ForeThought* 4.1.x, and you want to upgrade all the equipment that is running services to *ForeThought* 5.2.x **without** using DLE, use the instructions found in Section 3.8.
- If you are configuring LANE for the first time and all of your equipment is running *ForeThought* 5.2.x, use the instructions that follow here in Section 3.6.

To configure an ELAN on a switch, you must log into AMI on a switch running *ForeThought* 5.2 and use the commands found under configuration lane.



More information about each of these commands may be found in the *AMI Configuration Commands Reference Manual*.

There are three major steps that the system administrator should follow in order to configure and maintain ELANs:

- 1. Configure an LECS configuration database file.
- 2. Start the LAN Emulation Services (LECS and LES/BUS).
- 3. Start the LEC(s) and join an ELAN.



Steps 1 may be performed using a text editor on any system. However, the resulting file can be **used** only on systems running under Solaris 2.5, 2.5.1, or 2.6, or on a FORE switch running *ForeThought* 4.0 or greater.

The remainder of this section gives a practical example of configuring and administering an ELAN using *ForeThought* 5.2.

# 3.6.1 Configuring an LECS Configuration Database File

The LECS uses a text configuration file to contain the configuration information needed by LECs that wish to participate in an ELAN. The LECS configuration file may be built and edited using a text editor such as vi or emacs.



ForeThought VLAN Manager, a product available separately from FORE Systems, greatly simplifies the administration of ELANs. Its graphical user interface controls the content of the LECS configuration file transparently to the user. For more information, please refer to the ForeView VLAN Manager User's Manual.

### 3.6.1.1 Before You Begin

Before building or modifying the LECS configuration file, you should first determine the topology of the ELAN or ELANs that you want to administer. You must supply the following information when building or editing the LECS configuration file:

- Provide the name of each ELAN (engineering, marketing, etc.).
- Provide the LAN type (Ethernet or Token Ring) for each ELAN.
- Provide the MTU size for each ELAN.
- Provide the ATM address of the LES for each ELAN. If you are using DLE, this address must be the anycast address for the DLE peer servers in each ELAN. Be sure to choose a distinct anycast address for each ELAN in the network. It must be unique within the first 19 bytes.
- Provide the address of each LEC that may participate in each ELAN.
- Provide the MPOA control parameters if you wish to run MPOA.
- If you wish LECs to use a default ELAN, the default LES information must also be included.
- Provide various other configurable parameters.

#### **CAUTION**



Do not attempt to edit an existing functional LECS configuration file without first making a backup copy of the file. Incorrect modification of the configuration file could result in loss of communication between one or more members of a defined ELAN.



You may make changes to the LECS configuration file while the LECS process is running. The configuration file is reread periodically by the LECS process (the default period is ten minutes). Consequently, any changes that you make to the configuration file are not recognized until the file is reread.

### 3.6.1.2 LECS Configuration File Syntax

Each line that you enter in the configuration file takes the general form:

[[group].]key : value

The group field may represent:

- ELANs (by name) ELAN names are case-sensitive, and may not exceed 32 characters in length
- clients ATM or MAC addresses
- miscellaneous LECS control information specified by using a group name of LECS

The key field is used to denote an individual parameter within a group.

The value field contains the value assigned to the key.

Omitting the group implies that the key and value apply to all groups in the configuration file. Leading and trailing spaces, as well as spaces on either side of the ":", are ignored.

For example, to specify a maximum frame size of 1516 bytes for the ELAN named engineering, enter the following:

engineering.Maximum Frame Size: 1516

Similarly, to specify a default maximum frame size of 1516 bytes for all ELANs defined in a given configuration file, enter the following:

.Maximum Frame Size : 1516

Table 3.1 defines the various key parameters that may be entered in the configuration file. The acceptable range of values and the default value for each parameter is also given.

**Table 3.1 -** LECS Configuration File Parameters

Parameter	Definition		
.LAN_Type: Ethernet/IEEE 802.3	Identifies the type of ELAN, either Ethernet/IEEE 802.3 or IEEE 802.5. The default is Ethernet/IEEE 802.3.		
.Maximum_Frame_Size: 1516	Specifies the length (in number of bytes) of the largest frame. Selections are: 1516, 1580, 4544, 9234, and 18190. The default is 1516 for Ethernet and 4544 for Token Ring.		
.Control_TimeOut: 120	Specifies the timing out of request/response control frame interactions, in seconds. The minimum is 10 seconds and the maximum is 300 seconds. The default is 120 seconds.		
.Maximum_Unknown_Frame_Count: 1	Limits the number of unicast frames sent to the BUS. The minimum is 1 frame per 60 seconds and the maximum is 10 frames per second. The default is 1 frame per second.		
.Maximum_Unknown_Frame_Time: 1	Limits the number of unicast frames sent to the BUS in the specified number of seconds. The default is 1 second.		
.VCC_TimeOut_Period: 1200	Specifies the length of time that an idle data connection remains open before being closed. The default value is 1200 seconds.		
.Maximum_Retry_Count: 1	Limits the number of LE_ARP retransmission requests. The minimum is 0 and the maximum is 2. The default is 1.		
.Aging_Time: 300	Specifies the period that LE_ARP cache table entries remain valid, in seconds. The minimum is 10 and the maximum is 300. The default value is 300 seconds.		
.Forward_Delay_Time: 15	Specifies the timing out of non-local ARP cache entries in seconds. The minimum is 4 and the maximum is 30. The default value is 15 seconds.		
.Expected_LE_ARP_Response_Time: 1	Specifies the maximum time a LEC expects an LE_ARP request/response will take, in seconds. The minimum is 1 and the maximum is 30. The default value is 1 second.		
.Flush_TimeOut: 4	Specifies the maximum time a LEC expects an LE_FLUSH request/response will take, in seconds. The minimum is 1 and the maximum is 4. The default value is 4 seconds.		

**Table 3.1 -** LECS Configuration File Parameters (Continued)

Parameter	Definition		
.Path_Switching_Delay: 6	Specifies the minimum time between switching BUS an data paths, in seconds. The minimum is 1 and the max mum is 8. The default value is 6 seconds.		
.Local_Segment_ID	The segment ID of the emulated LAN for an IEEE 802.5 source routing bridge.		
.ELAN_Identifier: auto	Specifies a non-zero ELAN identifier. The default is auto, which signals the LES to generate an ELAN identifier itself.		
.Multicast_Send_VCC_Type: Best Effort	Specifies the multicast send mode, either Best Effort, Variable, or Constant. The default is Best Effort.		
.Connection_Complete_Timer: 4	Specifies the time period in which data or READY_IND is expected, in seconds. The minimum is 1 and the maximum is 10. The default is 4 seconds.		
.ShortCut_Protocols: IP	Specifies the set of protocols on which to perform flow detection. Also, specifies the set of protocols for which MPOA resolution is supported. The default is {}. Currently, the only supported value is IP.		
.ShortCut_Threshold: 10/1	Specifies the number of frames per second that a LANE/MPOA client (LEC/MPC) forwards to the same destination via the default forwarding path before which it should begin using a shortcut. The minimum number of frames is 1 frame and the maximum is 65,535 frames. The minimum rate of speed at which the frames are forwarded is 1 second and the maximum is 60 seconds. The default is 10 frames per second.		
.Resolution_Initial_Retry_Time: 5	Specifies the initial retry time interval after which a LEC/MPC may send another MPOA Resolution Request if an MPOA Resolution Reply has not been received for the initial request. The minimum is 1 second. The maximum is 300 seconds. The default is 5 seconds.		
.Resolution_Maximum_Retry_Time: 40	Specifies the maximum retry time interval after which a LEC/MPC assumes an MPOA Resolution Request has failed. The minimum is 10 seconds. The maximum is 300 seconds. The default is 40 seconds.		

**Table 3.1 -** LECS Configuration File Parameters (Continued)

Parameter	Definition		
.Resolution_Hold_Down_Time: 160	Specifies the minimum amount of time to wait before reinitiating an MPOA Resolution Request after a failed resolution attempt. This value is usually greater than the Resolution_Maximum_Retry_Time. The minimum is 30 seconds. The maximum is 1,200 seconds. The default is 160 seconds.		
.MPOA_KeepAlive_Time: 10	Specifies how often an MPS must send MPOA Keep-Alive messages to all LEC/MPCs for which it has created and is maintaining Egress Cache Entries. The minimum is 1 second. The maximum is 300 seconds. The default is 10 seconds.		
.MPOA_KeepAlive_Lifetime: 35	Specifies the length of time that a Keep-Alive message is considered valid. It is recommended that this value be at least twice the value of the MPOA_KeepAlive_Time. The minimum is 3 seconds. The maximum is 1,000 seconds. The default is 35 seconds.		
.Resolution_GiveUp_Time: 40	Specifies the minimum amount of time to wait before an MPS should give up on a pending resolution request. The minimum is 5 seconds. The maximum is 300 seconds. The default is 40 seconds.		
.Resolution_Default_Holding_Time: 1200	Specifies the default holding time for use in NHRP Resolution Replies. The minimum is 60 seconds. The maximum is 7,200 seconds. The default is 1,200 seconds.		
<pre>.IP_Flows: <ip-flow> {, <ip-flow>}</ip-flow></ip-flow></pre>	Specifies a set of IP flows that are defined by the parameters listed below.		
<ip-flow></ip-flow>	Specifies the IP information that is defined by the parameters listed below. The format is <pre><pre><pre><pre><pre><pre><pre><pre></pre></pre></pre></pre></pre></pre></pre></pre>		
<src-dst></src-dst>	Specifies the source and destination addresses of an IP flow. The format is <ip-addr> <ip-addr> with the source address first and the destination address second.</ip-addr></ip-addr>		
<ip-addr></ip-addr>	Specifies the format for the IP addresses <ip-addr> used above. The format is <dotted-addr>/<pre>/<pre>prefix-leng&gt;.</pre></pre></dotted-addr></ip-addr>		

**Table 3.1 -** LECS Configuration File Parameters (Continued)

Parameter	Definition
<dotted-addr></dotted-addr>	Specifies the format for the IP addresses <dotted-addr> used above. The format is <octet>.<octet>.<octet>.<ctet>. The range of each octet is from 0 to 255.</ctet></octet></octet></octet></dotted-addr>
<pre><prefix-leng></prefix-leng></pre>	Specifies the prefix length to indicates how much of a given IP address is significant. The range is 0 through 32. 0 means any IP address can be matched and 32 means the entire IP address is significant.
<pre><pre><pre><pre></pre></pre></pre></pre>	Optionally specifies for which protocol this flow applies (ICMP, IGMP, TCP, or UDP). The format is one of: ICMP, IGMP, TCP <port> <port>, or UDP <port> <port>. If no protocol is specified, then this flow applies to all protocols.</port></port></port></port>
<port></port>	Specifies the source and destination port number, respectively, to be used with the <pre><pre>ctively</pre>, to be used with the <pre>cproto&gt;</pre> parameter. 0 indicates a match for any port number. The range is 0 through 65535.</pre>
<thresholds></thresholds>	Specifies a traffic threshold that must be reached before a shortcut VCC is established for a flow. Also, specifies a threshold below which a shortcut VCC is no longer considered valid, is torn down, and the default forwarding path is used. Both thresholds are expressed in frames per second. The format is <threshold> <threshold> <threshold> <threshold> <threshold> down threshold first, the tear-down threshold second, and the set-up threshold for the default forwarding path third.</threshold></threshold></threshold></threshold></threshold>
<threshold></threshold>	Specifies the format for entering the threshold data as described for <thresholds>. The format is <frame-count>/<frame-time>. For set-up, a zero count with a non-zero time means a shortcut VCC should be established upon the first frame and a non-zero count with a zero time means a shortcut VCC is never established. For tear-down, a zero count with a non-zero time means a shortcut VCC is never torn down.</frame-time></frame-count></thresholds>
<frame-count></frame-count>	Specifies the format for the number of frames for the set- up and tear-down thresholds. The range is 0 through 65535.

**Table 3.1 -** LECS Configuration File Parameters (Continued)

Parameter	Definition
<frame-time></frame-time>	Specifies the format for the number of seconds for the set-up and tear-down thresholds. The range is 0 through 65535.
<qos></qos>	Specifies the parameters used in signalling a connection for the specified flow. The format is <qos-flags> <qos-class>. Currently, <qos-flags> is unused and its value should be 0.</qos-flags></qos-class></qos-flags>
<qos-class></qos-class>	Specifies the Quality of Service (QoS) class for this flow. Shared specifies a UBR connection will be shared among all flows to the same destination. UBR specifies a non-shared UBR connection. The format is UBR <rate>. CBR specifies a CBR connection. The format is CBR <rate>. VBR specifies a VBR connection. The format is VBR <rate> <rate> <burst> (for CLP=0+1 cells) <rate> <burst> (for CLP=0 cells). ABR specifies an ABR connection. There is no format currently defined for ABR.</burst></rate></burst></rate></burst></rate></burst></rate></burst></rate></burst></rate></burst></rate></burst></rate></rate></rate></rate>
<rate></rate>	Specifies the rate for the QoS class specified. For UBR and CBR, this is the Peak Cell Rate (PCR) in cells per second. For VBR and NRT-VBR, this is the PCR for cells that are CLP=0+1 in cells per second, the Sustainable Cell Rate (SCR) for cells that are CLP=0+1 in cells per second, and the SCR for cells that are CLP=0 in cells per second.
<burst></burst>	Specifies the Maximum Burst Size (MBS) in cells for the QoS class specified (VBR and NRT-VBR only). This is the MBS for cells that are CLP=0+1, in cells, and the MBS for cells that are CLP=0, in cells.

Lines beginning with # may be inserted if you wish to include comments or to improve the clarity of the presentation when the file is viewed or printed. These lines are ignored when the file is read. Lines may be continued by escaping the end-of-line with a backslash  $^{\circ}$  (do not enter the quote marks).

### 3.6.1.3 Defining an ELAN

Each ELAN is defined by an address statement whose value denotes the ATM address of the ELAN's LES. For example:

engineering.Address: c5.0005.80.ffe100.0000.f21a.01b9.0020480605b2.00



To configure DLE for an ELAN, use the anycast address in this statement. Be sure to use a distinct anycast address for each ELAN in the network. It must be unique within the first 19 bytes.

In addition, you may instruct a given ELAN to override any of the default values. For example, the engineering ELAN could override a Maximum\_Frame\_Size setting of 1516; thus:

engineering.Maximum\_Frame\_Size: 4544

If you want to control which clients may or may not join a given ELAN, two additional keys, Accept and Reject, whose values are comma-separated lists of matching elements, may be used.



The Accept values must be supplied if you are going to enable ELAN access control.

These values may be:

a MAC address.

engineering.Accept: 0020480605b2 , 002048080011 , 0020481020ef an ATM address and equal-length bit mask,

or an ATM address containing "don't-care" semi-octets denoted by an "X":

marketing.Accept: 47.0005.80.FFE100.XXXX.XXXX.XXXX.002048XXXXXX.XX

The last two forms of ATM-address matching elements are functionally the same. The latter is shorter but only allows for masks whose semi-octets are all ones or all zeros, while the former allows for arbitrary masks. A prospective-client address is "captured" by an ELAN name if the client's address matches one of the Accept elements but not one of the Reject elements (if present). Finally, an ELAN may be configured to accept any client that wishes to join by including the following statement:

The order in which to apply the Accept and Reject rules is given by a Match.Ordering group.key statement, whose value is a comma-separated list of ELAN names. For example:

Match.Ordering: default, engineering, marketing

The names of all ELANs that have Accept keys must be included in Match. Ordering.

The LE\_CONFIGURE\_REQUEST frame contains an ATM address and an optional MAC address or route descriptor (which is always present for ELAN access control requests). The Accept/Reject checking proceeds in two distinct phases: first, for the MAC address or route descriptor, if present, and second, for the ATM address. So, even though a client might be rejected by its MAC address, it can be accepted by its ATM address. Therefore, when configuring the Accept and Reject rules, ensure that you write them either as explicit lists of only MAC addresses or route descriptors, or only as ATM address matches.

# 3.6.1.4 Defining a Client

Clients need not be defined in the LECS configuration file. Typically, you would define a client for the purpose of overriding one or more of the default configuration parameters for that particular client.

A client is defined by using its ATM address, MAC address, or route descriptor in the group field, and perhaps giving the name of its ELAN as the value of the LAN\_Name key. For example:

Configuration parameter overrides can also be given on a per-client basis. For example, the following statements override the default VCC\_TimeOut\_Period and Aging\_Time configuration parameters for a client whose MAC address is 002048080011 on the engineering ELAN:

002048080011.LAN\_Name:engineering 002048080011.VCC\_TimeOut\_Period:1200 002048080011.Aging\_Time: 30

#### 3.6.1.5 LECS Control Parameters

Specifying values for keys in the LECS group provides control over the operation of the LECS process.



If you change the values of the LECS control parameters while the LECS process is running, the new values do not take effect until the LECS process is stopped, and then restarted.

When a client contacts the LECS, the connections established are known as Configuration Direct VCCs. To override the default value of the VCC\_TimeOut\_Period key (the number of seconds before an idle Configuration Direct VCC is automatically closed by the LECS), enter a statement similar to the following:

LECS.VCC\_TimeOut\_Period: 1200

The LECS periodically checks whether its configuration file has been modified, and, if it has, the file is reread. The length of this period, in seconds, is given by the Reload\_Period key:

LECS.Reload\_Period: 600

The Permanent\_Circuits key holds a comma-separated list of VPI.VCI pairs denoting the local ends of 0.17 PVCs on which the LECS should listen. For example:

LECS.Permanent\_Circuits: 0.42, 0.112

The LECS can provide the client with a fourteen-bit pattern to permute the MAC-address generation algorithm. This bit pattern is specified with the MAC\_Address\_Base key.

LECS.MAC Address Base: 38fe

#### 3.6.1.6 LECS MPOA Parameters

MPOA requests use similar, and in many cases, the same database keys as LANE requests. However, there are some MPOA-specific keys that can be specified. LEC/MPC parameters can be specified for shortcuts. For example, the following group of parameters indicates that shortcuts should be established for IP flows, but only after a threshold of 10 frames per second is surpassed.

```
.ShortCut_Protocols: IP
.ShortCut Threshold: 10/1
```

LEC/MPC parameters can also be specified for resolution requests. When a LEC/MPC sends an MPOA Resolution Request, it sets a timer to a Resolution\_Initial\_Retry\_Time. If an MPOA Resolution Reply is not received in that amount of time, a retry may be sent. Each time a retry is sent, the timer is set to the Resolution\_Initial\_Retry\_Time value \* a retry multiplier. If the value exceeds the Resolution\_Maximum\_Retry\_Time value, the LEC/MPC assumes the request has failed. A new request may not be sent until the Resolution\_Hold\_Down\_Time has been surpassed.

For example, the following parameters indicate that an initial MPOA Resolution Request should be retried after 5 seconds, backing off to a maximum retry time of 40 seconds, and then the MPOA Resolution Request process is re-initialized after 160 seconds.

```
.Resolution_Initial_Retry_Time: 5
.Resolution_Maximum_Retry_Time: 40
.Resolution_Hold_Down_Time: 160
```

MPOA server (MPS) parameters can also be specified. For example, the following parameters indicate that the MPS must send Keep-Alive messages every 10 seconds. Each of these messages is valid for 35 seconds. The MPS must wait 40 seconds before giving up on a pending resolution request, and should use 1200 seconds as the default holding time in NHRP Resolution Replies.

```
.MPOA_KeepAlive_Time: 10
.MPOA_KeepAlive_Lifetime: 35
.Resolution_GiveUp_Time: 40
.Resolution_Default_Holding_Time: 1200
```

Parameters can also be specified for flow descriptors which determine whether and when to trigger creation of shortcut circuits. The LECS sends the LEC/MPC this set of parameters. These parameters consist of the following elements in the following order: a source/destination specifier, flow establishment thresholds, and a QoS descriptor.

For example, you could specify that telnet traffic to the Class C 202.19.88.0 subnet should be sent on a UBR VCC with a peak rate of 10,000 cells per second, but only after the traffic on that connection exceeds 20 frames per second. If the VCC is idle for more than 10 minutes, then the shortcut should be torn down.

The parameters in the IP flow example are defined as follows:

0.0.0.0/0

202.19.88.0/24	The destination address/prefix length.			
TCP 0 23	The protocol to which this flow applies, and the source and destination ports, respectively.			
20/1 1/600 0/1	The set-up threshold, the tear-down threshold, and the set-up threshold for the default forwarding path, respectively.			
0x0 UBR 10000	The QoS flag (currently unused and should be set to 0), the QoS class, and the PCR for the specified QoS class, respectively.			

The source address/prefix length.

# 3.6.2 Sample LECS Configuration File

#### CAUTION



Do not attempt to edit an existing functional LECS configuration file without first making a backup copy of the file. Incorrect modification of the configuration file could result in loss of communication between one or more members of a defined ELAN.



For a detailed discussion of how to configure an LECS configuration file similar to the one given in this section, please refer to Section 3.6.1.

The sample LECS configuration file shown at the end of this section in Figure 3.19 and Figure 3.20 defines three ELANs:

- default
- engineering
- marketing

The Match.Ordering statement specifies the ELAN names in the order that prospective clients will attempt to match. The default configuration parameters are shown with their default values. These values apply to all ELANs in this configuration file, unless overridden for a particular ELAN or client.

ELAN default is configured to accept any client that wishes to join. The ATM address of the default LES is listed in the default. Address statement. If DLE is being configured for ELAN default, then this address must be the anycast address that allows the clients to reach any of the DLE peer servers for ELAN default.



Be sure to choose a distinct anycast address for each ELAN in the network. It must be unique within the first 19 bytes.

ELAN engineering has overridden the default Maximum\_Frame\_Size with a new size of 4544 bytes. Consequently, this frame size applies only to traffic on the engineering ELAN. The default and marketing ELANs continue to use the default frame size of 1516 bytes.

Two LECs, whose MAC addresses are 002048080011 and 0020481020ef, are identified as acceptable clients for the engineering and marketing ELANs.

```
# The search ordering of elan names
Match.Ordering: default, engineering, marketing
# the default configuration parameters
.Control_TimeOut: 120
.Maximum_Unknown_Frame_Count: 1
.Maximum_Unknown_Frame_Time: 1
.VCC_TimeOut_Period: 1200
.Maximum_Retry_Count: 1
.Aging_Time: 300
.Forward_Delay_Time: 15
.Expected_LE_ARP_Response_Time: 1
.Flush_TimeOut: 4
.Path_Switching_Delay: 6
.Multicast_Send_VCC_Type: Best Effort
.Connection_Complete_Timer: 4
.LAN_Type: Ethernet/IEEE 802.3
.Maximum_Frame_Size: 1516
.ShortCut_Protocols:IP
.ShortCut_Threshold:10/1
.Resolution_Initial_Retry_Time:5
.Resolution_Maximum_Retry_Time:40
.Resolution_Hold_Down_Time:160
.MPOA_KeepAlive_Time:10
.MPOA_KeepAlive_Lifetime:35
.Resolution_GiveUp_Time:40
.Resolution_Default_Holding_Time:1200
.IP_Flows: 0.0.0.0/0 202.19.88.0/24 \setminus
            TCP 0 23 \
            20/1 1/600 0/1 \
            0x0 UBR 10000
# Parameters for the default elan
default.Address: 47.0005.80.ffe100.0000.f21a.21b8.0097036324b2.25
```

Figure 3.19 - Sample LECS Configuration File (Part One of Two)

```
#
# Parameters for elan: engineering
#
engineering.Address: 47.0005.80.ffe100.0000.f21a.01b9.0020480605b2.11
engineering.Accept: 002048080011 , 0020481020ef
engineering.Maximum_Frame_Size: 4544
#
# Parameters for elan: marketing
#
marketing.Address: 47.0005.80.ffe100.0000.f21a.01b9.0020480605b2.21
marketing.Accept: 002048080011 , 0020481020ef
#
```

Figure 3.20 - Sample LECS Configuration File (Part Two of Two)

# 3.6.3 Starting the LAN Emulation Services

LAN Emulation services include the LECS and the LES/BUS. Once the LECS configuration database file has been configured, these services must be started so that they are available for LECs to attempt to use. Using *ForeThought* 5.2, the LES/BUS services must run in the same device.

### 3.6.3.1 Starting the LECS

Once an LECS configuration file has been configured, you need to retrieve the LECS configuration database file that you built elsewhere and put it on the switch that is going to run the LECS.

1. Use the following AMI command to retrieve the file:

```
configuration lane lecs get <host>:<remotefile> [<localfile>]
```

For example, you would enter something similar to the following:

```
configuration lane lecs get 198.29.22.46:lecs.cfg
```

2. After you have retrieved the LECS configuration database file, use the following AMI command to start the LECS service on the switch:

```
conf lane lecs new <LECS Selector byte (HEX)> [-db <LECS database
    file>]\[-default <LES atm address>] [<LECS-wka> | none]
```

For example, to start the LECS service using the -db option and using the ATM Forum's well-known address you would enter something similar to the following:

configuration lane lecs new 0x0c -db lecs.cfg



If you want to use an address other than the ATM Forum's address, you would enter that address at the end of the command. If you want to disable the well-known address so that the LECS can only be contacted by using the switch's actual address (with selector byte) type none at the end of the command.

3. Use the following AMI command to verify that the LECS has been started and is running:

#### configuration lane lecs show

Index	AdminStatus	OperStatus	Selector	WKA	Database
1	up	up	0x0c	atm-forum	lecs.cfg

The OperStatus field shows up, meaning that the LECS is running. Now you must start the DLE peer servers as described in the next section.



If you used an address other than the ATM Forum's address, the WKA field would show other and that address would be displayed below the entry. If you disabled the well-known address the WKA field would show none.



If you used the -default option, then that LES address would be displayed below this entry in a field titled Default LES.

### 3.6.3.2 Starting the DLE LES/BUS Peer Servers

The LES and BUS services must be started for the ELAN. This example assumes you are using DLE; therefore, you must enter a distinct anycast address (that is unique within the first 19 bytes) for the LECs to use to contact the LES, and the address of each of the DLE peer servers so that this server can communicate with its peers.



Although this example illustrates how to set up DLE, to facilitate the configuration of DLE, it is recommended that you use *ForeView* instead of AMI.

1. To start the services, use the following AMI command on the switch that is going to run one of the DLE peer servers:

For example, you would enter something similar to the following:

```
conf lane les new 90 engineering -anycast c5.0005.80.ffe100.0000.f21a.3596.0020481a3596.f0 -peers 47.0005.80.ffe100.0000.f21a.10bb.0020481a10bb.90 47.0005.80.ffe100.0000.f21a.3552.0020481a3552.10 47.0005.80.ffe100.0000.f21a.3218.0020481a3218.44
```



This command creates a co-located LES and BUS using a single AMI command. You cannot create a BUS separately using *ForeThought* 5.0 or greater. The conflane bus commands are only useful in providing backwards compatibility with switches that are running earlier versions of *ForeThought* software.



If no ELAN type is entered, the switch assumes Ethernet and uses 1516 as the MTU size. If Token Ring is used as the type, but no MTU size is entered, 4544 is used as the size.



You must enter the address of each of the DLE peer servers when you are starting DLE; e.g., if you want four peers, then all four must be configured with the addresses of the other three peers, as well as their own LES address, at the time that each LES/BUS is started. If you want to add new peers to the list, you can use the conflane lespeeradd command.

2. Use the following AMI command to verify that the LES and the BUS have been started and are running.

#### configuration lane les show

Index AdminStatus OperStatus LesSel Type			MTU	ELAN	SECURE	TLVs		
1	up	up	0x90	ethernet	1516	engineering	disable	enable
	LES	:0x47.0005.8	0.ffe100	.0000.f21a	a.10bb.0	020481a10bb.	90	
	BUS :0x47.0005.80.ffe100.0000.f21a			.10bb.0	020481a10bb.	90 (Co-Lo	ocated)	
		:c000580ffe1	000000f2	1a35960020	)481a359	6f0 (ANYCAST	)	
	PEER	:0x47.0005.8	0.ffe100	.0000.f21a	a.3552.0	020481a3552.	10	
	PEER	:0x47.0005.8	0.ffe100	.0000.f21a	a.3218.0	020481a3218.	44	

The OperStatus field shows up, meaning that the LES and the BUS are running.

3. You should then configure the DLE peer servers. Open an AMI session on one of the other switches that will run a DLE peer server and enter something similar to the following:

```
conf lane les new 10 engineering -anycast c5.0005.80.ffe100.0000.f21a.3596.0020481a3596.f0 -peers 47.0005.80.ffe100.0000.f21a.3552.0020481a3552.10 47.0005.80.ffe100.0000.f21a.10bb.0020481a10bb.90 47.0005.80.ffe100.0000.f21a.3218.0020481a3218.44
```

4. Open an AMI session on the third switch that will run a DLE peer server and enter something similar to the following:

```
conf lane les new 44 engineering -anycast c5.0005.80.ffe100.0000.f21a.3596.0020481a3596.f0 -peers 47.0005.80.ffe100.0000.f21a.3218.0020481a3218.44 47.0005.80.ffe100.0000.f21a.10bb.0020481a10bb.90 47.0005.80.ffe100.0000.f21a.3552.0020481a3552.10
```

Once all of the peers have been created, use the **conf lane les show advanced** command to verify that the peers have established point-to-point and point-to-multipoint connections to each other.

# 3.6.4 Starting the LEC(s) and Joining an ELAN

Now that the ELAN services have been started, you can have LECs join the ELAN that you have created.



The switch software only allows you to create an instance of a LEC on a switch. To create an instance of a LEC on a host, you must use the *ForeRunner* VLAN Manager or use a *ForeRunner* host adapter. Please refer to the respective User's Manual for instructions.



A LEC created on the switch cannot join a Token Ring ELAN. It can only join an Ethernet ELAN.

1. To start a LEC that will attempt to join the ELAN, use the following AMI command on the switch that is going to be a LEC:



The recommended (and default) method for starting a LEC is to use the wellknown mode, meaning that the LEC will attempt to contact the LECS on the "well-known" address as defined by the ATM Forum's LAN Emulation standards (47.0079.00.0000000.0000.0000.0000.000 A03E000001.00).

For example, to start a LEC that attempts to join the ELAN called engineering, enter the following:

configuration lane lec new 2 engineering -ip 192.168.61.25 -mask 255.255.25.0



If you want to use the manual mode, you must enter either a LECS address other than the well-known address or you must enter a LES address. If you enter a LES address, this means that the LEC bypasses the LECS and directly contacts the specified LES.



In order to use DLE, the LES address to be used must be the anycast address of the DLE peer server, not the server NSAP address.

2. Verify that the LEC has joined the ELAN by using the following AMI command:

#### configuration lane lec show

	Admin	Oper					
Index	Status	Status	Sel	Mode	MACaddress	IfName	ELAN
1	up	up	0x02	wellknown	00204815096b	el0	engineering
	LECS:0x						
	LES :c000580ffe1000000f21a35960020481a3596f0						

The OperStatus field shows up, meaning that the LEC has successfully joined the ELAN.



If the OperStatus field shows joining, this means that the LEC is still registering with the ELAN. Wait a few seconds and check it again. When it has finished, the OperStatus field displays up.

You can also use the command conf lane lec show advanced to display more information about the LEC.

3. After the first LEC has joined the ELAN, you can perform the same steps in this section on another switch to allow a LEC to run on that switch. You can also use the VLAN Manager or the host software to add more LECs to this ELAN. Once all the LECs have joined, the ELAN is complete.

# 3.7 Upgrading an ELAN to Use DLE

This section describes how to upgrade your ELAN if you had previously configured LANE, you want to upgrade some or all of the clients from *ForeThought 4.1.x* to *ForeThought 5.2.x*, and you want to upgrade all the equipment that is running services to *ForeThought 5.2.x* using DLE.



If you used the VLAN Manager to create your ELAN, you should use it to upgrade the ELAN. Refer to the *ForeView VLAN Manager User's Manual* for the appropriate steps.

The following basic steps are involved in upgrading your ELAN to use DLE. Each of these steps is described in detail in the following sections. It is recommended that you read the entire section before attempting to upgrade to DLE. These steps describe the commands for a switch. Refer to the corresponding User's Manual for the appropriate steps for commands for hosts, *PowerHubs*, ES-3810s, or the VLAN Manager.

- 1. Edit the LECS.CFG file.
- Delete the old LES and BUS.
- 3. Upgrade the switches that are running services.
- 4. Create the DLE peer servers.
- 5. Transfer the updated LECS.CFG file.
- 6. Restart the LECS.
- 7. Recreate the LECs.
- 8. Create the last DLE peer.
- 9. Add the last DLE peer to each list of peers.
- 10. Change the last failover ELAN information in the LECS.CFG file.
- 11. Transfer the final LECS.CFG file.
- 12. Restart the LECS.

### 3.7.1 Edit the LECS.CFG File



Before you edit the LECS.CFG file, you may wish to back it up to a host using the <code>operflashput</code> command. If you are using TFTP as the transfer protocol (this is the default), the remote host to which the FLASH file will be sent must be running the TFTP server code and must have on an empty file in the <code>/tftpboot</code> directory on the remote host to receive the FLASH file. See Chapter 3 in the <code>ATM Management Interface</code> (AMI) Manual for more information. If you are using FTP as the transfer protocol, you only need to enter the command shown below to perform the transfer.

Transfer the LECS.CFG file that is currently being used by the LECS to a workstation (other than the one to which you backed up the file) so you can edit the file.

```
myswitch::operation> flash put lecs.cfg 169.144.85.195:/tftpboot/lecs.cfg
Transferred 2323 bytes of fs:/lecs.cfg
```

The information about the LES addresses in your old LECS.CFG file may look something like this before editing:

Match.Ordering: Mktg | 0, Mktg | 1

On the workstation that has the LECS.CFG file, use a text editor to make the following changes:

1. Add a new ELAN to the beginning of the Match.Ordering list that is named like the old failover ELANs, but without the |n|. In this example, there are currently two failover ELANS named Mktg |0| and Mktg |1|. Therefore, add an ELAN named Mktg.

- 2. Give the anycast address of the DLE peer servers to the new ELAN. Also, replace the LES address of the old ELANs with this anycast address, except for the last of the |n failover ELANs (in this case, Mktg|1). Leave Mktg|1 with old LES address for now so the LECs can use it until they are all changed over. (It will be changed at the end of the process.)
- 3. If there were any non-colocated BUSs (none shown in this example), delete any lines that refer to them, since each BUS will co-located with a LES after you upgrade to *ForeThought* 5.2.x.

After following the steps above, the information about the LES addresses in your LECS.CFG file would look like this:

### 3.7.2 Delete the LES and BUS

Display the LES information so that you can find the index number of the LES.

Administer the LES and BUS down for Mktg | 0 using the following AMI command:

```
myswitch::configuration lane les> admin 1 down

Delete the co-located LES and BUS.
```

```
myswitch::configuration lane les> delete 1
```

#### Verify the LES has been deleted.

myswitch::configuration lane les> show
No LES information is available.



The LES and BUS in this example were co-located. If yours are not co-located, you need to administer the BUS down using conf lane bus admin <index> down and to delete the BUS using conf lane bus delete <index>. The BUS will be automatically created as a co-located BUS when you use the conf lane les new command.

Do the same for each old LES and BUS, except for the last LES and BUS (in this case, Mktg | 1). All clients will still be using the last LES/BUS temporarily until the DLE peer servers have been established. The last LES and BUS will be changed over at the very end of the process.

# 3.7.3 Upgrade the Switches Running Services

Upgrade each of the switches that is going to be running services to *ForeThought* 5.2.x using the following AMI command:

oper upgrade <remotehost>:<full path to remotefile>

Do not upgrade the switch that is running the last LES and BUS (in this case, Mktg  $\mid$  1). The LECs will still be using the last LES/BUS temporarily until the DLE peer servers have been established.

### 3.7.4 Create the DLE Peer Servers



You must enter the address of each of the DLE peer servers when you are starting DLE; e.g., if you want four peers, then all four must be configured with the addresses of the other three peers, as well as their own LES address, at the time that each LES/BUS is started. If you want to add new peers to the list, you can use the conflane lespeeradd command.

Create a DLE peer server (LES/BUS pair) with the new ELAN name (in this case, Mktg) on each switch that is to become a DLE peer server.

```
myswitch::configuration lane les> new 0x00 Mktg -anycast c5.0005.80.ffe100.0000.f21c.126b.0020481c126b.66 -peers 47.0005.80.ffe100.0000.f21a.24f9.0020481a24f9.00 47.0005.80.ffe100.0000.f21c.10bb.0020481c10bb.90
```

Use the **show** command to verify the information that you entered:

Create each of the other peer servers using the same anycast address and giving them the peer address(es) of each peer in the ELAN. (In this example, there is only one peer.)

```
myswitch::configuration lane les> new 0x90 Mktg -anycast
c5.0005.80.ffe100.0000.f21c.126b.0020481c126b.66 -peers
47.0005.80.ffe100.0000.f21c.10bb.0020481c10bb.90
47.0005.80.ffe100.0000.f21a.24f9.0020481a24f9.00
```

Use the **show** command to verify the information that you entered:

# 3.7.5 Transfer the Updated LECS.CFG File

Transfer the updated LECS.CFG file back to the switch that is running the LECS.

```
myswitch::operation flash> get 169.144.85.195:/tftpboot/lecs.cfg lecs.cfg
```

### 3.7.6 Restart the LECS

Administer the LECS down and back up again. This forces any active clients to re-establish their connection with the LECS and forces the LECS to read and use the new LECS.CFG file.

```
myswitch::configuration lane lecs> admin 1 down
myswitch::configuration lane lecs> admin 1 up
```

Use the **show** command and look at the OperStatus field to verify that the LECS has come up again.

```
myswitch::configuration lane lecs> show

Index AdminStatus OperStatus Selector WKA Database

1 up up 0x00 atm-forum lecs.cfg
```

#### 3.7.7 Recreate the LECs

If you are going to upgrade any or all of the clients to *ForeThought* 5.2.x, you should do so now. Use the following AMI command to upgrade any of the clients that are running on switches:

oper upgrade <remotehost>:<full path to remotefile>



The rest of these instructions apply whether you upgraded the LEC on the switch to *ForeThought* 5.2.x or not. For commands to recreate LECs that are on platforms other than a switch, refer to the appropriate User's Manual.

Use the following command to get the interface name from the IfName field as follows:

```
myswitch::configuration lane lec> show advanced
       Admin
             Oper
Index Status Status Sel Mode MACaddress IfName
                                                            ELAN
    1 up up 0x11 wellknown 0a20481a2c78 el17 Mktg|0
      LECS:0x47.0079.00.000000.0000.0000.0000.00a03e000001.00
       LES :c5000580ffe1000000f21c126b0020481c126b66
       BUS: 47000580ffe1000000f21c126b0020481c126b66
       LEC ID : 13057
                                           Discovered ELAN name : Mktg | 0
       Configure Direct VCC: 0.323
                                          Maximum Frame Size : 1516
       Control Direct VCC: 0.667
                                          Control Distribute VCC: 0.203
       Multicast Send VCC: 0.674
                                          Multicast Forward VCC: 0.205
       Last Error :
```

Configure that LEC down.

```
myswitch::configuration> ip admin el17 down
```

Delete <u>all</u> instances of the LEC (including anything that had been < ELAN name > | n).

```
myswitch::configuration lane lec> delete 1
myswitch::configuration lane lec> delete 2
```

Recreate the LEC. (You now only need one instance per LEC and you will not specify |n| in the name anymore.)

```
myswitch::configuration lane lec> new 0x11 Mktg -ip 192.168.61.25 -mask 255.255.255.0
```

Use the following command to verify that the LEC has joined the ELAN by looking at the OperStatus field as follows:

```
myswitch::configuration> lane lec show advanced
       Admin Oper
Index Status Status Sel Mode
                                       MACaddress
                                                    IfName
                                                              ELAN
    1 up up 0x11 wellknown 0a20481a2c78 el17
                                                              Mktg
       LECS: 0x47.0079.00.000000.0000.0000.0000.00a03e000001.00
       LES :c5000580ffe1000000f21c126b0020481c126b66
       BUS: 47000580ffe1000000f21c126b0020481c126b66
       LEC ID : 13057
                                             Discovered ELAN name : Mktg
       Configure Direct VCC: 0.323
                                           Maximum Frame Size : 1516
       Control Direct VCC: 0.667
                                            Control Distribute VCC: 0.203
       Multicast Send VCC: 0.674
                                            Multicast Forward VCC: 0.205
       Last Error :
```

Repeat all of the steps in this section for each LEC. Refer to the corresponding User's Manual for the appropriate steps for re-creating LECs that are running on hosts, *PowerHubs*, or ES-3810s. At this time, you may also add any new clients to the ELAN.

### 3.7.8 Create the Last DLE Peer

After all of the LECs have been changed over and any new LECs have been added, go back and change the remaining failover LES to a DLE peer. Create the DLE peer as follows:

```
myswitch::configuration lane les> new 0x02 Mktg -anycast
c5.0005.80.ffe100.0000.f21c.126b.0020481c126b.66 -peers
47.0005.80.ffe100.0000.f21a.24f9.0020321a26e5.02
47.0005.80.ffe100.0000.f21a.24f9.0020481a24f9.00
47.0005.80.ffe100.0000.f21c.10bb.0020481c10bb.90
```

# 3.7.9 Add the Last DLE Peer to Each Peer List

Now add the new peer to the list of each of the existing peers using the following command on each of those peers:

```
myswitch::configuration lane les> new -peeradd 1 47.0005.80.ffe100.0000.f21a.24f9.0020321a26e5.02
```

# 3.7.10 Update the LECS.CFG File

After you add the last DLE peer, go back and delete the old remaining failover ELAN information from the LECS.CFG file. Transfer the file to a workstation for editing.

```
myswitch::operation> flash put lecs.cfg 169.144.85.195:/tftpboot/lecs.cfg
Transferred 2323 bytes of fs:/lecs.cfg
```

On that workstation, use a text editor to delete all lines referring to any of the |n| failover ELANs (in this case, Mktg |0| and Mktg |1|). If the LES address portion of your LECS.CFG file looked like this before editing:

Match.Ordering: Mktg, Mktg | 0, Mktg | 1

#### It will look like this after editing:

Match.Ordering: Mktg

### 3.7.11 Transfer the Final LECS.CFG File

Transfer the final LECS.CFG file back to the switch that is running the LECS.

myswitch::operation flash> get 169.144.85.195:/tftpboot/lecs.cfg lecs.cfg

### 3.7.12 Restart the LECS

Administer the LECS down and back up again. This forces the active LECs to re-establish their connection with the LECS and forces the LECS to read and use the new LECS.CFG file.

```
myswitch::configuration lane lecs> admin 1 down
myswitch::configuration lane lecs> admin 1 up
```

Use the **show** command and look at the OperStatus field to verify that the LECS has come up again:

The transfer of your ELAN to DLE services is now complete.

# 3.8 Upgrading an ELAN without Using DLE

This section describes how to upgrade your ELAN if you had previously configured LANE, you want to leave the clients running *ForeThought 4.1.x*, and you want to upgrade all the equipment that is running services to *ForeThought 5.2.x* **without** using DLE.

The following basic steps are involved in upgrading your ELAN *ForeThought* 5.2.x **without** using DLE. Each of these steps is described in detail in the following sections. It is recommended that you read the entire section before attempting the upgrade.

- 1. If the LES and BUS are not co-located, they must be deleted and re-created as co-located because *ForeThought* 5.2.x on the switch only supports co-located services. If the LES and BUS are already co-located, skip to step 2.
- 2. Upgrade the switches that are running services.
- 3. Re-create the LES and BUS if they were not co-located. If the LES and BUS were already co-located, skip to step 4.
- 4. Administer up the switches that are running services.



If you used the VLAN Manager to create your ELAN, you should use it to upgrade the ELAN. Refer to the *ForeView VLAN Manager User's Manual* for the appropriate steps.



These steps describe commands for services running on a switch. Refer to the corresponding User's Manual for the appropriate steps for commands for any services that are running on hosts, *PowerHubs*, or ES-3810s.

# 3.8.1 Deleting the Non Co-located Services

There are several steps involved in changing the non co-located services.

#### 3.8.1.1 Administer Down the Services

Administer down the LECS.

myswitch::configuration lane> lecs admin <LECS index> down
Administer down the LES and BUS.

```
myswitch::configuration lane> les admin <LES index> down
myswitch::configuration lane> bus admin <BUS index> down
```

#### 3.8.1.2 Delete the Non Co-located LES and BUS

Delete the LES and BUS.

```
myswitch::configuration lane> les delete <LES index>
myswitch::configuration lane> bus delete <BUS index>
```

#### 3.8.1.2.1 Edit the LECS.CFG File



Before you edit the LECS.CFG file, you may wish to back it up to a host using the oper flash put command. If you are using TFTP as the transfer protocol (this is the default), the remote host to which the FLASH file will be sent must be running the TFTP server code and must have an empty file in the /tftpboot directory on the remote host to receive the FLASH file. See Chapter 3 in the ATM Management Interface (AMI) Manual for more information. If you are using FTP as the transfer protocol, you only need to enter the command shown below to perform the transfer.

Transfer the LECS.CFG file that is currently being used by the LECS to a workstation (other than the one to which you backed up the file) so you can edit the file.

```
myswitch::operation> flash put lecs.cfg 169.144.85.195:/tftpboot/lecs.cfg
Transferred 2323 bytes of fs:/lecs.cfg
```

On the workstation that has the LECS.CFG file, use a text editor to delete any lines that refer to any non-colocated BUSs.

After you have made the changes, transfer the LECS.CFG file back to the LECS.

```
myswitch::operation flash> get 169.144.85.195:/tftpboot/lecs.cfg lecs.cfg
```

# 3.8.2 Upgrade the Switches Running Services

Upgrade the switch(es) running each of the services to *ForeThought* 5.2.x using the following AMI command:

```
oper upgrade <remotehost>:<full path to remotefile>
```

# 3.8.3 Recreate the LES and BUS Together

If your LES and BUS are already co-located, skip to step 4. If your LES and BUS were not co-located, re-create each LES and BUS using the following single AMI command:

```
myswitch::configuration lane> les new <LES Selector Byte (HEX)> <LES name>\
    [-bus <BUS Selector Byte (HEX)>]\
    [-type (ethernet | token-ring)]\
    [-mtu (1516 | 1580 | 4544 | 9234 | 18190)]\
    [-secure wka | <LECS ATM Address>]\
    [-registertlvs (enable | disable)]\
    [-anycast <LES Anycast ATM Address>]\
    [-peers <atm-addr> ...]
```

You need to specify the LES selector byte (the BUS uses the same selector byte by default) and you need to give the same LES name that you used before. You may optionally specify any of the parameters, except for the anycast address and the peer addresses since you are not using DLE.

# 3.8.4 Administer the Services Up

Administer up the LES/BUS pair and the LECS on the switch(es) running each of the services.

```
myswitch::configuration lane> les admin <LES index> up
myswitch::configuration lane> lecs admin <LECS index> up
```

The transfer of your ELAN to ForeThought 5.2.x is now complete.

# Configuring an Emulated LAN

# CHAPTER 4 MPOA

This chapter provides an overview of LAN Emulation (LANE) and Multi-Protocol Over ATM (MPOA) as implemented in FORE Systems' *ForeThought* software.

## 4.1 Overview of LANE/MPOA

FORE System's *ForeThought* software is compliant with the ATM Forum's *LAN Emulation Over ATM Version 1.0* specification. LANE allows higher level protocols and LAN applications to interoperate, without modifications, with an ATM network.

The LANE components, running on the ATM network, interact to emulate an Ethernet or Token Ring LAN. This emulated Ethernet or Token Ring LAN is called an *emulated LAN* (*ELAN*). The ELAN components resolve MAC addresses to ATM addresses, replace the connectionless operation of legacy LANs with point-to-point connections, and provide broadcast and multicast services. The ELAN consists of a LANE/MPOA Client (LEC/MPC) running on each host in the ELAN, and the following LANE Services:

- the LAN Emulation Server (LES)
- the Broadcast and Unknown Server (BUS)
- the LAN Emulation Configuration Server (LECS)

The LANE services may operate on a FORE Systems ATM switch, *PowerHub* 7000, or Solaris workstation. *ForeThought* also provides support for Distributed LAN Emulation (DLE) which provides load-sharing and improved fault-tolerance within an ELAN.

## 4.2 LANE Primer

LANE is the foundation on which MPOA is built. Therefore, before presenting an explanation of MPOA, an understanding of LANE components and their operation in an ELAN is needed.

## 4.2.1 LANE Components

An ELAN includes the following components:

LANE/MPOA Client (LEC/MPC)

The LEC/MPC can wear two different "hats." When wearing its LEC "hat," it simply communicates with other ELAN components (the LES and BUS) to resolve MAC addresses into ATM addresses. When it puts on its MPC "hat," the additional function of the LEC/MPC in an MPOA-aware network is to source and sink internetwork shortcuts.

LAN Emulation Configuration Server (LECS)

Runs on a Solaris workstation or a FORE Systems ATM switch. Maintains information about all ELANs within the administrative domain. When the LEC/MPC successfully communicates with the LECS, the LECS provides a list of ELANs which the LEC/MPC can join. The LECS may be configured with various MPOA parameters. LEC/MPCs that connect to LANE/MPOA services through an MPOA-aware LECS are configured with these centrally-supplied MPOA parameters. LEC/MPCs that connect through an LECS that does not contain MPOA parameters still perform flow analysis and attempt inter-ELAN shortcuts according to their user-editable or factory-default settings.

#### **LAN Emulation Server (LES)**

Runs on a *PowerHub* 7000, a *ForeRunner* ATM switch, or a Solaris workstation. Maintains information about the LEC/MPCs within a single ELAN and performs address resolution. The LES can be configured to support or disable MPOA operation in an ELAN. The LES accepts MPOA parameters from registering LEC/MPCs and MPSs, and also distributes MPOA parameters to LEC/MPCs in response to queries. (This is the mechanism used by LEC/MPCs to determine whether routers in the ELAN are MPOA-aware).

# Broadcast and Unknown Server (BUS)

Runs on a *PowerHub* 7000, a *ForeRunner* ATM switch, or a Solaris workstation. Provides services within a single ELAN allowing broadcasts, multicasts, and unknown unicasts. The BUS is MPOA-ignorant.

## 4.2.2 An Example LANE Configuration

Figure 4.1 shows an example configuration of a single ELAN in a FORE network. The ELAN includes:

- PC Workstations, each running a LEC/MPC. Each has a ForeRunner ATM adapter, the ForeRunner driver for the adapter, and one or more ForeRunner ELAN drivers installed.
- Two ForeRunner ASX-200BX switches running LESs, BUSs, and LECs. Each switch
  is also running an LECS. The LES/BUS pairs are configured as peers under Distributed LAN Emulation. The peer configuration allows the LECs associated with
  a particular LES/BUS automatically to reconnect to the remaining functional peer
  if their "home" LES/BUS fails.
- A PowerHub 7000 running a LEC/MPC, and providing access to non-ATM networks.

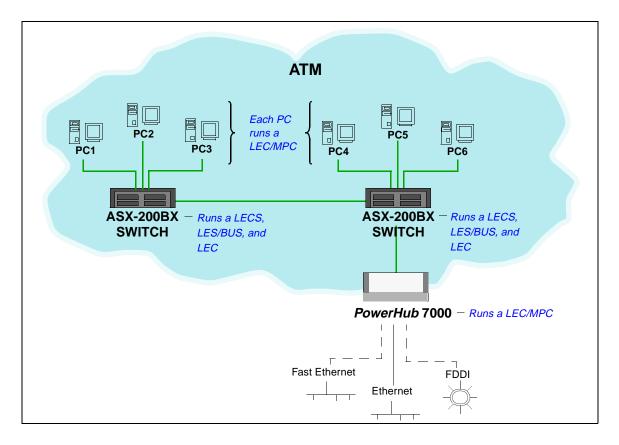


Figure 4.1 - An Example of an ELAN

#### 4.2.2.1 The Initialization Process

Each LEC/MPC goes through the following process when it starts up:

- 1. The LEC/MPC obtains its own ATM address via address registration. Optionally, you can manually specify the ATM address.
- 2. The LEC/MPC establishes a connection to a LECS using an address obtained via ILMI, a well-known address, or PVC (0,17). Optionally, you can manually specify the LECS address.

MPOA

- 3. The LEC/MPC requests the information needed to join a specified ELAN or the default ELAN. The LECS has information about available ELANs, what ELANs each LEC/MPC can join, and which ELAN the LEC/MPC should attempt to join first.
  - If a LECS is not available, or if you choose not to use it, you can manually specify the information required to join a specific ELAN.
- 4. The LEC/MPC contacts the LES associated with the ELAN it wants to join and registers its MAC-ATM address pair. It also contacts the BUS associated with the ELAN. At this point, the LEC/MPC and the LES have the information required to allow this host to communicate with other hosts on the ELAN as if it were an Ethernet (or Token Ring) network. Refer to the following section for a description of how the LEC/MPC connects to other hosts on the ELAN.

#### 4.2.2.2 The Connection Process

To send packets to another host on the ELAN:

- 1. The LEC/MPC calls the LES to map the MAC destination address into an ATM address. (The LES maintains a mapping table of the address of all LEC/MPCs on the ELAN.)
- 2. If the LES finds an entry in its table for the destination MAC address, it returns the destination ATM address to the LEC/MPC.
- 3. The LEC/MPC then opens up a point-to-point ATM connection to the destination host to send the packet.

#### 4.2.2.3 Multicast and Broadcast Packets

The LEC/MPC sends outgoing multicast and broadcast packets to the BUS which uses a point-to-multipoint connection to send the packets to multiple ATM addresses in the ELAN.

## 4.2.2.4 Accessing Fast Ethernet and FDDI Networks

Note that the diagram in Figure 4.1 shows dotted lines from the *PowerHub* 7000 to the Fast Ethernet and FDDI networks. This shows how the *PowerHub* can provide access to non-ATM networks.

#### 4.2.2.5 Multiple ELANs

It is possible to set up more than one ELAN in a FORE network. For each new ELAN, you must configure another LES/BUS instance for that LAN. On the access devices, bridge groups must be used to associate physical ports with ELANs on the ATM side. An end station in the ELAN with a *ForeRunner* adapter can connect to up to 16 ELANs simultaneously.

#### 4.2.2.6 Distributed LAN Emulation

To provide greater resilience, support larger ELANs, and support separated clusters of users in an ELAN, *ForeThought* software provides Distributed LAN Emulation (DLE). DLE allows the LES/BUS functions to be distributed among multiple interconnected LES/BUS instances called *peers*. In the example ELAN shown in Figure 4.1, the two LES/BUS pairs running in the switches function as peers in the same ELAN. The LEC/MPCs are distributed such that they are not all connected to the same server. With this arrangement, should one of the peer servers fail, the clients connected to the remaining server continue to maintain connectivity; while the clients that were connected to the failed server automatically re-establish connectivity to the ELAN within 60 seconds.

#### 4.2.2.7 Automatic ELAN Selection

To simplify configuration of the ELAN, a host is allowed to join an ELAN without specifying an ELAN name. If the LECS has been configured to provide the required information, and you do not manually specify an ELAN name to join when you configure the host's ELAN driver, the host initially attempts to join the ELAN specified by the LECS. The host successfully joins the ELAN if the LECS is available, the proper LES address for the ELAN has been specified in the LECS, and the LES and BUS are available.

#### 4.2.2.8 Intelligent BUS

This feature reduces broadcast traffic by using the MAC address information in the LES. When an intelligent BUS receives a unicast frame, the BUS first checks the LES's mapping table to see if the MAC address is registered there. If it is, the BUS forwards the frame directly to the destination, instead of broadcasting.

# 4.3 An Introduction to Multi-Protocol Over ATM

MPOA builds upon the foundation of LANE.

#### 4.3.1 LANE Without MPOA

ATM networks co-exist with and support network applications which may not be ATM-aware. Consequently, ATM protocols are needed to monitor legacy network protocol (IP, IPX, Appletalk, etc.) packets and perform translation into ATM cells and circuits. This monitoring and translation can be performed in one of the following ways:

- in a host protocol stack after packet construction and before packet transmission
- in a LAN-to-ATM edge device as packets move through the network

LANE is one example of such a protocol. It resolves datalink layer addresses into ATM addresses and establishes circuits to the destination addresses. Network addresses *within* a subnet can be learned using LANE's broadcast support.

However, LANE relies on physical routers to deliver packets *across* subnets (see Figure 4.2). Because routers *must* examine – and modify – *every* packet, ATM cells *must* be reassembled into packets, modified, and re-segmented at *every* router hop. This process imposes significant transmission delays between the source and destination of the network traffic.

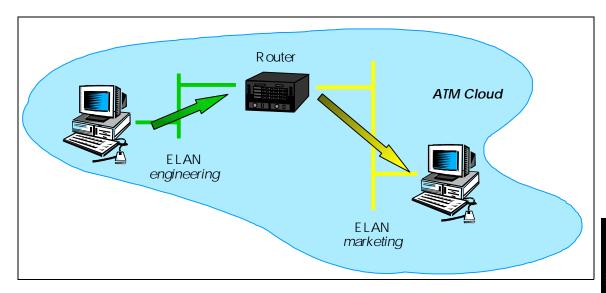


Figure 4.2 - LANE without MPOA

In addition to LANE, protocols such as IP can operate over an ATM network via the IETF Internetworking Over NBMA Networks (ION) Working Group's *Next Hop Resolution Protocol* (NHRP). NHRP allows the ATM network to be divided into Logical IP Subnets (LISs). Using NHRP, routers are still required to interconnect these subnets; but NHRP permits intermediate routers to be bypassed on the *data* path. NHRP allows entities called Next Hop Clients (NHCs) to send queries between different subnets. These queries are propagated using Next Hop Servers (NHSs) via paths found using standard routing protocols. Consequently, NHRP enables the establishment of VCC data paths across subnet boundaries *without requiring physical routers in the data path*.

## 4.3.2 Why MPOA?

The ATM Forum developed the Multi-Protocol over ATM (MPOA) specification to address the limitations of LAN Emulation. MPOA extends ATM support of legacy networks into the network layer. The main objective of MPOA is the efficient transfer of unicast data between subnet(s).

MPOA introduces LANE/MPOA Clients (LEC/MPCs) and MPOA Servers (MPSs) and defines the protocols that are required for LEC/MPCs and MPSs to communicate. LEC/MPCs issue queries for ATM addresses, and receive replies from the MPS using these protocols. MPOA also maintains interoperability with the existing infrastructure of routers. MPOA Servers reside in routers that run standard Internetwork Layer routing protocols such as OSPF, thus providing integration with existing networks.

ForeThought software implements MPOA shortcuts for IP traffic. It does this by adding capabilities to LANE, not by replacing LANE. LANE/MPOA client drivers are extended LANE drivers. When handling traffic within the same ELAN and subnet, they function like LECs. However, when handling traffic that crosses subnets, LEC/MPCs initially work with MPOA servers (MPSs) to use MPS-established hop-by-hop circuits. Then, for traffic flows that exceed configurable limits, shortcut circuits are built that allow the traffic to traverse the route without the necessity of the router(s):

- · reassemble packets from ATM cells
- modify the packets
- and then re-segment the packets for transmission to the next hop

Consequently, traffic flowing through a shortcut VCC moves at essentially wire speed from source to destination (see Figure 4.3).

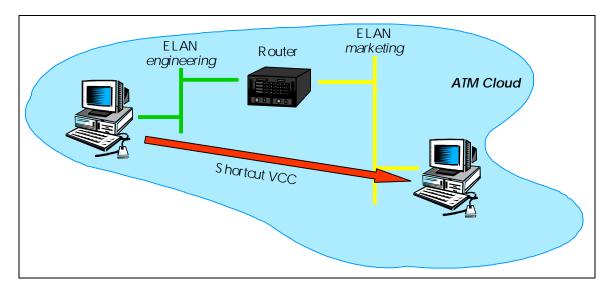


Figure 4.3 - LANE with MPOA

## 4.3.3 MPOA Components

MPOA requires LANE services for both ELAN traffic handling and MPOA configuration.

LANE/MPOA Client (LEC/MPC)

The LEC/MPC can wear two different "hats." When wearing its LEC "hat," it simply communicates with other ELAN components (the LES and BUS) to resolve MAC addresses into ATM addresses. When it puts on its MPC "hat," the additional function of the LEC/MPC in an MPOA-aware network is to source and sink internetwork shortcut circuits. A LEC/MPC that is the source of a shortcut is known as an *ingress* LEC/MPC. A LEC/MPC that is the sink of a shortcut is known as an *egress* LEC/MPC. The LEC/MPC includes an NHRP Client (NHC).

An ingress LEC/MPC monitors traffic flow that is being forwarded over an ELAN to a router that contains an MPS. When the ingress LEC/MPC recognizes a flow rate (configurable) that could benefit from a shortcut (and thus bypass the routed path), it requests a shortcut to the destination. If a shortcut is available, the ingress LEC/MPC sets up a shortcut VCC, and forwards traffic for the destination over the shortcut.

An egress LEC/MPC receives internetwork traffic from other LEC/MPCs to be forwarded to its local interfaces/users. For traffic received over a shortcut, the egress LEC/MPC adds the appropriate encapsulation and forwards them via a LAN interface (that may be a bridge port, an internal host stack, etc.).

MPOA Server (MPS)

An MPS includes an NHRP Server (NHS) and is the logical component of a router that provides internetwork layer forwarding information to LEC/MPCs. The MPS answers MPOA queries from ingress LEC/MPCs and provides encapsulation information to egress LEC/MPCs.

The MPS also converts between MPOA requests and replies, and NHRP requests and replies on behalf of LEC/MPCs.

## 4.3.4 MPOA Example

The following are the basic requirements for establishing a shortcut across an MPOA-enabled network:

- There must be LEC/MPCs at each end of the network between which a shortcut is desired.
- The local router interface at each end must be running an MPS.
- A Next Hop Resolution Protocol (NHRP) path must exist between MPSs.

The following example illustrates a typical ATM network that allows MPOA shortcuts to be employed.

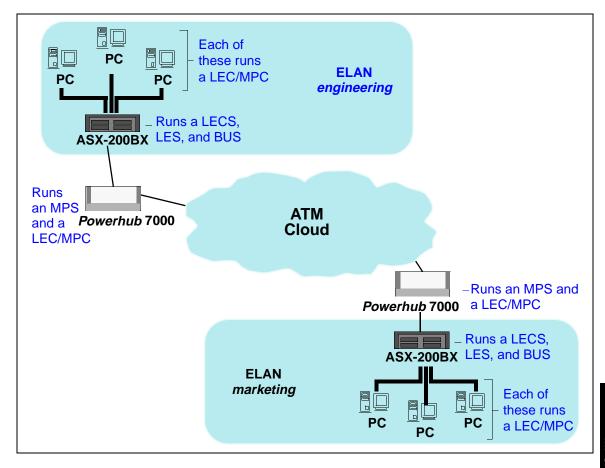


Figure 4.4 - MPOA Example Network

#### 4.3.4.1 MPS Configuration

The network administrator must configure each MPS with the site-specific IP address matching the gateway address being used by LEC/MPCs in its ELAN.

The MPS on each *PowerHub* 7000 is configured as follows:

- 1. For each LANE/MPOA virtual port, specify an ELAN name. The LECS configuration must also be updated to allow the MPS to join these ELANs.
- 2. For each LANE/MPOA virtual port, specify an IP address.
- 3. Enable LANE/MPOA support.
- 4. Enable routing.
- 5. Save the configuration and reboot the MPS if necessary to make the changes effective.

Router table information need not be configured. The MPS will instead gather this information through routing protocol exchanges with other MPS's.

#### 4.3.4.2 Initialization

When its host boots, *each* LEC/MPC automatically goes through the following sequence to establish a connection to the MPS.

- 1. The LEC/MPC registers via ILMI with the switch to which it is attached.
- The LEC/MPC connects to an LECS to which it sends its own ATM address and
  the name of the ELAN it wishes to join (the ELAN name is an empty string unless
  the LEC/MPC has been site-configured with an ELAN name). The LEC/MPC also
  supplies a LANE 1.0 compliant parameter identifying itself as an MPOA-aware
  client.
- 3. Next, the LEC/MPC receives the following from the LECS:
  - the name of the ELAN to which it is assigned
  - the ATM address of the LES for the ELAN it is joining
  - the parameters containing the flow detection and shortcut establishment policies it is to use
- 4. The LEC/MPC then connects to its assigned LES, and provides the LES with a parameter identifying itself as MPOA-aware.
- 5. Finally, the LEC/MPC connects to the ELAN's BUS.

Once the LANE/MPOA connections are established, third-party network-layer protocol drivers on the host can establish network-layer connectivity. The methods these upper-layer drivers use to determine host IP addresses, default gateway, and backup gateway addresses vary depending on the third-party product. For example, the LANE/MPOA driver itself permits these drivers to use BOOTP or DHCP to obtain IP configuration information.

#### 4.3.4.3 Flow Analysis

On a LEC/MPC's host or edge device, IP packets with destinations *within* the host's subnet are sent using LANE 1.0 methods; i.e., the client puts on its LEC "hat" and works directly with its ELAN's services to connect with local destinations. Packets destined for *remote* subnets cause the LEC/MPC to put on its MPC "hat." This client is then referred to as an "ingress LEC/MPC."

Ingress LEC/MPCs associate destination IP addresses with shortcut circuits. Ingress LEC/MPCs use configurable parameters called *flow descriptors* to determine *whether* and *when* to trigger creation of shortcut circuits. The ingress LEC/MPC also monitors the most recent use of a shortcut circuit to determine when to tear down the shortcut. Specifically, when an ingress LEC/MPC sends a packet:

- 1. If a shortcut circuit *already exists* to the IP destination, the LEC/MPC sends the packet over this circuit.
- 2. If no shortcut circuit exists, the LEC/MPC determines *whether* shortcuts to this IP address are allowed. If shortcuts to the destination IP address are *not* allowed, the LEC/MPC sends the packet to the gateway router.
- 3. If no shortcut circuits exist, *and* shortcuts to the IP address *are* allowed, the LEC/MPC determines if the packet traffic flow exceeds the shortcut enable trigger value (set by the flow descriptors) for the destination IP address's flow. If the flow exceeds the trigger value, the LEC/MPC tries to establish a shortcut circuit to the destination LEC/MPC (called the egress LEC/MPC). If the flow does not exceed the trigger value, the ingress LEC/MPC simply sends the packet traffic to the gateway router.

## 4.3.4.4 Making a Shortcut

When the ingress LEC/MPC determines that the packet traffic flow exceeds the shortcutenable trigger value, the ingress LEC/MPC tries to establish a shortcut circuit to the egress LEC/MPC. The following describes how a shortcut is set up:

- The ingress LEC/MPC initiates the shortcut creation process by sending an MPOA
  resolution request to the MPS it uses as a gateway router (this MPS is called the
  ingress MPS). The MPS converts the request to a next hop resolution protocol
  (NHRP) request. This NHRP request includes the destination's IP address and asks
  for the corresponding ATM destination address.
- 2. This request is passed along hop-by-hop until it reaches the final MPS (called the *egress* MPS) on the route to the destination IP address.
- 3. The egress MPS sends a cache imposition request to the egress LEC/MPC. The egress LEC/MPC sends a cache imposition reply, which is converted to a NHRP response by the egress MPS. The ingress MPS converts the NHRP response to an MPOA resolution response and transmits it to the ingress LEC/MPC.

4. When the ingress LEC/MPC receives the NHRP response containing the destination's ATM address, it first checks if a shortcut circuit to that ATM address already exists. If a shortcut circuit to that address already exists, it sends the packets via the existing shortcut circuit. If no shortcut circuit exists it opens a new shortcut circuit and begins sending packets over it to the destination.

#### 4.3.4.5 Shortcut Teardown

Application programs and networking protocol stacks are MPOA-ignorant and, therefore, do not tear down shortcut circuits when the shortcut is no longer needed. Therefore, the MPOA layer itself tears down seldom-used shortcuts to avoid circuit exhaustion in the client and network. When a shortcut is idle for a period exceeding a set limit, the shortcut is torn down.

# CHAPTER 5

# ForeThought PNNI

PNNI (Private Network Node Interface or Private Network-to-Network Interface) is an ATM Forum approved protocol which defines interoperability between private ATM switches. PNNI defines both the routing and signalling standards for inter-switch interoperability.

This chapter provides an overview of FORE Systems' pre-standard version of PNNI, *ForeThought* PNNI (FT-PNNI), and its use in a multiple-switch network. FT-PNNI is a scalable routing and signalling protocol used in networks containing multiple *ForeRunner* switches. FT-PNNI simplifies large network topologies by organizing the nodes (switches) in that network into smaller groups.

It is this reorganization of the network topology that makes FT-PNNI's simplified routing possible. By segmenting a large network into smaller peer groups of nodes, FT-PNNI reduces the amount of network topology information that those very nodes must maintain.

# 5.1 FT-PNNI Routing

The FT-PNNI routing protocol serves to distribute topology and address reachability information between switches and groups of switches in a network. This topology and addressing information is used by switches to compute paths through the network. The functions of the FT-PNNI routing protocol include the following:

- · Hello Protocol
- Topology Database Exchange
- Flooding
- Hierarchical Routing

#### 5.1.1 Hello Protocol

FT-PNNI hello packets are exchanged between neighboring nodes. Each switch (node) transmits a hello indication on each of its FT-PNNI routing channels at regular intervals. The time between these hello indications is called the Hello Indication Interval. When a switch receives a hello indication from one of its neighbors, it stores the logical link (loglink) from that neighbor to itself in the topology database.

Based on the hello indications, loglinks are refreshed periodically. Since loglinks discovered as a result of hello indications are unidirectional, each switch stores unidirectional loglinks with its immediate neighbor as the source and itself as the destination.

## 5.1.2 Topology Database Exchange

Each switch sends to each of its neighbors a group of loglinks from its topology database at regular intervals, called the NSAP map indication interval. This exchange of information between neighboring switches ensures that the topology databases of the switches stays synchronized.

# 5.1.3 Flooding

Flooding is a reliable hop-by-hop propagation of loglinks throughout a peer group. Whenever a new loglink is discovered by a switch, the switch immediately broadcasts the existence of this link to all of its neighbors. The neighboring switches then broadcast the existence of the same link to all of their neighbors. Very quickly, the existence of the new loglink is flooded throughout all of the switches in the network.

Similarly, when a link goes down, or when a *significant* change is seen in the metrics of a loglink, the latest state of the loglink is propagated immediately throughout the network.

# 5.1.4 Hierarchical Routing

FT-PNNI operates in a hierarchical topology. The structure of the hierarchy is defined by the peer group ID in a routing domain. Address assignment in FT-PNNI, therefore, corresponds to this hierarchical topology, providing increased scalability.

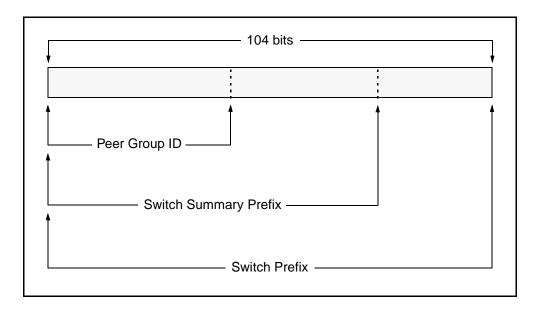


Figure 5.1 - Example of a 13-byte Switch Prefix

## 5.1.4.1 Hierarchical Addressing

FT-PNNI uses private ATM address prefixes (NSAP prefixes) as node identifiers. FT-PNNI does not distinguish between node identifiers and reachability information. Thus, the IDs of nodes in the FT-PNNI addressing map are NSAP prefixes. In the default case at the lowest level peer group, the switches have a 13-byte prefix as their node ID and end systems (hosts) have a 19-byte prefix as their node ID.

#### 5.1.4.1.1 Switch Prefix

Each switch in a FT-PNNI network is configured with a 13-byte prefix called the switch prefix. Hosts that are attached to the switch are presented with this prefix during ILMI address registration. In this way, end systems are configured with a private ATM address that includes the 13-byte switch prefix.

#### 5.1.4.1.2 Switch Summary Prefix

Each switch is configured with a switch mask (*swmask*) which gives the length of the switch summary prefix within the switch prefix. The *swmask* gives the number of most significant bits of the switch prefix that constitute the switch summary prefix. Since all end system addresses attached to a switch have the same switch summary prefix, their reachability information can be summarized by this prefix (i.e., by the switch summary prefix).

#### 5.1.4.1.3 Peer Group ID

Each switch is configured with a peer group mask (*pgmask*) which gives the length of the peer group ID within the switch summary prefix. The *pgmask* gives the number of most significant bits of the switch summary prefix that constitute the peer group ID.

Each peer group has a peer group ID that uniquely identifies it from every other peer group. Every node (switch or end system) in a particular peer group shares that same unique peer group ID, thereby indicating membership to that peer group.

A simple example of summarizing by peer group ID can be seen in Figure 5.3, where every switch (and end system, although not shown) in peer group A is identified starting with "A."

## 5.1.4.2 Path Computation

Path computation is performed on demand whenever FT-PNNI signalling requests a path to a given destination. The Bellman-Ford Shortest Path algorithm is used to compute the shortest path tree of all nodes in the topology with the local node as the source. The administrative weight metric in the loglinks is used as the minimizing criterion in computing the shortest path route. In the case of finding multiple equal-cost paths to a given destination, available cell rate is used to break ties.

# 5.2 The Physical Network

In an ATM network, data is sent and received over virtual circuits, or circuits that only exist when needed. This communication over these virtual circuits is made possible by signalling that occurs between the switches in the network.

In a network of *ForeRunner* switches, any new addition to the topology is recognized immediately by all nodes (switches) having a direct connection to the new node. Then each switch that has recognized the new switch sends a message to each of its direct connections, and so on. Eventually, and within a very brief period of time, every switch in the network is aware of the new addition and the links by which that new addition can be reached. This topology is stored by each switch in its local topology database.

In a small, local area network (LAN), the topology is relatively simple, meaning that the switches in the LAN have a relatively small topology database to maintain. As the LAN grows, however, and more switches are added, that same database can grow to be very large in a short period of time.

As this topology database grows, the amount of information a switch must look up when searching for an address also grows. In the end, this can result in delayed connection set-up, congestion in the network, and even lost data.

Figure 5.2 depicts a typical ATM network, containing 21 *ForeRunner* switches (). The hierarchy of this network is flat, meaning that each switch must be aware of all the other switches in the network, as well as all the possible routes to those switches. As more switches are added to this network, the hierarchy will become more complex and the switches will have to contend with larger topology databases.

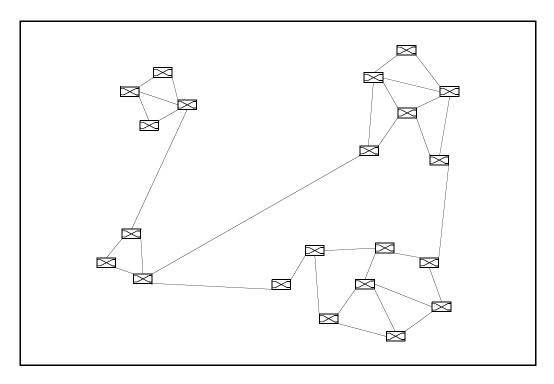


Figure 5.2 - Private ATM Network with 21 Switches and 34 Bidirectional Links

It is in these large, single-level networks that FT-PNNI is most useful, because it lets you simplify large network topologies by creating a two-level hierarchy. In this hierarchy, clusters of contiguous switches are grouped together and they are collectively summarized by a single, logical node.

Figure 5.3 shows the same network as in Figure 5.2 after being organized into peer groups, now having a two-level hierarchy. The subsections that follow explain the organization of these peer groups, how they simplify the overall network topology, and how they change the logical view of the network.

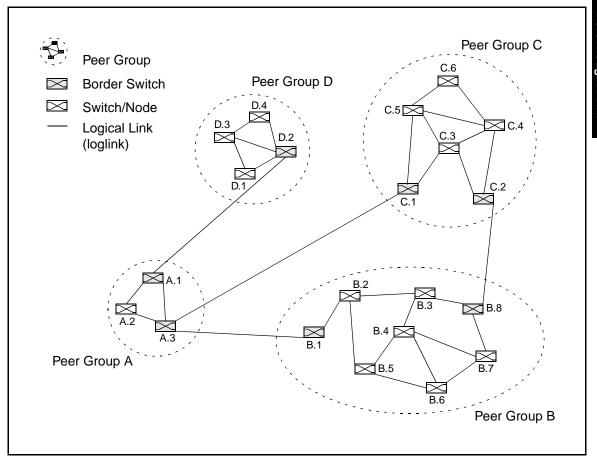


Figure 5.3 - Example of FT-PNNI Hierarchy Showing Lowest-Level Peer Groups

## 5.2.1 Peer Groups

The FT-PNNI hierarchy begins with a network of switches, organized into peer groups. A peer group is a collection of interconnected switches that are organized into a group. Peer group organization can be determined by a network administrator, but switches that are located close to one another are usually made into a peer group.

The network shown in Figure 5.3 is organized into four lowest-level peer groups: A, B, C, and D. The switches within a certain peer group are numbered according to that particular group. For example, the switches in peer group A are identified as A.1, A.2, and A.3.

Peer groups have a peer group identifier (ID), assigned at configuration time and exchanged in hello indication messages. Neighboring switches can determine if they belong to the same peer group by comparing these peer group IDs.

Switches in a peer group are aware of the topology of their own peer group and the existence of all other peer groups. They recognize the links between their peer group and others, but they are not aware of internal topology information of other peer groups. Instead, the topologies of other peer groups are summarized as a single, reachable location, known as a peer group summary node (PGSN).

# 5.2.2 Peer Group Topology

To maintain an accurate and updated view of its relative location and status, a switch periodically sends a hello indication message to every other switch with which it has a direct connection. These hello indications contain the switch prefix, peer group membership information, and link metrics (attributes) for the physical link between the two switches.

Through this regular exchange of messages, each switch learns which switches are its immediate neighbors, to what peer groups they belong, and whether or not the link between itself and its neighbors is valid.

## 5.2.3 Border Switches

A border switch is any switch that has at least one link to a switch in another peer group. Border switches play an important role in FT-PNNI because they are responsible for summarizing reachability information for their respective peer groups, appropriately filtering the flow of topology database information across peer group boundaries, and building the lowest level source route for call setups entering the peer group.

## 5.2.4 Peer Group Summary Node (PGSN)

A PGSN is a virtual (logical) or imaginary node that summarizes a peer group's reachability information. The PGSN has the peer group ID of its peer group as its switch summary prefix. Each border switch in the peer group advertises a logical link (loglink) to the PGSN. The PGSN is a logical representation of the switches contained in a peer group.

## 5.2.5 Backbone Topology

Loglinks between border switches and loglinks from border switches to PGSNs are called backbone links and considered part of the backbone topology. Information regarding these backbone links is propagated across peer group boundaries during database exchange and flooding.

## 5.2.6 Single Switch Perspective

The main reason for grouping switches in large networks is to simplify each individual switch's view of the topology. For example, each switch in peer group A is aware of every other switch in peer group A, the border switches in the rest of the network, the links between them, and the backbone topology. Switches in peer group A are not aware, however, of the internal topology of other peer groups. Instead, individual switches see a PGSN (see Figure 5.4).

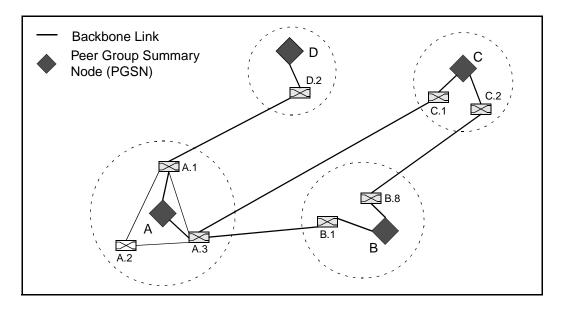


Figure 5.4 - View of the Network from Switches in Peer Group A

# CHAPTER 6

# **ATM Forum PNNI**

PNNI (Private Network Node Interface or Private Network-to-Network Interface) is a scalable protocol that defines both the routing and signalling standards for interoperability between private ATM switches. PNNI simplifies large network topologies by organizing the nodes in a network into smaller groups. Through this re-organization, PNNI reduces the amount of topology information that those nodes must maintain.

This chapter provides an overview of FORE Systems' implementation of ATM Forum PNNI (hereafter referred to as PNNI), and its use in a multiple-switch network. To understand PNNI, it is important to first understand how PNNI's routing and signalling protocols work.

# 6.1 PNNI Routing Protocol

The PNNI routing protocol enables the distribution of topology and address reachability information between switches and groups of switches. *Reachability information* is the data that explains how to contact a single ATM address or a group of ATM addresses that are summarized by a single prefix. This information is used to compute routes across the network. The key elements of the PNNI routing protocol are as follows:

- Hello Protocol
- Database Exchange Protocol
- Flooding Protocol
- Path Computation
- Hierarchical Routing

# 6.1.1 Hello Protocol

The Hello protocol is used to discover and verify the identity of neighbor nodes and to determine the status of the logical links to those nodes. Logical links between nodes in the same peer group are called *horizontal links*. Hello packets are exchanged at regular intervals over links connecting neighboring nodes. Initially, these packets are used to exchange identities between the nodes and to determine if they belong to the same peer group. After that, Hello packets are used as keep-alive messages for the links over which the packets are exchanged. After a configured number of hello intervals have passed without receiving a hello message, the link is declared to be down.

## 6.1.2 Database Exchange Protocol

When neighboring nodes determine that they are peers; i.e., that they belong to the same peer group and they are running compatible versions of the PNNI protocol, they begin a database exchange process. At regular intervals, each node tells the other what information it has in its topology database and requests similar information from the other neighboring nodes. This exchange of information ensures that the topology databases of all switches in the network stay synchronized.

## 6.1.3 Flooding Protocol

Topology information is organized into smallest units called PNNI Topology State Elements (*PTSEs*). PNNI routing uses a flooding mechanism to propagate these PTSEs throughout the network in a hop-by-hop fashion. PTSEs are encapsulated into PNNI Topology State Packets (*PTSPs*) for transmission. PTSPs always contain the node identifier (ID) and the peer group ID of the switch that originated all of the PTSEs.

When a PTSP is received, its component PTSEs are examined. Each PTSE is acknowledged by sending a PTSE acknowledgment packet (containing possibly a bundle of PTSE acknowledgments) to the sending peer. If the PTSE is more recent than the copy that the receiving node has in its topology database, the database copy is updated and is re-flooded to all neighboring peers except the one from which the PTSE was received. A PTSE sent to a neighboring peer is periodically retransmitted until it is acknowledged.

In this fashion, when a link is created or goes down, this information is propagated immediately throughout the network. Similarly, when a *significant* change is seen in the metrics of a link, this state change is flooded throughout the network.

# 6.1.4 Path Computation

Path computation is done on the source (or originating) switch. All of the network topology information obtained and updated by the higher layers of the protocol is used to determine a path through the network which is efficient and which satisfies all of the relevant Quality of Service (QoS) parameters.

The path is fully specified by the source switch all the way to the destination. Path information is stored in a Designated Transit List (DTL), which is forwarded along with the setup message to establish the circuit for a particular call.

# 6.1.5 Hierarchical Routing

Hierarchical routing allows contiguous switches to be organized into peer groups. The switches within each peer group exchange detailed topology information about their own group that is not visible to switches outside the peer group. Similarly, switches within each peer group do not receive detailed topology information about switches outside of their peer group. Instead, each peer group is represented externally by its Peer Group Leader (PGL). The PGLs exchange among themselves summary topology and reachability information pertaining to their respective peer groups.

Although *ForeThought* 5.2.x does not support the dynamic configuration of ATM Forum PNNI PGL hierarchical networks, FORE switches running *ForeThought* 5.2.x can be placed in and interact appropriately within PNNI hierarchical PGL networks. For more information about how *ForeThought* 5.2.x can be configured to connect two ATM Forum PNNI peer groups, see Section 6.3.1.

Each peer group is summarized by its PGL as a single group node. A *node* is a logical entity that resides in a switch and performs routing operations such as discovering other nodes in the network, maintaining a topology database of its peer group, exchanging that database with its neighbors, and computing paths from itself to other nodes in the network.

The interconnection of logical group nodes form a *higher-level topology*. Logical group nodes can, in turn, be grouped into higher-level peer groups that, in turn, are represented by logical group nodes in an even higher level topology, and so on. PGLs of two adjacent peer groups communicate through an SVC that is established between them. To better support QoS and multicast, each switch has knowledge of all higher-level topology information, in addition to the detailed topology information of its own peer group. Higher-level topology information is propagated down the levels by the PGLs.

For example, suppose a switch A.1.1 belongs to the lowest-level peer group A.1, which, in turn, belongs to a higher-level peer group A. Then, A.1.1 has detailed topology information of peer group A.1, as well as higher-level topology information of peer group A. Suppose peer group A.2 also belongs to peer group A. Then, switch A.1.1 knows about the logical group node representing A.2, although it does not know the detailed topology inside A.2.

# 6.2 PNNI Signalling Protocol

PNNI signalling is used to establish point-to-point and point-to-multipoint connections across the network. This protocol is based on ATM Forum UNI signalling with mechanisms added to support source routing and crankback.

## 6.2.1 Source Routing

A DTL is a source route which specifies the preferred call routing path that the *ForeThought* PNNI or PNNI router should use when setting up an SVC. Each DTL is a source route and each entry in the DTL represents a single hop in that source route. Each hop is represented by a *ForeThought* PNNI or PNNI node and the logical output port at that node.

#### 6.2.2 Crankback

During PNNI signalling, a call being processed according to a DTL may encounter a blocked node or link along the designated route. Crankback allows a partial reroute of such a rejected call so that it does not have to be cleared the whole way back to the source. Additionally, an indication of the blockage and subsequent reroute information is sent to the originator of the DTL.

# 6.3 Internetworking between PNNI and FT-PNNI

When ATM Forum PNNI is deployed, some networks running FT-PNNI will be connected to networks running PNNI. To provide internetworking of these protocols, *ForeThought* 5.2.x allows reachability information to be leaked dynamically between FT-PNNI and PNNI peer groups.

## 6.3.1 Gateway Switches and Split Switches

PNNI or FT-PNNI interfaces are attached to nodes (PNNI or FT-PNNI node). When you upgrade a switch to *ForeThought* 5.2.x, it has a single FT-PNNI node by default. *ForeThought* 5.2.x allows you to create a second node on that same switch that runs PNNI, so you have both a single FT-PNNI node and a single PNNI node running concurrently within a single switch. This is called a *gateway switch*. A gateway switch connects a PNNI network with a FT-PNNI network. You can also create two PNNI nodes instead on single switch. This is called a *split switch*. A split switch connects two PNNI areas.

The FT-PNNI and PNNI nodes within a gateway switch and the two nodes within a split switch do not share topology information, but they dynamically exchange reachability information to facilitate connectivity between a source in one network and a destination in another. The means by which this exchange of reachability information occurs is discussed in the next section.

As far as Designated Transit List (DTL) processing is concerned, a gateway switch or a split switch appears as either a DTL terminating or DTL originating node. For example, Figure 6.1 shows that S3 is a DTL terminating node for the FT-PNNI DTL and a DTL originating node for the PNNI DTL for a connection going from switch S1 to switch S4.

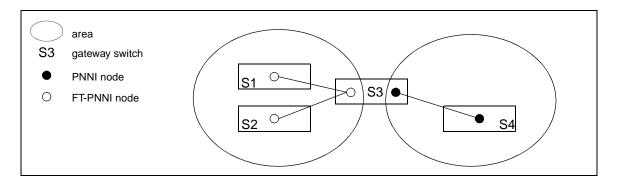


Figure 6.1 - Internetworking of FT-PNNI and PNNI

# 6.3.2 Dynamic Leaking of Reachability Information

This feature allows dynamic leaking of reachability information between the single FT-PNNI node and the single PNNI node within a gateway switch. It may also be used to control the leaking of reachability information between two PNNI peer groups. Using split switches, it possible to connect and exchange reachability information between any number of PNNI peer groups. This dynamic leaking is controlled using the following user-configurable settings: grouping of nodes within a switch into areas and domains, and configuring policies that allow you to summarize, suppress, or advertise reachability addresses between peer groups. See Section 6.3.2.3.1 for more information about policies.

#### 6.3.2.1 Areas

An *area* is a subset of nodes within a domain that are contiguous and that together execute a link state routing protocol to exchange reachability information dynamically among themselves. Because of the database exchange protocol, two nodes that belong to an area will have identical copies of the link-state topology database. Using *ForeThought* 5.2.x, an area can be one of the following:

- a non-hierarchical PNNI network (i.e., a single PNNI peer group because only the lowest-level peer group is implemented)
- a hierarchical FT-PNNI network (i.e., a multiple FT-PNNI peer group system using hierarchy)
- a non-hierarchical FT-PNNI network (i.e., a single FT-PNNI peer group)

Split switches can be used to connect two PNNI areas and gateway switches can be used to connect an FT-PNNI area and a PNNI area, as shown in Figure 6.2. The only restriction in configuring areas is that a FT-PNNI area cannot adjoin another FT-PNNI area.

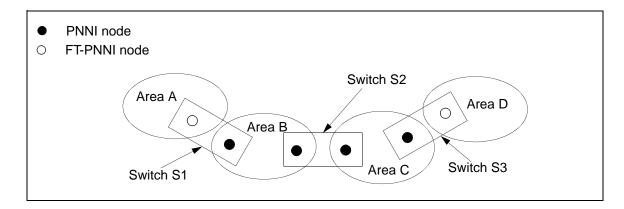


Figure 6.2 - Split Switches and Gateway Switches Connecting Areas

#### 6.3.2.1.1 Peer Groups in Areas

Multiple peer groups can exist within an area. Peer groups within an area are connected to each other by border links as shown in Figure 6.3.

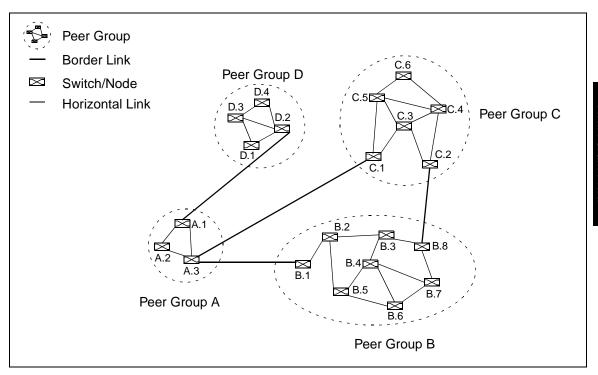


Figure 6.3 - Peer Groups Connected by Border Links

#### 6.3.2.1.2 Area IDs

Each area has an area ID which is unique only within a domain. The area ID is only useful at the border of the area. Since area IDs are included in some of the routing packets distributed throughout a domain, all switches belonging to the same area must be configured with the same area ID. It is especially important for the nodes in split switches and gateway switches to be configured with the correct area ID.

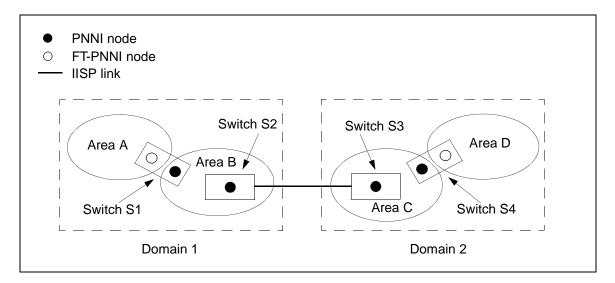
#### 6.3.2.1.3 Levels

Each areas has a *level* associated with it. Levels are used to control the flow of reachability information in a multi-level routing hierarchy (which is discussed in the next section). An N-level hierarchy has N discrete values of levels associated with it, namely,  $L_1$ ,  $L_2$ ,...,  $L_N$ . Each of the areas in the hierarchy has one of these N values as its level. Numerical values are assigned to levels in order from the lowest to the highest; i.e.,  $L_1 < L_2 < L_3 < L_{N_1}$ , so  $L_N$  is the highest level.

An area of level  $L_i$  can adjoin any number of areas at level  $L_{i-1}$ , but can only adjoin at most one area at level  $L_{i+1}$  or higher. This means that on a given switch, there can never be areas configured that belong to more than two distinct levels. Furthermore, there can only be one area that belongs to the higher level.

#### 6.3.2.2 Domains

A *domain* is a group of areas that are configured to dynamically exchange reachability information with one another. This allows connectivity between end systems belonging to different areas without configuring static routes *between areas*. However, reachability information is exchanged *between domains* only if routes are statically configured between the two domains. Figure 6.4 shows an example of how a static route can connect two domains.



**Figure 6.4 -** Static Route Connecting Two Domains

In Figure 6.4, the two nodes configured in switch S1 belong to different areas, and so they do not share the same link-state topology database. The same is true for the nodes configured in switch S4. Dynamic reachability leaking takes place among the nodes in switch S1, enabling end systems in Area A to reach other end systems in Area B, and vice versa within Domain 1. The same applies to the nodes in switch S4.

The node in switch S2 and the node in switch S3 belong to two different areas, each of which is part of a different domain. Exchange of reachability information between the nodes will occur over the IISP link that has been statically configured between the domains.

#### 6.3.2.2.1 Configuring Domains

Domains only need to be configured on split switches and gateway switches; i.e., switches that belong to multiple areas. Other switches operate properly without domain configuration.

One FT-PNNI domain (with domain index 1) is pre-configured by default on switches running *ForeThought* 5.2.x. This domain cannot be deleted, but the information within the domain can be changed. Multiple domains can be configured on a switch. For each domain, the following parameters can be changed:

- The domain name, which is optional, can be configured for easier manageability. However, the name of the default domain cannot be changed.
- The domain protocol, which sets the routing protocol for the switch, can be set to ftpnni, pnni, or gateway. This parameter can also be changed on the default domain. The default setting is ftpnni. If you want to configure a split switch, you must change the domain protocol to pnni. If you want to configure a gateway switch, you must change the domain protocol to gateway.
- The domain prefix can be changed. This sets the default 13-byte NSAP prefix used by the switch's routing protocol(s).
- The default summarization parameter can also be modified. This tells the switch's routing protocol(s) to use the peer group mask (in the case of FT-PNNI) or the PNNI level (in the case of PNNI) to determine the length of the mask to use on the prefix to form the peer group ID.
- A signalling interface (port/VPI) can be associated with a particular domain, while the routing interface (port/VPI) for that signalling interface can be associated with a particular routing node within that domain.

For more information about changing these parameters, see Chapter 1 of the AMI Configuration Commands Reference Manual.

#### 6.3.2.3 Propagation of Reachability Information

The leaking of reachability information between two areas is constrained by two things: policy and scope.

#### 6.3.2.3.1 Policies

*ForeThought* software lets you configure a policy as a flexible means of enforcing security across the network topology. When a node discovers an address that is leaked by another node, it checks the *policy* to determine if the address is to be advertised within its own area. There are three types of policies:

- summary Addresses matching a summary policy cause just the summarized prefix of the address to be announced to the node's peer group.
- suppress Addresses matching a suppress policy are <u>not</u> announced to the node's peer group at all.
- advertise Addresses matching an advertise policy cause the entire address to be announced to the node's peer group.

In each node's policy table, the switch chooses the best address match, meaning the longest prefix match, to determine which policy applies to an address. Therefore, the suppress and advertise policies provide limited filtering since addresses matching a broad prefix can be filtered at the split switch or gateway switch through the suppress policy, then a particular service, device, or switch can be exempted by advertising an address within the suppressed range. The summary policy is best for scalability since one address prefix is shared, rather than dozens or hundreds of more specific addresses.

For example, suppose you have an area A which contains peer groups A.1 and A.2. Suppose you create two policies: one says suppress all address that are area A and the other policy says advertise all address that are peer group A.1. If an address comes in to the area that has the prefix for A.1, the policy for advertise takes precedence over the policy for suppress (the advertise policy is a longer prefix match). Therefore, the address will be advertised. However, if an address comes in to the area that has a prefix for A.2, it will be suppressed (since there is no advertise or summary match for A.2).

#### 6.3.2.3.2 Scope

Each piece of reachability information has a *source area ID* and a *scope* associated with it. The source area ID of reachability information originated in a given area is 0 within that area, meaning that is local. The source area ID of reachability information advertised from, for example, area A1 into area A2 has source area ID A1.

When a leaked address is to be advertised or summarized, a node also determines the scope with which to advertise the address (or its summary). The scope denotes the highest level in the PNNI routing hierarchy at which an address can be advertised. Scoping ensures that reachability information does not loop. A reachable address being advertised with a scope of l can only be leaked to areas of level lower (larger in numerical value) than l. The scope of reachability information originated in a given area is 0 within that area, meaning that it is local. The scope of reachability information advertised from, for example, area A1 (at level  $l_1$ ) into area A2 (at level  $l_2$ ) is  $l_2$  if  $l_2$  is a lower level than  $l_1$ , but is  $l_1$  if  $l_2$  is a higher level than  $l_1$ .

#### 6.3.2.3.3 The Process for Leaking Reachability Information

The following process is used by a node to determine if and how a leaked reachable address is to be advertised to its peer group:

- 1. The node finds the longest matching policy prefix for the address.
- 2. If there is no matching policy prefix, the leaked address is advertised as is. If there are multiple such addresses with different scopes, then the one with the widest (smallest in numerical value) scope is used to determine the scope with which to advertise the address.
- 3. If the longest matching policy prefix is a policy suppress, the leaked address is not advertised at all.
- 4. If the longest matching policy prefix is a policy advertise or summary, the policy prefix is advertised or summarized. If there are multiple leaked addresses with different scopes for which the policy prefix is longest matching, then the one with the widest scope is used to determine the scope with which to advertise or summarize the policy prefix.

The following rule is used to determine the scope with which to advertise an address or a policy prefix:

• If the address used to determine the scope (chosen above) was leaked from a strictly lower level (numerically larger) area, then the scope of the address is used; otherwise, the level of the area to which the advertising node belongs is used.

Metrics can optionally be configured with a policy prefix. In this case, the configured metrics are included in the advertisement of the policy prefix.

#### 6.3.2.4 VP Trunk QoS Extension

This feature allows QoS parameters (CTD, CDV, and CLR) to be configured for VP trunks. For a PNNI interface that is configured on a VP trunk, the associated QoS parameters are computed as the combination of the parameters configured on the VP trunk and those that are part of the switch and network modules. This feature also allows a particular service category to be configured as "not supported," in which case the associated PNNI link does not advertise the service category, thus preventing connections with the given service category from using the link.

#### ATM Forum PNNI

# CHAPTER 7 Signalling

This chapter contains information that pertains to the signalling portion of *ForeThought* software. This signalling information is described in the following sections:

- Section 7.1 VCI Allocation Range
- Section 7.2 Signalling Scope
- Section 7.3 Signalling Channel Auto Configuration Procedures
- Section 7.4 Allowable Combination of Traffic Parameters

# 7.1 VCI Allocation Range

The VCI allocation range is the range of VCI values in a path within which a signalling channel will allocate VCs. This range defaults to the VCI-range of the virtual path containing the signalling channel (also referred to as the *containing path*).



If the signalling channel is a link-scope signalling channel, then this range applies to the containing path and all other paths within which the signalling channel has allocated connections.

A VCI allocation range can be specified when a signalling channel is created using the ATM Management Interface (AMI) command conf signalling new. The range is computed differently depending on whether or not you have ILMI running.

# 7.1.1 Determining the VCI Allocation Range with ILMI Down

If ILMI is not operational and if the VCI allocation range of the signalling channel is specified, then the range is determined as an intersection of the following:

- the VCI-range of the path containing the signalling channel
- · the range specified for the signalling channel

For example, if you create a signalling channel with ILMI down on a VP that has no VCIs reserved, and you specify a VCI range:

```
conf signalling new 1A3 65 -ilmi down -minvci 70 -maxvci 200
```

the actual range is computed as an intersection based on the range you enter (70 - 200) and the range available in that VP (1 - 511).

To display what the actual VCI range is, enter the following command:

#### conf signalling show atm

				Admin	Admin	Oper	Oper		
Port	VPI	SigVCI	ILMIVCI	MinVCI	MaxVCI	MinVCI	MaxVCI	InSigUpc	OutSigService
1A1	0	5	16	32	511	32	511	0	vbr
1A2	0	5	16	32	511	32	511	0	vbr
1A3	0	5	16	32	511	32	511	0	vbr
1A3	65	5	16	70	200	70	200	0	vbr
1A4	0	5	16	32	511	32	511	0	vbr
1CTL	0	5	16	32	1023	32	1023	0	vbr

In this display, the range that you have requested is shown in the Admin MinVCI and Admin MaxVCI fields. The range that is computed and allowed is displayed in the Oper MinVCI and Oper MaxVCI fields. In this example, no VCIs are reserved on VP 65, so the requested range of 70 - 200 is allowed.

In another example, if there are VCIs already reserved on VP 65, as shown below in the MinVCI and MaxVCI fields:

#### configuration vpt show

	Input		Outpu	ıt						
	Port	VPI	Port	VPI	ResBW	CurBW	MinVCI	MaxVCI	VCs	Protocol
	1A1	0	termi	nate	N/A	37.3K	1	511	35	pvc
	1A2	0	termi	nate	N/A	0.8K	1	511	6	pvc
	1A3	0	termi	nate	N/A	0.8K	1	511	6	pvc
	1A3	65	termi	nate	N/A	0.8K	1	120	6	pvc
	1A4	0	termi	nate	N/A	0.8K	1	511	6	pvc
	1CTL	0	termi	nate	N/A	45.4K	1	1023	85	pvc
	origi	nate	1A1	0	N/A	1.7K	1	511	7	pvc
	origi	nate	1A2	0	N/A	0.8K	1	511	6	pvc
	origi	nate	1A3	0	N/A	0.8K	1	511	6	pvc
	origi	nate	1A3	65	N/A	0.8K	1	120	6	pvc
	origi	nate	1A4	0	N/A	0.8K	1	511	6	pvc
]	Press :	retur	n for	more,	q to qu	it: q				

and you enter the following:

conf signalling new 1A3 65 -ilmi down -minvci 70 -maxvci 200

then the actual range is computed as an intersection based on the range you enter (70 - 200) and the range available in that VP (1 - 120).

To see what the actual VCI range is, enter the following command:

#### conf signalling show atm

			Admin	Admin	Oper	Oper		
VPI	SigVCI	ILMIVCI	MinVCI	MaxVCI	MinVCI	MaxVCI	InSigUpc	OutSigService
0	5	16	32	511	32	511	0	vbr
0	5	16	32	511	32	511	0	vbr
0	5	16	32	511	32	511	0	vbr
65	5	16	70	200	70	120	0	vbr
0	5	16	32	511	32	511	0	vbr
0	5	16	32	1023	32	1023	0	vbr
	0 0 0 65 0	0 5 0 5 0 5 65 5	0 5 16 0 5 16 0 5 16 65 5 16 0 5 16	VPI SigVCI ILMIVCI MinVCI 0 5 16 32 0 5 16 32 0 5 16 32 65 5 16 70 0 5 16 32	VPI         SigVCI         ILMIVCI         MinVCI         MaxVCI           0         5         16         32         511           0         5         16         32         511           0         5         16         32         511           65         5         16         70         200           0         5         16         32         511	VPI         SigVCI         ILMIVCI         MinVCI         MaxVCI         MinVCI           0         5         16         32         511         32           0         5         16         32         511         32           0         5         16         32         511         32           65         5         16         70         200         70           0         5         16         32         511         32	VPI         SigVCI         ILMIVCI         MinVCI         MaxVCI         MinVCI         MaxVCI           0         5         16         32         511         32         511           0         5         16         32         511         32         511           0         5         16         32         511         32         511           65         5         16         70         200         70         120           0         5         16         32         511         32         511	VPI         SigVCI         ILMIVCI         MinVCI         MaxVCI         MinVCI         MaxVCI         InSigUpc           0         5         16         32         511         32         511         0           0         5         16         32         511         32         511         0           0         5         16         32         511         32         511         0           65         5         16         70         200         70         120         0           0         5         16         32         511         32         511         0

In this example, since VCIs 1 - 120 are already reserved on VP 65, the range of 70 - 120 is the actual range allowed.

# 7.1.2 Determining the VCI Allocation Range with ILMI Up

If ILMI is operational, then the peer switch's VCI allocation range is obtained by reading its ILMI MIB variable atmfAtmLayerMaxVccs. The VCI allocation range of the signalling channel is then determined as an intersection of the following:

- the VCI-range of the path containing the signalling channel
- the range specified for the signalling channel (if applicable)
- the VCI-range of the peer signalling entity

For example, if your peer has no VCIs reserved, and you create a signalling channel with ILMI up on a VP that has no VCIs reserved, and you specify a VCI range:

the actual range is computed as an intersection based on the range you enter (70 - 200), the range available in that VP (1 - 511), and the range of the peer signalling entity (32 - 511).

To display what the actual VCI range is, enter the following command:

#### conf signalling show atm

				Admin	Admin	Oper	Oper		
Port	VPI	SigVCI	ILMIVCI	MinVCI	MaxVCI	MinVCI	MaxVCI	InSigUpc	OutSigService
1A1	0	5	16	32	511	32	511	0	vbr
1A2	0	5	16	32	511	32	511	0	vbr
1A3	0	5	16	32	511	32	511	0	vbr
1A3	65	5	16	70	200	70	200	0	vbr
1A4	0	5	16	32	511	32	511	0	vbr
1CTL	0	5	16	32	1023	32	1023	0	vbr

In this display, the range that you have requested is shown in the Admin MinVCI and Admin MaxVCI fields. The range that is computed and allowed is displayed in the Oper MinVCI and Oper MaxVCI fields. In this example, no VCIs are reserved on VP 65 and the signalling peer supports the range of 32 - 511, so the requested range of 70 - 200 is allowed.

In another example, if the peer supports the range of 32 - 511, but there are VCIs already reserved on VP 65, as shown below in the MinVCI and MaxVCI fields:

#### configuration vpt show

	Input		Outpu	ıt						
	Port	VPI	Port	VPI	ResBW	CurBW	MinVCI	MaxVCI	VCs	Protocol
	1A1	0	termi	nate	N/A	37.3K	1	511	35	pvc
	1A2	0	termi	nate	N/A	0.8K	1	511	6	pvc
	1A3	0	termi	nate	N/A	0.8K	1	511	6	pvc
	1A3	65	termi	nate	N/A	0.8K	1	120	6	pvc
	1A4	0	termi	nate	N/A	0.8K	1	511	6	pvc
	1CTL	0	termi	nate	N/A	45.4K	1	1023	85	pvc
	origi	nate	1A1	0	N/A	1.7K	1	511	7	pvc
	origi	nate	1A2	0	N/A	0.8K	1	511	6	pvc
	origi	nate	1A3	0	N/A	0.8K	1	511	6	pvc
	origi	nate	1A3	65	N/A	0.8K	1	120	6	pvc
	origi	nate	1A4	0	N/A	0.8K	1	511	6	pvc
]	Press :	retur	n for	more,	q to qu	it: q				

and you enter the following:

conf signalling new 1A3 65 -ilmi up -minvci 70 -maxvci 200

then the actual range is computed as an intersection based on the range you enter (70 - 200), the range available in that VP (1 - 120), and the range supported by the peer (32 - 511).

To see what the actual VCI range is, enter the following command:

#### conf signalling show atm

				Admin	Admin	Oper	Oper			
Port	VPI	SigVCI	ILMIVCI	MinVCI	MaxVCI	MinVCI	MaxVCI	InSigUpc	OutSigService	
1A1	0	5	16	32	511	32	511	0	vbr	
1A2	0	5	16	32	511	32	511	0	vbr	
1A3	0	5	16	32	511	32	511	0	vbr	
1A3	65	5	16	70	200	70	120	0	vbr	
1A4	0	5	16	32	511	32	511	0	vbr	
1CTL	0	5	16	32	1023	32	1023	0	vbr	

In this example, since the peer supports VCIs 32 - 511, but VCIs 1 -120 are already reserved on VP 65, the range of 70 - 120 is the actual range allowed.

In a third example, if the peer only supports the range of 32 - 180, and VCIs 1 - 120 are already reserved on VP 65, and you enter the following:

```
conf signalling new 1A3 65 -ilmi up -minvci 70 -maxvci 200
```

then the actual range is computed as an intersection based on the range you enter (70 - 200), the range available in that VP (1 - 120), and the range supported by the peer (32 - 180).

To see what the actual VCI range is, enter the following command:

#### conf signalling show atm

				Admin	Admin	Oper	Oper		
Port	VPI	SigVCI	ILMIVCI	MinVCI	MaxVCI	MinVCI	MaxVCI	InSigUpc	OutSigService
1A1	0	5	16	32	511	32	511	0	vbr
1A2	0	5	16	32	511	32	511	0	vbr
1A3	0	5	16	32	511	32	511	0	vbr
1A3	65	5	16	70	200	70	120	0	vbr
1A4	0	5	16	32	511	32	511	0	vbr
1CTL	0	5	16	32	1023	32	1023	0	vbr

In this example, since the peer supports VCIs 32 - 180, and VCIs 1 - 120 are already reserved on VP 65, the range of 70 - 120 is the actual range allowed.

# 7.2 Signalling Scope

A signalling channel can be configured to control either VCs within its own VP (VP-scope, the default) or to control VCs in its own VP, in addition to controlling VCs in other VPs on the link (link-scope).

The following criteria must be satisfied in order to use link-scope signalling channels:

- There can be only one link-scope signalling channel on a given link.
- A link-scope signalling channel must be created within VP 0 on the given link.
- The path containing the link-scope signalling channel must be an elastic path.
- A backplane signalling channel in an ASX-1000 cannot be a link-scope signalling channel.

# 7.2.1 VC-Space

The VC-space of a signalling channel is the VPI/VCI space within which a signalling channel can allocate connections. This is determined by the allocation scope of the signalling channel and the VCI allocation range (see Section 7.1) for the signalling channel.

The VC-space of a VP-scope signalling channel is the VCI allocation range of the signalling channel applied only to the path that contains the signalling channel.

The VC-space of a link-scope signalling channel is comprised of the aggregate VCI allocation range of the signalling channel applied to all paths controlled by the signalling channel (i.e., the path containing the signalling channel and all other paths in which the signalling channel has allocated connections).

# 7.2.2 Dynamic Paths

A link-scope signalling channel allocates new virtual channel connections based on the availability of virtual channel connections within the VC-space of the link-scope signalling channel at that time. A link-scope signalling channel allocates connections in a new originating/terminating path (called a *dynamic path*, because signalling uses such a path on the fly) when its current VC-space is exhausted.



Every time a link-scope signalling channel uses a new path for allocating connections, its VC-space increases.

To avoid conflicts of ownership of the path, a link-scope signalling channel never uses an existing provisioned terminating/originating path (apart from the path that contains the signalling channel) to allocate virtual channel connections.

Dynamic paths inherit the characteristics of the virtual path that contains the signalling channel (called the *containing path*). Therefore, dynamic paths have the same VCI-range as the containing path and are elastic paths.

To see which paths are dynamic paths, look at the Protocol field associated with each path under the AMI command conf vpt show. This field displays q2931 for dynamic paths and displays pvc for provisioned terminating/originating paths.

A dynamic path is controlled only by the link-scope signalling channel that uses it. Therefore, a dynamic path is not open to any provisioning such as the creation of a PVC connection, SPANS signalling path, etc.

# 7.3 Signalling Channel Auto Configuration Procedures

This section provides an overview of signalling channel auto configuration and describes some rules for configuring signalling interfaces. Auto configuration requires that ILMI is up on that signalling interface.

# 7.3.1 Overview of Signalling Channel Auto Configuration

Signalling support for UNI 3.0, UNI 3.1, and UNI 4.0 is provided in *ForeThought* 5.2.x. The switch can be configured to detect which interfaces are UNI 3.0, which are UNI 3.1, and which are UNI 4.0 automatically. The switch also translates signalling messages between interfaces of different types.

When a signalling channel is created using conf signalling new -version auto, then the creation is delayed until a response is obtained from the peer or until the ILMI keep alive requests time out. Based on the peer response, the appropriate SSCOP stack is initialized. If the peer does not support the atmfAtmLayerUniVersion variable, then the switch software defaults to UNI 3.0. Similarly, if the ILMI keep alive requests time out, then the switch software defaults to UNI 3.0. The switch periodically checks the version supported on the peer by querying the atmfAtmLayerUniVersion variable. If the peer version has changed since the last query, the signalling stack is recreated to reflect the change.

The following two tables summarize the above information. Table 7.1 highlights the actions taken based on the configurations of the FORE switches on both sides of an interface. *ForeThought* versions prior to *ForeThought* 4.1.x always default to UNI 3.0. Both UNI 3.1 and UNI 4.0 use SSCOP version 31.

Table 7.1 - Action Taken Based on Both Switches' Signalling Channel Configurations

Configured Version	Peer's Configured Version	Action Taken
UNI 3.0	UNI 3.0 UNI 3.1 UNI 4.0 PNNI 1.0 Auto	SSCOP 30 stack initialized SSCOP 30 stack initialized; error message displayed on syslog SSCOP 30 stack initialized; error message displayed on syslog SSCOP 30 stack initialized; error message displayed on syslog SSCOP 30 stack initialized; peer changes version to UNI 3.0
UNI 3.1	UNI 3.0 UNI 3.1 UNI4.0 PNNI 1.0 Auto	SSCOP 31 stack initialized; error message displayed on syslog SSCOP 31 stack initialized SSCOP 31 stack initialized; error message displayed on syslog SSCOP 31 stack initialized; misconfiguration SSCOP 31 stack initialized; peer changes version to UNI 3.1
PNNI 1.0	UNI 3.0 UNI 3.1 UNI4.0 PNNI 1.0 Auto	SSCOP 31 stack initialized; error message displayed on syslog SSCOP 31 stack initialized; misconfiguration SSCOP 31 stack initialized; misconfiguration SSCOP 31 stack initialized SSCOP 31 stack initialized; peer changes version to PNNI 1.0 if it supports PNNI
UNI 4.0	UNI 3.0 UNI 3.1 UNI4.0 PNNI 1.0 Auto	SSCOP 31 stack initialized; error message displayed on syslog SSCOP 31 stack initialized; error message displayed on syslog SSCOP 31 stack initialized SSCOP 31 stack initialized; misconfiguration SSCOP 31 stack initialized; peer changes version to UNI 4.0
Auto	UNI 3.0 UNI 3.1 UNI4.0 PNNI 1.0	SSCOP 30 stack initialized SSCOP 31 stack initialized SSCOP 31 stack initialized SSCOP 31 stack initialized SSCOP 31 stack initialized Both stacks initialized as SSCOP 31 if both sides support UNI 3.1 or UNI 4.0; otherwise, both sides are initialized to SSCOP 30. If both sides support PNNI then, both sides are initialized to SSCOP 31 and the version is set to PNNI 1.0.

An alternative way of looking at the auto configuration procedure is to look at the atmfAtmLayerUniVersion variable. Table 7.2 shows actions taken based on the atmfAtmLayerUniVersion variable supported by the peer switch.

**Table 7.2 -** Action Taken Based on the Peer's Supported MIB Variable

Configured UNI Version	atmfAtmLayer UNI Version	Action Taken
UNI 3.0	Version 3.0 Version 3.1 Version 4.0 Version 2.0 MIB variable not supported	SSCOP 30 stack initialized SSCOP 30 stack initialized; error message displayed on syslog SSCOP 31 stack initialized; error message displayed on syslog SSCOP 30 stack initialized; version 2.0 is not supported and error message displayed on syslog SSCOP 30 stack initialized
UNI 3.1	Version 3.0 Version 3.1 Version 4.0 Version 2.0  MIB variable not supported	SSCOP 31 stack initialized; error message displayed on syslog SSCOP 31 stack initialized SSCOP 31 stack initialized; error message displayed on syslog SSCOP 30 stack initialized; version 2.0 is not supported and error message displayed on syslog SSCOP 31 stack initialized; error message displayed on syslog
UNI 4.0	Version 3.0 Version 3.1 Version 4.0 Version 2.0 MIB variable not supported	SSCOP 30 stack initialized SSCOP 30 stack initialized; error message displayed on syslog SSCOP 31 stack initialized SSCOP 30 stack initialized; version 2.0 is not supported and error message displayed on syslog SSCOP 30 stack initialized
Auto	Version 3.0 Version 3.1 Version 4.0 Version 2.0  MIB variable not supported	SSCOP 30 stack initialized SSCOP 31 stack initialized SSCOP 31 stack initialized SSCOP 30 stack initialized; version 2.0 is not supported and error message displayed on syslog SSCOP 30 stack initialized

# 7.3.2 Rules for Signalling Channel Auto Configuration

In *ForeThought* 5.0.x and greater, there are two rules for signalling channel auto configuration. The first deals with specifying the signalling interface type and signalling interface version. The second rule concerns specifying the signalling scope and mode.

### 7.3.2.1 Specifying the Type and Interface Version

In AMI, the signalling interface type and signalling interface version are now treated as a pair. The pertinent portion of the AMI syntax is listed here for reference:

```
myswitch::configuration signalling> new <port> <vpi>
        [-version (auto | uni30 | uni31 | pnni10 | uni40)]
         [-type (auto | publicUNI | IISP | privateNNI | privateUNI)]
```

The type and version either both have to be auto or both have to be given a parameter other than auto. Specifying the type as auto and hard configuring the version or vice-versa is not allowed. If a type is not entered, then the type parameter defaults to auto. Similarly, the version parameter defaults to auto if a version is not entered.

The one exception to this rule (to allow backward compatibility) is for the type to be entered as publicUNI and the version to be entered as auto. Table 7.3 shows the valid combinations.

Туре	Version
auto	auto
privateUNI	uni30
privateUNI	uni31
privateUNI	uni40
publicUNI	auto <sup>1</sup>
publicUNI	uni30
publicUNI	uni31
publicUNI	uni40
IISP	uni30
IISP	uni31
IISP	uni40
privateNNI	pnni10

**Table 7.3 -** Valid Type and Version Combinations

<sup>&</sup>lt;sup>1.</sup> This is the only exception to the rule.

#### 7.3.2.1.1 Examples of Valid Configurations

The following are examples of valid configurations to enter into AMI:

```
conf sig new lal 0

conf sig new lal 0 -type privateUNI -version uni30

conf sig new lal 0 -type privateUNI -version uni31

conf sig new lal 0 -type privateUNI -version uni40

conf sig new lal 0 -type publicUNI

conf sig new lal 0 -type publicUNI -version auto

conf sig new lal 0 -type publicUNI -version uni30

conf sig new lal 0 -type publicUNI -version uni31

conf sig new lal 0 -type publicUNI -version uni31

conf sig new lal 0 -type publicUNI -version uni40

conf sig new lal 0 -type IISP -version uni31

conf sig new lal 0 -type IISP -version uni31

conf sig new lal 0 -type IISP -version uni40

conf sig new lal 0 -type publicUNI -version uni40

conf sig new lal 0 -type IISP -version uni40

conf sig new lal 0 -type privateNNI -version pnni10
```

The first example is the default that is most often used when configuring a signalling interface. Using this specification allows the interface to auto configure to any possible valid combination like privUNI/uni30, privUNI/uni31, privUNI/uni40, ftpnni/uni30, ftpnni/uni31, ftpnni/uni40, or privateNNI/pnni10. This is the most preferred input.

The fifth and sixth examples illustrate the exception to the rule.

## 7.3.2.1.2 Examples of Invalid Configurations

Table 7.4 shows examples of invalid configurations to enter into AMI and the reasons why they are invalid.

**Table 7.4 -** Invalid Type and Version Combinations

Invalid Combination	Reason
conf sig new 1a1 0 -version uni30	The type is not entered (auto).
conf sig new 1a1 0 -type auto -version uni30	The type is auto.
conf sig new 1a1 0 -version uni31	The type is not entered (auto).
conf sig new 1a1 0 -type auto -version uni31	The type is auto.
conf sig new 1a1 0 -version uni40	The type is not entered (auto).
conf sig new 1a1 0 -type auto -version uni40	The type is auto.
conf sig new 1a1 0 -version pnni10	The type is not entered (auto).
conf sig new 1a1 0 -type auto -version pnni10	The type is auto.
conf sig new 1a1 0 -type privateUNI	The version is not entered (auto).
conf sig new 1a1 0 -type privateUNI -version auto	The version is auto.
conf sig new 1a1 0 -type privateUNI -version pnni10	pnni10 is an NNI protocol.
conf sig new 1a1 0 -type publicUNI -version pnni10	pnni10 is an NNI protocol.
conf sig new 1a1 0 -type IISP	The version is not entered (auto).
conf sig new 1a1 0 -type IISP -version auto	The version is auto.
conf sig new 1a1 0 -type IISP -version pnni10	IISP and pnni10 are conflicting NNI protocols.
conf sig new 1a1 0 -type privateNNI	The version is not entered (auto).
conf sig new 1a1 0 -type privateNNI -version auto	The version is auto.
conf sig new 1a1 0 -type privateNNI -version uni30	The NNI type and the UNI version are conflicting protocols.
conf sig new 1a1 0 -type privateNNI -version uni31	The NNI type and the UNI version are conflicting protocols.
conf sig new 1a1 0 -type privateNNI -version uni40	The NNI type and the UNI version are conflicting protocols.

### 7.3.2.2 Specifying the Scope and Mode

In AMI, the signalling scope (designated in AMI as <code>-sig\_alloc</code>) and the signalling mode (designated in AMI as <code>-sig\_mode</code>) are now treated as a pair. The pertinent portion of the AMI syntax is listed here for reference:

```
myswitch::configuration signalling> new <port> <vpi>
protocol options:
    [-sig_alloc (vp | link | auto)]
    [-sig_mode (nonAssoc | vpAssoc | auto)]
```

The scope and mode either both have to be auto or both have to be given a parameter other than auto. Specifying the scope as auto and hard configuring the mode or vice-versa is not allowed. If a scope is not entered, then the scope parameter defaults to auto. Similarly, the mode parameter defaults to auto if a mode is not entered.

Table 7.5 shows the valid combinations.

**Table 7.5 -** Valid Scope and Mode Combinations

Scope	Mode
autoScope	autoMode
vpScope	vpAssoc
vpScope	nonAssoc1
linkScope	nonAssoc

<sup>1.</sup> This combination is valid as long as the -type is not configured as auto or privateNNI. If the -type is auto, the link could become a PNNI link, which would conflict with a -mode of nonAssoc and a -scope of vpScope.

### 7.3.2.2.1 Examples of Valid Configurations

The following are examples of valid configurations to enter into AMI:

```
conf sig new 1a1 0
conf sig new 1a1 0 -sig_alloc auto -sig_mode auto
conf sig new 1a1 0 -sig_alloc vp -sig_mode vpAssoc
conf sig new 1a1 0 -sig_alloc link -sig_mode nonAssoc
```

The first example is the default that is most often used when configuring a signalling interface. This is the most preferred mode since the auto configuration of the link and scope procedures can correctly determine what is valid for a given interface.

### 7.3.2.2.2 Examples of Invalid Configurations

Table 7.6 shows examples of invalid configurations to enter into AMI and the reasons why they are invalid.

**Table 7.6 - Invalid Scope and Mode Combinations** 

Invalid Combination	Reason
conf sig new 1a1 0 -sig_mode vpAssoc	The scope is not entered (auto).
conf sig new 1a1 0 -sig_alloc auto -sig_mode vpAssoc	The scope is auto.
conf sig new 1a1 0 -sig_mode nonAssoc	The scope is not entered (auto).
conf sig new 1a1 0 -sig_alloc auto -sig_mode nonAssoc	The scope is auto.
conf sig new 1a1 0 -sig_alloc vp	The mode is not entered (auto).
conf sig new 1a1 0 -sig_alloc vp -sig_mode auto	The mode is auto.
conf sig new 1a1 0 -sig_alloc link	The mode is not entered (auto).
conf sig new 1a1 0 -sig_alloc link -sig_mode auto	The mode is auto.
conf sig new 1a1 0 -sig_alloc link -sig_mode vpAssoc	This is an invalid combination for an interface.
conf sig new 1a1 0 -sig_alloc vp -sig_mode nonAssoc	This is an invalid combination if the type is auto (default).
conf sig new 1a1 0 -type privateNNI -version pnni10 -sig_alloc vp -sig_mode nonAssoc	VP scope and nonAssoc mode are incompatible with the type privateNNI.

There is a potential invalid configuration that can occur because the auto configuration of the operating scope and mode depend on whether a path is an elastic path or not. A non-elastic path has bandwidth reserved for it from the link bandwidth at the time that the path is created (when the **-reserved** option in **conf vpt new** is used.) An elastic path does not have reserved bandwidth at creation time. Instead, it uses the residual bandwidth available on the link.

The invalid configuration can occur when a signalling interface on VPI 0 that has both the scope and mode configured as auto becomes a PNNI interface (either by the user explicitly specifying -type privateNNI or by specifying -type auto). In this case, if the path is an elastic path, then the operating scope becomes link and the operating mode becomes nonAssoc. If the path is a non-elastic path, then the operating scope becomes vp and the operating mode becomes vpAssoc.

Therefore, if a signalling interface exists between two switches on VPI 0, with one half of this path being elastic on one switch and one half being non-elastic on the other switch, then the interface comes up with incompatible modes on either side if it becomes a PNNI interface. This results in no calls being set up on this interface.

To prevent this invalid configuration, be sure that both sides of your PNNI interface are either elastic or non-elastic. If this invalid configuration occurs, simply delete one side of the path and recreate it to match the other side.

# 7.4 Allowable Combination of Traffic Parameters

### 7.4.1 PNNI 1.0/UNI 4.0

The following information applies to both signalled connections and to PNNI SPVCs.

## 7.4.1.1 Service Categories

The ATM service categories assigned to connections in PNNI 1.0 and UNI 4.0 are defined in the ATM Traffic Management Specification, Version 4.0 (TM 4.0). For PNNI 1.0 interfaces, *ForeThought* 5.2.x supports the service categories defined in TM 4.0 with a few exceptions.

An explicit way of requesting a particular ATM service category is not provided in PNNI 1.0 and UNI 4.0 signalling. Instead, it must be derived from three pieces of information in the SETUP message.

- the broadband bearer class in octet 5 of the broadband bearer capability (BBC) information element
- the absence or presence of the ATM transfer capability (ATC) octet (octet 5a) in the BBC information element
- the value of the ATC, if present, and the absence or presence of the best effort indicator in octet 18 of the ATM Traffic Descriptor information element

Derivation of a service category from the above information is specified in Section A9.2 (Determination of ATM service Category) of the ATM User-Network Interface (UNI) Signalling Specification, Version 4.0 (UNI 4.0). However, *ForeThought* 5.2.x does not support the following from Table A9-1 of this specification:

• The Transparent VP-Service in the BBC information element is not supported.

### 7.4.1.2 Allowable Combination of Traffic Parameters

The parameters specified in the BBC information element, the ATM traffic descriptor information element, the Extended QoS parameter information element, the End-to-End transit delay information element, and the QoS parameter information element of the SETUP message should be consistent.

Table A9-2 of UNI 4.0 shows the allowable combinations. The following exceptions apply to *ForeThought* 5.2.x:

• In the Extended QoS parameters information element, for forward and backward Cell Loss Ratio (CLR), *ForeThought* 5.2.x does not distinguish between the (CLP=0) and CLP (0+1) traffic streams.

• The PNNI 1.0 and UNI 4.0 specifications allow incomplete traffic contracts for VBR in four instances (two for real time VBR and two for non-real time VBR). Check note 7 and 8 of Table A9-2 of the UNI 4.0 specification. *ForeThought* 5.2.x does not support these four combinations.

### 7.4.2 UNI 3.X

FORE Systems supports a subset of traffic contracts which are specified in Table F.1 of the UNI 3.1 specification. These traffic contracts are enforced on both signalled connections as well on PNNI SPVCs. The following are the allowable contracts supported by FORE's switch fabrics.

- CBR: PCR (0), PCR (0+1)
- CBR: PCR (0+1)
- VBR: PCR (0+1), SCR, MBS (0)
- VBR: PCR (0+1),SCR, MBS (0+1)
- UBR: PCR (0+1)



Traffic contracts that do not completely characterize the traffic characteristics of a given service category are incomplete and not supported. (For example, a VBR contract that is received with an ATM traffic descriptor information element that does not specify SCR and MBS cell rates does not fully characterize a VBR contract. Such a contract is, therefore, rejected by the switch fabric).

Table 7.7 shows the subset of the UNI 3.1 allowable combination of traffic parameters supported by *ForeThought* 5.2.x:

Table 7.7 - UNI 3.1 Allowable Combination of Traffic Parameters in ForeThought 5.2.x

Broadband Bearer Class	A	X	С	X	С	X	A	X	С	X
Traffic Type		CBR		&		&		CBR		&
Timing Required		Y		&&		&&		Y		&&
PCR (CLP=0)	S	S								
PCR (CLP=0+1)	S	S	S	S	S	S	S	S	S	S
SCR and MBS (CLP=0)			S	S						
SCR and MBS (CLP=0+1)					S	S				
Best Effort									S	S
Tagging	Y/N	Y/N	Y/N	Y/N	N	N	N	N	N	N
QoS Class	*	*	*	*	*	*	*	*	0	0

The following are used in the table above:

Y Yes

N No

**S** Specified

Y/N Either Yes or No is allowed.

- \* All QoS classes are supported.
- & The parameter is coded to either "no indication," or "VBR," or, for UNI 3.1, Octet 5a (Traffic Type / Timing Required) is absent; these 3 codings are treated as equivalent.
- **&&** The parameter is coded to either "no indication," or "No," or, for UNI 3.1, Octet 5a (Traffic Type / Timing Required) is absent; these 3 codings are treated as equivalent.

A blank entry indicates that the parameter is not present.

# CHAPTER 8 Security

ForeThought software (that is version 5.0.x or greater) provides various forms of switch security. Security can be implemented by creating userids. There are also security methods that prevent access to the switch, which include IP filtering and NSAP filtering. Each of these security methods is described in the following sections.

# 8.1 Configuring Userids

The network administrator creates and assigns a userid for each user or for a group of users via the conf security login new command in the ATM Management Interface (AMI). (See the AMI Configuration Commands Reference Manual for more information about this command.) A userid consists of the following:

- a method of login authentication
- · a level of AMI command privileges
- a level of AMI access levels
- · an optional password

On a new switch running ForeThought 5.0.x or greater, there are two separate default userids: ami and asx. Both are configured with the local authentication method, with admin privileges (meaning you are allowed to use all AMI commands), and all access (meaning you are allowed to login to the switch using all the possible methods). Both userids are assigned a null password.

The network administrator should configure userids for all people who will have access to each switch. Changing a userid only changes it in the switch's local login file. If the administrator wants to use the same set of userids on each switch in the network, he or she should configure the entire set of userids on any one switch. Then, he or she can backup the login file to a host using the conf security login backup command. Then, the administrator can copy the login file to each of the other switches in the network using the conf security login restore command on each switch. (See the AMI Configuration Commands Reference Manual for more information about these commands.)

If the administrator wants to change, add, or delete one or more userids and propagate these changes to the other switches, he or she should use the same method. First, make all of the modifications on one switch, back up the login file to a host, and restore the login file to each of the other switches.

There are several different scenarios that can occur, such as the SecurID server being down or a userid not being listed on a particular switch. Table 8.1 shows the different login scenarios and the action taken by the switch for each. In this table, port refers to the access method assigned to the userid (e.g., telnet or serial port).

**Table 8.1 -** Possible Login Scenarios

Login Scenarios	Action
Userid not listed and SecurID server not accessible	Reject
Userid not listed; SecurID server accessible; and SecurID server rejects userid	Reject
Userid not listed; SecurID server accessible; and SecurID server accepts userid	Accept
Userid listed and port not OK	Reject
Userid listed; port OK; local authentication; and supplied password does not match	Reject
Userid listed; port OK; local authentication; and supplied password matches	Accept
Userid listed; port OK; SecurID authentication; SecurID server accessible; and SecurID server rejects userid	Reject
Userid listed; port OK; SecurID authentication; SecurID server accessible; and SecurID server accepts userid	Accept
Userid listed; port OK; SecurID authentication; SecurID server not accessible; and no local password	Reject
Userid listed; port OK; SecurID authentication; SecurID server not accessible; local password; and supplied password does not match	Reject
Userid listed; port OK; SecurID authentication; SecurID server not accessible; local password; and supplied password matches	Accept

# 8.1.1 Login Authentication Method

The network administrator can configure two different forms of login authentication: local authentication and SecurID authentication. The administrator may employ either method for all users, or he or she may choose to employ the local method for some users and the SecurID method for other users. Each method is described in the following sections.

#### 8.1.1.1 Local Authentication

When a user is configured for local password authentication, he or she is prompted for a login ID (userid) and a password which is stored locally in the switch whenever he or she attempts to open an AMI session either via telnet or via the serial port. (The login password is not required if the user tries to open an AMI session via a remote switch or via *ForeView*.) After a validation check is made based on the scenarios listed in Table 8.1, an AMI session is started (provided that a local AMI session is not already running).

For specific information and examples of how to log in to the switch via telnet, via the serial port, or via a remote switch, see the *ATM Management Interface (AMI) Manual*. For information about logging in via *ForeView*, see the *ForeView Network Management User's Manual*.

#### 8.1.1.2 SecurID Authentication

When a user is configured for SecurID authentication and he or she attempts to log in, the user is prompted for a login ID (userid) and a SecurID passcode. The two-part passcode consists of: a secret, memorized personal identification number (PIN) and the current code generated by the user's assigned SecurID token. (The passcode is not required if the user tries to open an AMI session via a remote switch or via *ForeView*.) After a validation check is made based on the scenarios listed in Table 8.1, an AMI session is started (provided that a local AMI session is not already running).

#### 8.1.1.2.1 SecurID Protection on Switches

Security Dynamics ACE/server and client software<sup>1</sup> prevents a user from logging into AMI locally on a switch until the passcode entered by the user has been validated using an external security server. The server uses SecurID tokens to validate the identity of users, and allows access only to authorized users on valid clients (switches).



Because SecurID does not protect SNMP, ILMI, remote AMI, or *ForeView* access to the switch, it is recommended that users either employ IP filtering as a selective mechanism to allow SNMP changes, or disable SETs from the network entirely.

8 - 3

<sup>&</sup>lt;sup>1.</sup> The client software is already provided on FORE ATM switches, but the server software must be purchased separately from Security Dynamics.

#### 8.1.1.2.2 SecurID Passcode

This authentication method provides a high level of security because the SecurID passcode that allows access to the protected switches is comprised of two parts:

- a secret, memorized personal identification number (PIN)
- the current code generated by the user's assigned SecurID token

#### 8.1.1.2.2.1 PIN Number

The PIN is known only by the user. It can be either alpha-numeric or strictly numerical, and can be either a fixed or variable length from 4 - 8 characters, depending on how the system administrator configures the server.

#### 8.1.1.2.2.2 SecurID Tokens

The second part is a unique code from the SecurID token that only the user possesses and which cannot be counterfeited. Each authorized user on a protected system is assigned a SecurID token to use when accessing a protected switch. SecurID tokens are small, hand-held devices that use microprocessors to calculate and display random codes. These codes change at a specified interval, which is usually once every minute. The random code displayed on a user's token is the same code the server software has generated for that moment.

#### 8.1.1.2.3 SecurID Server

The server can run on a UNIX system or on a Windows NT system. Each FORE switch has a defined set of authorized users. (See the *AMI Configuration Commands Reference Manual* for more information.) When a user is designated as someone who has SecurID authorization, that user is not permitted access to that switch until his or her identity is validated based on the scenarios listed in Table 8.1.

When a user attempts to login with a SecurID passcode, the SecurID software running on the switch verifies the passcode with the server; verifies the authenticity of the server so that no other machine can pretend to be the server in order to capture security data; and encrypts and decrypts messages sent between the switch and the server.

#### 8.1.1.2.3.1 Slave Server

A backup, or slave server can be installed to ensure that authentication services are not interrupted, even if the server goes down. When the slave detects that the master has failed, the slave takes over authentication services. This failover is transparent to the user.

#### 8.1.1.2.3.2 Server Database

The server's database includes records for all tokens, a list of switches to be protected, an audit trail of SecurID and administrative activity, and a list of users who are authorized to access each switch.

#### 8.1.1.2.3.3 Data Encryption between the Server and Switches

Messages sent between the server and the switches are encrypted either using the DES algorithm or Security Dynamics proprietary encryption algorithm. The server can use either method, but all switches must use the same algorithm as the one configured on the server.

This data encryption method protects communications between the server and the switches because the first time the switch contacts the server, it receives a node secret file, which is a string of about 16 bytes. This string, which is known only to the server and this switch, is used in encrypting messages between the server and the switch. Additionally, communications between the master and slave servers and stored token information are encrypted. No one, including system administrators, can breach security by inspecting secured token data.

#### 8.1.1.2.4 SecurID AMI Commands

A new AMI menu called conf security login securid has been added to allow the configuration of SecurID on the switch. These commands are described in detail in the AMI Configuration Commands Reference Manual.

#### 8.1.1.2.5 Installing SecurID on a Switch

The following sections describe how to install SecurID on a switch.

#### 8.1.1.2.5.1 Installing the Server Software

A SecurID server must be run to implement the SecurID authentication method for users of the switch. Refer to the Security Dynamics ACE/Server or Client user's manuals for instructions about installing the software on the server.

### 8.1.1.2.5.2 Transferring the Configuration File

As described earlier, the server and the switches need to maintain some common configuration parameters. The desired configuration information is specified in the <code>sdconf.rec</code> file when the server is installed. Once you have installed the server software, copy this file to the switch using the AMI command <code>conf security login securid get</code>. This command uses the tftp or ftp protocol, depending on what is set under <code>conf system protocol</code>, to transfer the specified file from the specified server to the switch.



This get command can be executed only by the users with admin privileges.

This configuration file is read and information is stored in the FLASH so that they persist across reboots. Additionally, the first time the switch contacts the server, it receives a node secret file, which is a string of about 16 bytes. This string, which is known only to the server and this switch, is used in encrypting messages between the server and the switch and is also stored in the FLASH.



Even though this information is stored in the FLASH, it cannot be accessed using any of the oper flash commands. The conf security login securid get, delete, and show commands must be used.

If the configuration file is not found, if the wrong file was copied, or if the file is corrupted, the SecurID service does not work, an appropriate error message is logged to the console, and the user trying to log in is denied access.

#### 8.1.1.2.5.3 Editing the Server Configuration File

Once SecurID is running, it possible to make changes to the sdconf.rec file for the server(s)/clients. The file can be modified on the server itself by running sdsetup again and changing the desired parameters.



After the changes are made on the server, use the conf security login securid get AMI command to copy the updated sdconf.rec file back to all of the switches that are using that file.

A sample of the contents of a sdconf.rec file and a description of the parameters follows:

Server License and Configuration

LICENSE CREATION: Mon May 13 15:06:08 1996

LICENSE ID: 96012648

ACE/Server VERSION: v 2.1.104

FILE OWNERSHIP: root SETUID BIT: off

CLIENT RETRY: 5 times CLIENT TIMEOUT: 5 sec

DES ENCRYPTION: allowed and enabled

DURESS MODE: not allowed
NetSP CLIENTS: not allowed
EXTENDED TACACS: disabled
TACACS PLUS: disabled

MASTER SERVER: linus

MASTER SERVER ADDRESS: 204.95.89.107 SLAVE SERVER: allowed but not configured

AUTHENTICATION SERVICE: securid

PORT NUMBER: 1024

ADDRESSES: By name in host file or name service

BAD PASSCODES: 3
RESPONSE DELAY: 2
TOKENS IN LICENSE: 25
LICENSE CONFIGURATION

This license was created for:

Fore Systems, Inc. 1000 FORE Drive Warrendale, PA 15086

## These parameters are defined as follows:

Parameter	Description
LICENSE CREATION	The date the ACE/server license was created.
LICENSE ID	The ACE/server license ID.
ACE/Server VERSION	The version number of the server software.
FILE OWNERSHIP	The account that owns the ACE/server files; i.e., that has permissions to run the ACE/server administration programs. root is the default, but the login of any other administrator can be specified here.
SETUID BIT	If set, allows anyone who can run the programs to run them with the permissions of the files' owner.
CLIENT RETRY	Specifies how many times a client attempts to establish communications with the server before reporting an error.
CLIENT TIMEOUT	Specifies how many seconds should elapse between attempts to establish client-server communications.
DES ENCRYPTION	Shows if DES or Security Dynamics encryption is used for client-server communications.
DURESS MODE	Allows a user to enter a special PIN to signal that he or she is being forced to log in by an unauthorized person.
NetSP CLIENTS	Used for IBM's Network Security Program.
EXTENDED TACACS	Indicates if the server supports XTACACS clients.
TACACS PLUS	Indicates if the server supports TACACS+ clients.
MASTER SERVER	The name of the master server.
MASTER SERVER ADDRESS	The IP address of the master server.
SLAVE SERVER	Indicates that the slave server may be installed. If it is installed, its name and IP address become part of the configuration record.
AUTHENTICATION SERVICE	The name of this authentication service as it appears in the /etc/services or in a NIS Services file.
PORT NUMBER	The UDP port number that has been assigned for the use of this service.

Parameter	Description
ADDRESSES	Indicates how network devices shall have the IP addresses resolved.
BAD PASSCODES	The number of failed login attempts with incorrect passcodes, after which the server puts the token associated with that userid into Next Token Mode. When the token is in this mode, the next time the same user tries to log in and gets the passcode correct, he or she is also prompted to enter the next token code; i.e., the one that appears next on the token after 60 seconds.
RESPONSE DELAY	Indicates how long an authentication request should be held (by the server) before a response is returned to the client.
TOKENS IN LICENSE	The number of tokens for this license.
LICENSE CONFIGURATION	Information about the owner of the license.

#### 8.1.1.2.5.4 An Example Login Using SecurID

Once the SecurID server is set up and the configuration file has been copied to the switch, all users who are configured for securID will see a passcode prompt after the AMI login prompt when they attempt to log in to that switch. In the following example, a userid eng has been created on a switch that has the sdconf.rec file in FLASH and that has the MCA method for userid eng set to securid. The following is a transcript of how the user logs in, with the switch output in plain courier font and the user input in bold courier font:

```
localhost::> telnet fishtank
Trying 169.144.48.45 ...
Connected to bx02.
Escape character is `^]'.
S_ForeThought_5.1.0 (1.1144) (asx200bx) (fishtank)
login: eng <ENTER>
Enter PASSCODE: <PIN><Code on SecurID Token> <ENTER>
```

If the passcode entered by the user is accepted, an AMI session to the switch is started. If the server denies access to the user, the following message appears on the screen: Login incorrect. After three unsuccessful attempts, the telnet connection is torn down by the switch.

# 8.1.2 AMI Command Privileges

The network administrator can configure two different levels of AMI command privileges for userids: admin privileges and user privileges. Each method is described in the following sections.

### 8.1.2.1 Admin Privileges

A person whose userid is configured with admin privileges is allowed to access and use all of the AMI commands. This level of privileges should be reserved for the system administrator.

### 8.1.2.2 User Privileges

A person whose userid is configured with user privileges is allowed to access and use all AMI commands, except the following: conf security login backup, conf security login delete, conf security login modify, conf security login new, conf security login password (to modify passwords other than your own), conf security login restore, conf security login show, all conf security login securid commands, all conf security ipaccess commands, all conf security nsapfiltering commands, conf snmp sets, and all debug commands. The default is admin.

### 8.1.3 AMI Access Levels

The network administrator can configure four different levels of AMI access for userids: serial, network, all, and none. Each method is described in the following sections.

### 8.1.3.1 Serial Access

A person whose userid has a serial level of access is allowed to login to the switch only via the serial port.

### 8.1.3.2 Network Access

A person whose userid has a network level of access is allowed to login to the switch only via telnet.

### 8.1.3.3 All Access

A person whose userid has an all level of access is allowed to login to the switch via the serial port and via telnet.

### 8.1.3.4 No Access

A person whose userid has a none level of access is not allowed to log in to the switch at all.

### 8.1.4 Userid Password

The network administrator may also assign a password to a userid using the conf security login password command. (This command replaces the old oper password command.) When a user logged in with user privileges wants to modify the password, he or she must correctly enter the old password first before typing the new password. However, a user logged in with admin privileges can change any userid password without first entering the old local password, except their own. The maximum size is for a password is 512 characters. Any characters are allowed, except the colon (:) character.

# 8.1.5 Privilege Level for Unlisted Users

One of the login scenarios in Table 8.1 allows users who are not listed in the switch to login, provided that the SecurID server is accessible and provided that the SecurID server accepts the userid. The network administrator may assign a privilege level of user (access to only a certain subset of the AMI commands) or a privilege level of admin (access to all AMI commands) for these users via the conf security login upriv command.

# 8.2 IP Filtering

The IP filtering feature lets the network administrator limit access to the control port of the switch to prevent unauthorized access to the switch. The switch performs filtering on incoming IP packets by determining if there is a match between the packet's header source address and this table of authorized incoming IP addresses. If the addresses match, the packets are accepted, provided that they meet the requirements set up by the other IP filtering flags; otherwise, they are rejected. Statistics are kept of the number of rejected IP packets and about the last IP packet that was rejected.

### 8.2.1 Authorized IP Address Table

Using the conf security ipaccess accept command, the administrator can create an entry in a table of authorized IP addresses from which IP packets will be accepted. When the administrator creates an IP address entry in the table, he or she can apply a mask to specify a wildcard range of allowable addresses. For example, an IP address of 163.26.54.6 with a mask of 255.255.255.255 means only address 163.26.54.6 is accepted against that entry. However, an IP address of 163.26.54.6 with a mask of 255.255.0.0 means addresses 163.26.\*.\* are accepted against that entry. The table can contain a maximum of 32 entries.

The administrator can also delete entries from the table and display the contents of the table. The table will persist across a reboot. (See the *AMI Configuration Commands Reference Manual* for more information.)

### CAUTION



When the authorized IP address table is empty, all addresses are accepted. This is the default state. Therefore, it is recommended that at least one address be entered into the table. Otherwise, anyone may access the switch via the control port. The address you enter must be the address of the machine you are using. Otherwise, you will lock yourself out of the switch.

# 8.2.2 IP Filtering Flags

There are three IP filtering flags that can be configured to limit IP access in other ways: ssr, lsr, and all. These flags are set to allow or disallow IP packets using the conf security ipaccess ssr, lsr, and all commands. (See the AMI Configuration Commands Reference Manual for more information about these commands.)

### 8.2.2.1 Strict Source Routing Flag

If the ssr flag is set to allow, all incoming IP packets that are strict source routed are accepted, provided that they match an IP address in the table of authorized addresses. If the ssr flag is set to disallow, all incoming IP packets that are strict source routed are rejected, even if they match an IP address in the table of authorized addresses.

### 8.2.2.2 Loose Source Routing Flag

If the lsr flag is set to allow, all incoming IP packets that are loose source routed are accepted, provided that they match an IP address in the table of authorized addresses. If the lsr flag is set to disallow, all incoming IP packets that are loose source routed are rejected, even if they match an IP address in the table of authorized addresses.

### 8.2.2.3 All Flag

If the all flag is set to allow, all incoming IP packets are accepted, provided that they match an IP address in the table of authorized addresses. Allowing all is the default setting. If the all flag is set to disallow, all incoming IP packets are rejected, even if they match an IP address in the table of authorized addresses.

## 8.2.3 IP Access Statistics

The network administrator can display the total number of IP packets that have been filtered since the switch was rebooted using the statipaccess command. This command also displays information about the last IP packet that was dropped. This information includes the following:

- the reason that the last IP packet was dropped
- the system time at which the last IP packet was dropped
- · the name of the interface on which the last dropped IP packet was received
- the IP address contained in the source field of the header of the last IP packet that was dropped

# 8.3 NSAP Filtering

This feature provides a mechanism for filtering calls based on a combination of the calling (source) and called (destination) addresses, as well as the incoming and outgoing UNIs.

Each UNI may have one address filter for incoming call setups, and one for outgoing call setups. If a call setup is routed from incoming UNI A to outgoing UNI B, A's incoming call filter and B's outgoing call filter are applied. A call setup message must be accepted by both filters if both are present.



SPVCs and PVCs are not supported by NSAP filtering.

This feature also has a filter lookup mechanism and it provides statistics about calls that were filtered and information about the last call that was rejected by the filtering process.

# 8.3.1 Filters and Templates

Each filter is composed of an ordered set of templates. A template consists of the following:

- an incoming UNI (port/VPI)
- a source NSAP address/mask
- an outgoing UNI (port/VPI)
- a destination NSAP/Mask
- an action to either accept or reject the UNI or NSAP address

Masks may be wildcarded using a 0. Ports, VPIs, and NSAP addresses may be wildcarded using an asterisk (\*). If \* is specified for the port, then any port is accepted by the filter. If \* is specified for the VPI, then any VPI is accepted by the filter. If \* is specified for the NSAP address, or if 0 is specified for the mask, then any NSAP address is accepted by the filter.

Templates within a filter are applied in the order in which they appear in the filter, not by maximum prefix match. See the *AMI Configuration Commands Reference Manual* for more information about how to create, modify, and delete filters and templates.

# 8.3.2 NSAP Filtering Lookup

A filter lookup mechanism is provided which allows you to enter components of a call setup message to test whether a call setup attempt with the supplied addresses and ports would be accepted or rejected by a specific filter. The switch returns an answer of accepted or rejected and the index number of the template that accepted or rejected the information entered. If the information does not match any of the existing templates, then an answer of rejected and address unknown is given. When this feature is used, none of the statistics are incremented and no traps are sent.

# 8.3.3 NSAP Filtering Statistics

To display statistics for NSAP filtering, use the stat nsapfilter command. See the AMI Configuration Commands Reference Manual for more information about this command. NSAP filtering statistics and traps include:

- · counts of accepted calls
- calls rejected because of an explicit match with a reject template
- calls rejected because of no match with any template

A trap is sent when rejected calls occur more frequently than the user-specified limit.

# CHAPTER 9

# **Configuring Timing**

This chapter describes how to set up timing on a *ForeRunner* ATM Switch. Topics covered include:

- Section 9.1 Overview
- Section 9.2 Timing Modes
- Section 9.3 Switchclock
- Section 9.4 Port Level Timing
- Section 9.5 Timing Configuration Examples

### 9.1 Overview

For ATM networks to function well at full bandwidths, the equipment on the networks should be timed so that they are all transmitting and receiving data in a synchronized manner. Even though ATM is asynchronous, this only refers to the cell-level transmission that occurs in the ATM Layer. The physical layer underlying the cells is synchronous. Therefore, a time reference signal must be distributed to every element in the network to establish one cohesive entity. This time reference on *ForeRunner* ATM Switches is called a *switchclock*.

# 9.2 Timing Modes

There are two timing modes available on *ForeRunner* ATM Switches: TCM and switch. TCM mode means that all network modules that support distributed timing import their clock source from the Timing Control Module (TCM).



You must have a TCM installed in your switch in order to use this timing mode. See the *CEC-Plus User's Manual* for information about configuring timing with a TCM. All examples in this chapter assume that you do <u>not</u> have a TCM installed.

Switch mode means that all network modules that support distributed timing import their clock source from the designated switchclock port. This is the default mode for any switch that does not have a TCM installed. See Section 9.3 for more information about this mode.

# 9.3 Switchclock

If a TCM is not installed in the switch, all of the ports within a fabric use the switchclock as their timing reference. The switchclock can be any port that is able to recover a clock (i.e., the network module supports distributed timing). (To see if the network module supports distributed timing, use the command <code>conf module show</code> and look for <code>yes</code> under the <code>Timing</code> field.) In an ASX-1000, the port may be from another fabric in the same chassis.

When the switchclock is set to any of the ports within the same fabric, that port's clock is exported to all of the network modules within the same fabric. On an ASX-1000 only, the switchclock is also exported to the network modules in all of the other fabrics in this chassis. You can then configure each of the other fabrics to use that clock. In this way, all of the network modules in all of the fabrics will use the same timing source.

### 9.3.1 Failover of the Switchclock

You can configure a *primary* switchclock and a *secondary* switchclock. If the primary clock is valid, it is used. If the primary switchclock fails, the secondary clock is used. When the primary clock returns, it is used again as the clock source.

If both the primary and secondary values fail, the switch fabric uses the crystal of the first available timing network module as the switchclock, going from A to D. For example, if network module A supports distributed timing, then the crystal from A is used as the switchclock. As another example, if network module A is not installed, and network modules B and C do not support distributed timing, but network module D does support distributed timing, then the crystal from D is used as the switchclock.

Failover of the switchclock happens automatically when any of the following events occur:

- A carrier detect or carrier loss occurs on any port in this switch fabric.
- A network module is swapped in or out.
- A switch fabric is swapped in or out.

### 9.4 Port Level Timing

In addition to the switchclock, there is also a configurable parameter for each port called a *transmit* clock (*txclock*). Each port's txclock can be configured to use either the *network* clock or the *internal* clock.

- When the txclock is set to network, the clock that is recovered from the receive line of a port is used to drive the transmit line of that port.
- When the txclock is set to internal, the internal clock, or switchclock, is used to drive the transmit line of that port.



The Timing field under conf port sonet show displays N/A for all series of OC-12 network modules because they always use internal timing.

### 9.5 Timing Configuration Examples

This section provides examples of how to set the switchclock on various types of switches. For more information about the specific timing commands, see the *AMI Configuration Commands Reference Manual*.

If you have a TCM installed in your switch, see the *CEC-Plus User's Manual* for information about configuring timing with a TCM. All examples in this chapter assume that you do <u>not</u> have a TCM installed.

### 9.5.1 Configuring Timing on an ASX-200BX or ASX-200WG

This example assumes that you are going to use 1A1 as your primary clock and 1B1 as your secondary clock. Use the following ATM Management Interface (AMI) commands to set your primary and secondary switchclocks:

conf timing switchclock primary 1a1

conf timing switchclock secondary 1b1

### 9.5.2 Configuring Timing on an ASX-1000 (Single Timing Domain)

This example assumes that you are going to use 3A1 as your primary clock and 3B1 as your secondary clock. The only limitation on an ASX-1000 is that you must configure both the primary and secondary clocks to be on the same fabric. On all fabrics installed in the switch, use the following AMI commands to set your primary and secondary clocks as follows:

conf timing switchclock primary 3a1

conf timing switchclock secondary 3b1

# **Configuring Timing**

### 9.5.3 Configuring Timing on an ASX-1000 (Multiple Timing Domains)

This example assumes that you are going to use 1C1 as your primary clock and 1D1 as your secondary clock for fabrics 1 and 2 and that you are going to use 3A1 as your primary clock and 3B1 as your secondary clock for fabrics 3 and 4. The only limitation on an ASX-1000 is that you must configure both the primary and secondary export clocks to be on the same fabric.

1. On fabrics 1 and 2, use the following AMI commands to set your primary and secondary clocks as follows:

conf timing switchclock primary 1c1 conf timing switchclock secondary 1d1

2. On fabrics 3 and 4, use the following AMI commands to set your primary and secondary clocks as follows:

conf timing switchclock primary 3a1
conf timing switchclock secondary 3b1

### 9.5.4 Configuring Timing on a ForeRunnerLE 155

This example assumes that you are going to use 1D1 as your primary clock and 1D2 as your secondary clock. Use the following AMI commands to set your primary and secondary clocks as follows:

conf timing switchclock primary 1d1 conf timing switchclock secondary 1d2

The ports in interface groups A, B, C, and the Port Expansion Module (PEM) (interface D) use the crystal on the fabric as their timing source by default. Only the PEM can recover a clock from any of its ports and export it to the other interfaces. If the PEM is installed, the primary and secondary clocks be recovered from one of the PEM ports.

### Configuring Timing

The switch control software for the ForeRunner ATM switches includes an SNMP agent. The SNMP agent enables the remote monitoring and configuration of these switches.

### A.1 SNMP Indexing

There are two main SNMP indexing schemes used: software port indices and hardware port indices. Software port indices are single numbers starting at 0 for the first port, incrementing 8 ports per module on an ASX-200WG and on an ASX-200BX. For example, port A1 on an ASX-200WG or an ASX-200BX has a software port index of 0. Port C3 on an ASX-200WG, or an ASX-200BX has a software port index of 18, or 8 \* 2 + 2.

Hardware port indices are of the form {board}.{network module}.{port} or bnp notation. They start at 0.0.0 for the first port and increment across boards, network modules, and ports. For example, port C3 on an ASX-200WG or an ASX-200BX is hardware port 0.2.2.

Please refer to Table A.1 for a summary of the port numbering conventions used in *ForeRunner* switches and related SNMP indexing format.

**Table A.1 -** ASX-200WG/ASX-200BX Port Numbering

Port Name	Software Port Number	Board-Netmod- Port Index	Port Name	Software Port Number	Board-Netmod- Port Index
A1	0	0.0.0	C1	16	0.2.0
A2	1	0.0.1	C2	17	0.2.1
A3	2	0.0.2	C3	18	0.2.2
A4	3	0.0.3	C4	19	0.2.3
A5	4	0.0.4	C5	20	0.2.4
A6	5	0.0.5	C6	21	0.2.5
B1	8	0.1.0	D1	24	0.3.0
B2	9	0.1.1	D2	25	0.3.1
В3	10	0.1.2	D3	26	0.3.2
B4	11	0.1.3	D4	27	0.3.3
B5	12	0.1.4	D5	28	0.3.4
B6	13	0.1.5	D6	29	0.3.5
			CTL	56	0.7.0

### A.2 SNMP Traps

SNMP traps are used to update the state of the network automatically to remote network management hosts. The SNMP agent on the switch supports several SNMP traps.

The traps generated by the switch's SNMP agent can be sent to as many destinations as needed. These destinations are configurable via the ATM Management Interface (AMI). Each destination must be an IP address of a network management host. The network management host specified for a trap destination can be any host with which the switch has connectivity. This means that the host does not have to be a directly connected ATM host. It can be on any attached network. The following table describes the supported traps.

**Table A.2 - SNMP Traps Supported on the** *ForeRunner Switches* 

Trap Number	Trap Name	Description
0	asxSwLinkDown	An asxSwLinkDown trap signifies that the sending protocol entity recognizes a failure in one of the ATM Switch links that is connected to another switch.
1	asxSwLinkUp	An asxSwLinkUp trap signifies that the sending protocol entity recognizes that one of the ATM Switch links that is connected to another switch has come up.
2	asxHostLinkDown	An asxHostLinkDown trap signifies that the sending protocol entity recognizes a failure in one of ATM Switch links that is connected to a host.
3	asxHostLinkUp	An asxHostLinkUp trap signifies that the sending protocol entity recognizes that one of the ATM Switch links that is connected to a host has come up.
4	asxNetModuleDown	An asxNetModuleDown trap signifies that the sending protocol entity recognizes a failure in one of the ATM Switch network modules, that is identified by the board and the module numbers. This is probably caused by a hot-swap of a network module.
5	asxNetModuleUp	An asxNetModuleUp trap signifies that the sending protocol entity recognizes a new operational ATM Switch network module, that is identified by the board and the module numbers. This is probably caused by a hot-swap of a network module.

**Table A.2 - SNMP** Traps Supported on the *ForeRunner* Switches (Continued)

Trap Number	Trap Name	Description
6	asxPsInputDown	This trap alerts that one ATM switch power supply failed due to failure in the input voltage. The power supply that failed is identified by the power supply index. Note that an input voltage may be out of specification and may not cause a power supply failure if high loads are not applied.
7	asxPsInputUp	This trap alerts that one ATM switch power supply that had an AC input failure is up. The power supply that is back up is identified by the power supply index.
9	asxPsOutputDown	This trap alerts that one ATM switch power supply output or the power supply was physically removed. The power supply that failed is identified by the power supply index.
10	asxPsOutputUp	This trap alerts that one ATM switch power supply that had an output failure or was removed is now up. The power supply that is back up is identified by the power supply index.
22	asxFanBankDown	This trap alerts that one ATM switch fan bank failed. The fan bank that failed is identified by the fan bank index.
23	asxFanBankUp	This trap alerts that one ATM switch fan bank is up. The fan bank that is back up is identified by the fan bank index.
28	asxLinkDown	This trap alerts that the link that is identified by {hwPortBoard, hwPortModule, hwPortNumber} was configured up but lost its carrier (or the framing bit) and is currently down.
29	asxLinkUp	This trap alerts that the link that is identified by {hwPortBoard, hwPortModule, hwPortNumber} is back up.
30	asxSpansDown	This trap alerts that the SPANS signalling on the link that is identified by the sigPathPort and sigPathVPI failed.

**Table A.2 - SNMP Traps Supported on the** ForeRunner Switches (Continued)

Trap Number	Trap Name	Description
31	asxSpansUp	This trap alerts that the SPANS signalling on the link that is identified by the sigPathPort and sigPathVPI is up.
32	asxTempSensorOverTemp	This trap alerts that one of the temperature sensors indicates over temperature. The temperature sensor is identified by the temperature sensor index.
33	asxTempSensorRegularTemp	This trap alerts that one of the temperature sensors indicates regular temperature. The temperature sensor is identified by the temperature sensor index.
34	asxFabricTemperature OverTemp	This trap alerts that one of the temperature sensors indicates over temperature. The temperature sensor is identified by the temperature sensor index.
35	asxFabricTemperature RegularTemp	This trap alerts that one of the temperature sensors indicates regular temperature. The temperature sensor is identified by the temperature sensor index.
36	asxSonetLOSon	This trap indicates that the specified SONET port is experiencing Loss Of Signal. Bellcore Document TA-NWT-000253 Section 6.3.1.1.1 states that, "A SONET NE shall declare a LOS failure when the LOS defect persists for 2.5 ( $\pm$ .5) seconds, or when a LOS defect is present and the criteria for LOF failure declaration have been met."
37	asxSonetLOSoff	This trap indicates that the LOS condition identified by trap asxSonetLOSon has been cleared.
38	asxSonetPathLabelOn	This trap indicates that the specified SONET port is receiving and errored C2 Path Label byte. Reference Bellcore Document TA-NWT-000253 Section 3.3.2.3 and 6.3.1.1.8 the Path Label (C2) byte should have the value 0x13.
39	asxSonetPathLabelOff	This trap indicates that the Errored Path Label (C2) byte error condition signalled by the asxSonetPath-LabelOn trap has been cleared.

**Table A.2 - SNMP** Traps Supported on the *ForeRunner* Switches (Continued)

Trap Number	Trap Name	Description
40	asxSonetLineAISon	This trap indicates that the specified SONET port is receiving a Line level Alarm Indication Signal from the far-end equipment.
41	asxSonetLineAISoff	This trap indicates that the Line AIS error condition signalled by the asxSonetLineAISon trap has been cleared.
46	asxDS3PLCPYellowDetected	This trap indicates that the specified DS3 port has detected incoming Yellow Alarm.
47	asxDS3PLCPYellowCleared	This trap indicates that the specified DS3 port has detected clearance of incoming Yellow Alarm.
48	asxDS3PLCPLOFDetected	This trap indicates that the specified DS3 port has detected incoming LOF Alarm.
49	asxDS3PLCPLOFCleared	This trap indicates that the specified DS3 port has detected clearance of incoming LOF Alarm.
50	asxDS3LOFDetected	This trap indicates that Loss Of Frame(LOF) is detected on the incoming signal.
51	asxDS3LOFCleared	This trap indicates that Loss Of Frame is cleared on the incoming signal.
52	asxDS3AISDetected	This trap indicates that AIS Alarm is detected on the incoming signal.
53	asxDS3AISCleared	This trap indicates that AIS Alarm is cleared on the incoming signal.
60	asxDS1PLCPYellowDetected	This trap indicates that the specified DS1 port has detected an incoming Yellow Alarm.
61	asxDS1PLCPYellowCleared	This trap indicates that the specified DS1 port has detected clearance of an incoming Yellow Alarm.
62	asxDS1PLCPLOFDetected	This trap indicates that the specified DS1 port has detected an incoming LOF Alarm.
63	asxDS1PLCPLOFCleared	This trap indicates that the specified DS1 port has detected clearance of an incoming LOF Alarm.
64	asxDS1YellowDetected	This trap indicates that Yellow Alarm is detected on the incoming signal.

**Table A.2 - SNMP Traps Supported on the** ForeRunner Switches (Continued)

Trap Number	Trap Name	Description
65	asxDS1YellowCleared	This trap indicates that Yellow Alarm is cleared on the incoming signal.
66	asxDS1AISDetected	This trap indicates that AIS Alarm is detected on the incoming signal.
67	asxDS1AISCleared	This trap indicates that AIS Alarm is cleared on the incoming signal.
68	asxDS1LOSDetected	This trap indicates that LOS Alarm is detected on the incoming signal.
69	asxDS1LOSCleared	This trap indicates that LOS Alarm is cleared on the incoming signal.
70	asxDS1LOFDetected	This trap indicates that LOF Alarm is detected on the incoming signal.
71	asxDS1LOFCleared	This trap indicates that LOF Alarm is cleared on the incoming signal.
74	asxDS3FERFDetected	This trap indicates that FERF Alarm is detected on the incoming signal.
75	asxDS3FERFCleared	This trap indicates that FERF Alarm is cleared on the incoming signal.
78	asxE3YellowDetected	This trap indicates that the Yellow Alarm is being detected on the incoming signal.
79	asxE3YellowCleared	This trap indicates that Yellow alarm has cleared on the incoming signal.
80	asxE3OOFDetected	This trap indicates that Out Of Frame (OOF) is detected on the incoming signal.
81	asxE3OOFCleared	This trap indicates that Loss Of Frame is cleared on the incoming signal.
82	asxE3AtmLCDDetected	This trap indicates that the specified E3 port is experiencing Loss of Cell Delineation (LCD). An LCD failure is declared when the LCD defect persists for a period of 2.5 +/- 0.5 seconds.

**Table A.2 - SNMP** Traps Supported on the *ForeRunner* Switches (Continued)

Trap Number	Trap Name	Description
83	asxE3AtmLCDCleared	This trap indicates that the LCD failure identified by trap asxE3AtmLCDDetected has been cleared. An LCD failure is cleared when the LCD defect is absent for 10 +/- 0.5 seconds.
86	asxE3PLCPYellowDetected	This trap indicates that the specified E3 port has detected incoming PLCP Yellow Alarm.
87	asxE3PLCPYellowCleared	This trap indicates that the specified E3 port has detected clearance of incoming PLCP Yellow Alarm.
90	asxE1YellowDetected	This trap indicates that the Yellow Alarm is being detected on the incoming signal.
91	asxE1YellowCleared	This trap indicates that Yellow alarm has cleared on the incoming signal.
92	asxE1LOFDetected	This trap indicates that LOF is being detected on the incoming signal.
93	asxE1LOFCleared	This trap indicates that LOF is cleared on the incoming signal.
96	asxE1PLCPYellowDetected	This trap indicates that the specified E1 port has detected incoming PLCP Yellow Alarm.
97	asxE1PLCPYellowCleared	This trap indicates that the specified E1 port has detected clearance of incoming PLCP Yellow Alarm.
98	asxE1PLCPLOFDetected	This trap indicates that the specified E1 port has detected incoming PLCP LOF Alarm.
99	asxE1PLCPLOFCleared	This trap indicates that incoming PLCP LOF alarm has been cleared on the specified E1 port.
100	asxE1LOSDetected	This trap indicates that the specified E1 port has detected incoming LOS Alarm.
101	asxE1LOSCleared	This trap indicates that incoming LOS alarm has been cleared on the specified E1 port.
102	asxE1AISDetected	This trap indicates that the specified E1 port has detected incoming AIS Alarm.
103	asxE1AISCleared	This trap indicates that incoming AIS alarm has been cleared on the specified E1 port.

**Table A.2 - SNMP Traps Supported on the** ForeRunner Switches (Continued)

Trap Number	Trap Name	Description
104	asxE3AISDetected	This trap indicates that the specified E3 port has detected incoming AIS Alarm.
105	asxE3AISCleared	This trap indicates that incoming AIS alarm has been cleared on the specified E3 port.
106	asxE3LOSDetected	This trap indicates that the specified E3 port has detected incoming LOS Alarm.
107	asxE3LOSCleared	This trap indicates that incoming LOS alarm has been cleared on the specified E3 port.
108	asxE3PLCPLOFDetected	This trap indicates that the specified E3 port has detected incoming PLCP LOF Alarm.
109	asxE3PLCPLOFCleared	This trap indicates that incoming PLCP LOF alarm has been cleared on the specified E3 port.
112	asxJ2YellowDetected	This trap indicates that Yellow Alarm is detected on the incoming signal.
113	asxJ2YellowCleared	This trap indicates that Yellow Alarm is cleared on the incoming signal.
114	asxJ2AISDetected	This trap indicates that AIS Alarm is detected on the incoming signal.
115	asxJ2AISCleared	This trap indicates that AIS Alarm is cleared on the incoming signal.
116	asxJ2LOSDetected	This trap indicates that LOS Alarm is detected on the incoming signal.
117	asxJ2LOSCleared	This trap indicates that LOS Alarm is cleared on the incoming signal.
118	asxJ2LOFDetected	This trap indicates that LOF Alarm is detected on the incoming signal.
119	asxJ2LOFCleared	This trap indicates that LOF Alarm is cleared on the incoming signal.
120	asxDS3LOSDetected	This trap indicates that the specified DS3 port has detected incoming LOS Alarm.

**Table A.2 - SNMP** Traps Supported on the *ForeRunner* Switches (Continued)

Trap Number	Trap Name	Description
121	asxDS3LOSCleared	This trap indicates that the incoming LOS Alarm has been cleared on the specified DS3 port.
130	asxSonetLOFDetected	This trap indicates that the specified SONET port is experiencing Loss Of Frame (LOF) failure. An LOF failure is declared when the LOF defect persists for a period of 2.5 +/- 0.5 seconds, except when an LOS defect or failure is present.
131	asxSonetLOFCleared	This trap indicates that the LOF failure identified by trap asxSonetLOFDetected has been cleared. The LOF failure is cleared when the LOS failure is declared, or when the LOF defect is absent for 10 +/- 0.5 seconds.
132	asxSonetLineRDIDetected	This trap indicates that the specified SONET port is experiencing Line Remote Defect Indication (LRDI). A Line RDI failure is declared when the incoming Line RDI defects lasts for 2.5 +/- 0.5 seconds.
133	asxSonetLineRDICleared	This trap indicates that the Line RDI failure identified by trap asxSonetLineRDIDetected has been cleared. The Line RDI failure is cleared when no Line RDI defects are detected for 10 +/- 0.5 seconds.
134	asxSonetPathAISDetected	This trap indicates that the specified SONET port is experiencing Path Alarm Indication Signal (PAIS). A Path AIS failure is declared when the Path AIS defect persists for 2.5 +/- 0.5 seconds.
135	asxSonetPathAISCleared	This trap indicates that the Path AIS failure identified by trap asxSonetPathAISDetected has been cleared. A PAIS failure is cleared when the PAIS defect is absent for 10 +/- 0.5 seconds.
136	asxSonetPathLOPDetected	This trap indicates that the specified SONET port is experiencing Loss Of Pointer (LOP). A LOP failure is declared when the LOP defect persists for a period of 2.5 +/- 0.5 seconds.
137	asxSonetPathLOPCleared	This trap indicates that the LOP failure identified by trap asxSonetLOPDetected has been cleared. A LOP failure is cleared when the LOP defect is absent for 10 +/- 0.5 seconds.

**Table A.2 - SNMP Traps Supported on the** ForeRunner Switches (Continued)

Trap Number	Trap Name	Description
138	asxSonetPathUNEQDetected	This trap indicates that the specified SONET port is experiencing unequipped (UNEQ). A UNEQ failure is declared when the UNEQ defect persists for a period of 2.5 +/- 0.5 seconds.
139	asxSonetPathUNEQCleared	This trap indicates that the UNEQ failure identified by trap asxSonetUNEQDetected has been cleared. A UNEQ failure is cleared when the UNEQ defect is absent for 10 +/- 0.5 seconds.
140	asxSonetPathRDIDetected	This trap indicates that the specified SONET port is experiencing Path Remote Defect Indication (PRDI). A Path RDI failure is declared when the incoming Path RDI defects lasts for 2.5 +/- 0.5 seconds.
141	asxSonetPathRDICleared	This trap indicates that the Path RDI failure identified by trap asxSonetPathRDIDetected has been cleared. The Path RDI failure is cleared when no Path RDI defects are detected for 10 +/- 0.5 seconds.
142	asxSonetAtmLCDDetected	This trap indicates that the specified SONET port is experiencing Loss of Cell Delineation (LCD). A LCD failure is declared when the LCD defect persists for a period of 2.5 +/- 0.5 seconds.
143	asxSonetAtmLCDCleared	This trap indicates that the LCD failure identified by trap asxSonetAtmLCDDetected has been cleared. A LCD failure is cleared when the LCD defect is absent for 10 +/- 0.5 seconds.
160	asxDS3IdleDetected	This trap indicates that an Idle Maintenance Signal (IDLE) is detected on the incoming signal.
161	asxDS3IdleCleared	This trap indicates that an Idle Maintenance Signal (IDLE) is cleared on the incoming signal.
162	asxDS3AtmLCDDetected	This trap indicates that the specified DS3 port is experiencing Loss of Cell Delineation (LCD). An LCD failure is declared when the LCD defect persists for a period of 2.5 +/- 0.5 seconds.

**Table A.2 - SNMP** Traps Supported on the *ForeRunner* Switches (Continued)

Trap Number	Trap Name	Description
163	asxDS3AtmLCDCleared	This trap indicates that the LCD failure identified by trap asxDS3AtmLCDDetected has been cleared. An LCD failure is cleared when the LCD defect is absent for $10 +/-0.5$ seconds.
176	asxDS1PRBSDetected	This trap indicates that PRBS pattern is detected on the incoming signal.
177	asxDS1PRBSCleared	This trap indicates that PRBS pattern is cleared on the incoming signal.
178	asxDS1AtmLCDDetected	This trap indicates that the specified DS1 port is experiencing Loss of Cell Delineation (LCD). An LCD failure is declared when the LCD defect persists for a period of 2.5 +/- 0.5 seconds.
179	asxDS1AtmLCDCleared	This trap indicates that the LCD failure identified by trap asxDS1AtmLCDDetected has been cleared. An LCD failure is cleared when the LCD defect is absent for 10 +/- 0.5 seconds.
192	asxE3TrailChangeDetected	This trap indicates that a change in the trail trace message was detected on the incoming signal.
208	asxE1AtmLCDDetected	This trap indicates that the specified E1 port is experiencing Loss of Cell Delineation (LCD). An LCD failure is declared when the LCD defect persists for a period of 2.5 +/- 0.5 seconds.
209	asxE1AtmLCDCleared	This trap indicates that the LCD failure identified by trap asxE1AtmLCDDetected has been cleared. An LCD failure is cleared when the LCD defect is absent for 10 +/- 0.5 seconds.
224	asxJ2RLOCDetected	This trap indicates that Receive Loss of Clock (RLOC) is detected on the incoming signal.
225	asxJ2RLOCCleared	This trap indicates that Receive Loss of Clock (RLOC) is cleared on the incoming signal.
226	asxJ2HBERDetected	This trap indicates that High Bit Error Rate (HBER) is detected on the incoming signal.

**Table A.2 - SNMP Traps Supported on the** ForeRunner Switches (Continued)

Trap Number	Trap Name	Description
227	asxJ2HBERCleared	This trap indicates that High Bit Error Rate (HBER) is cleared on the incoming signal.
228	asxJ2PAISDetected	This trap indicates that Payload Alarm Indication Signal (PAIS) is detected on the incoming signal.
229	asxJ2PAISCleared	This trap indicates that Payload Alarm Indication Signal (PAIS) is cleared on the incoming signal.
230	asxJ2AtmLCDDetected	This trap indicates that the specified J2 port is experiencing Loss of Cell Delineation (LCD). An LCD failure is declared when the LCD defect persists for a period of 2.5 +/- 0.5 seconds.
231	asxJ2AtmLCDCleared	This trap indicates that the LCD failure identified by trap asxJ2AtmLCDDetected has been cleared. An LCD failure is cleared when the LCD defect is absent for 10 +/- 0.5 seconds.
232	asxJ2TLOCDetected	This trap indicates that Transmit Loss of Clock (TLOC) is detected.
233	asxJ2TLOCCleared	This trap indicates that Transmit Loss of Clock (TLOC) is cleared.
1024	asxOutputQueueCongested	This trap indicates that the output queue for the given priority has exceeded its dedicated length, and has begun overflowing into the shared buffer space on the network module.
1025	asxOutputQueueCellLoss	This trap indicates that the output queue for the given priority has overflowed and cells have been dropped.
1026	as x Extended Mode Violation	This trap indicates that a series A or B network module was inserted into a switch board running in extended mode.
1027	asxNonextendedMode- Warning	This trap indicates that a series C or greater network module was inserted into a switch board running in non-extended mode.

**Table A.2 - SNMP** Traps Supported on the *ForeRunner* Switches (Continued)

Trap Number	Trap Name	Description					
1028	q2931AVRejectTrap	This trap is generated whenever any UNI3.x with AddressValidation enabled rejects a Setup Request call more than q2931AVRejectTrapThreshold times in any given q2931AVRejectTrapPeriod.					
1029	crConfMemoryOflow	This trap is generated when the allocated call record memory (as indicated by crMemoryAllocated) is exceeded.					
1030	crXfrPrimaryXfrFailed	This trap is generated when the call record transfer to the primary host (as indicated by crXfrPrimaryUrl) fails.					
1031	crXfrSecondaryXfrFailed	This trap is generated when the call record transfer to the secondary host (as indicated by crXfrSecond- aryUrl) fails.					
1032	crConfMemAllocFail	This trap is generated when Callrecord functionality is unable to allocate memory as specified by crMemory-Allocated. This can happen when the crConfAdmin-Status changes state from "off" or when the switch reboots when Callrecords is configured "on".					
1033	crGeneralFailure	This trap is generated when any of the callrecord related functionality fails for any reason. One example would be when the Callrecord Module fails to schedule an interval timer.					
1034	asxDualScpSyncFailure	This trap indicates that automatic CDB synchronization is disabled due to failures.					
1035	asxDualScpSwitchOver	This trap indicates that the backup SCP has taken control of the switch.					
1036	asxDualScpHotSwap	This trap indicates that an SCP hotswap insertion or removal has occurred.					
1037	asxVPAISDetected	This trap indicates that the Alarm Indication Signal (AIS) is detected on the incoming (terminating) virtual path. This trap is generated once when the virtual path is declared to be in the active AIS state.					

**Table A.2 - SNMP Traps Supported on the** ForeRunner Switches (Continued)

Trap Number	Trap Name	Description				
1038	asxVPAISCleared	This trap indicates that the Alarm Indication Signal (AIS) has been removed from the incoming (terminating) virtual path. This trap is generated once when the virtual path is declared to be in the inactive AIS state.				
1039	asxVPRDIDetected	This trap indicates that the Remote Defect Indication (RDI) is detected on the incoming (terminating) virtual path. This trap is generated once when the virtual path is declared to be in the active RDI state.				
1040	asxVPRDICleared	This trap indicates that the Remote Defect Indication (RDI) has been removed from the incoming (terminating) virtual path. This trap is generated once when the virtual path is declared to be in the inactive RDI state.				
1041	asxNonextendedModeViolation	This trap indicates that a Series D network module was inserted into a switch board running in non-extended mode. Multicast will not work on the Series D module without removing all Series B modules and rebooting the switch.				
1042	asxUnsupportedNetwork- Module	This trap indicates that a unsupported network module was inserted into a switch.				
1049	asxIpFilterViolation	This trap occurs when an incoming IP packet is unauthorized to enter the switch control port and has been dropped.				
1053	q2931AFRejectKnown	This trap is generated whenever any q2931 UNI with Address Filtering enabled rejects a Setup request because the request matched a template with the action "reject." The variables sent in the trap identify the source and destination UNI for the call.				
1054	q2931AFRejectUnknown	This trap is generated whenever any q2931 UNI with Address Filtering enabled rejects a Setup request because the address matched no template. The variables sent in the trap identify the source and destination UNI for the call.				
1061	q2931CreationFailure	This trap is generated whenever a switch fails to create a UNI. This is most likely due to a resource limitation on the switch.				

**Table A.2 - SNMP** Traps Supported on the *ForeRunner* Switches (Continued)

Trap Number	Trap Name	Description					
1068	asxPsCurrentDown	This trap alerts that one ATM switch power supply had a current failure. The power supply that failed is identified by the power supply index.					
1069	asxPsCurrentUp	This trap alerts that one ATM switch power supply that had a current failure is now up. The power supply that is back up is identified by the power supply index.					
1070	asxPs5VoltDown	This trap alerts that one ATM switch power supply had a +5V failure. The power supply that failed is identified by the power supply index.					
1071	asxPs5VoltUp	This trap alerts that one ATM switch power supply that had a +5V failure is now up. The power supply that is back up is identified by the power supply index.					
1072	asxSwitchLoginDetected	An asxSwitchLoginDetected trap signifies that a user logged in on the switch.					
1073	asxSwitchLoginFailed	An asxSwitchLoginFailed trap signifies that a user's attempt to log in on the switch failed.					
1074	pnniTdbGuardbandResrv- Fail	This trap is generated when the guardband memory reserve for any of the PNNI TDB (Topology Database) related functionality like creation, modification, and deletion on objects like node, PTSE, flags, internal prefixes, external prefixes, etc. fails. The switch is low on memory.					
1075	pnniTdbInconsistentState	This trap is generated when the PNNI TDB (Topology Database) is in an unrecoverable error condition due to MALLOC and other TDB related failures. When this happens, the associated logical node is shut down and this trap is sent. The switch has to be rebooted to bring this PNNI logical node up again.					

**Table A.2 - SNMP Traps Supported on the** ForeRunner Switches (Continued)

Trap Number	Trap Name	Description
2000	frf8PVCStatus	This trap indicates when an interworking PVC has experienced an alarmed condition, either on the ATM network side or Frame Relay side. It is also generated when the PVC alarmed condition is cleared. It carries the operational status of the PVC by the frf8ConnOperStatus, as well as the reason why exiting, entering, or changing the alarmed state frf8ConnPVCAlarmReason.
		If the interworking PVC changes its status by an AdminStatus command (active/inactive/testing), causing the frf8ConnOperStatus to change (up/down), this trap is generated with frf8ConnPVCAlarmReason set to no defect. This trap gives the support to manage defects that might occur in either the Frame Relay network or the ATM network, even if the PVC Operational Status is already down.
2001	ifLinkDown	This trap indicates that the SNMP agent has detected that the ifOperStatus of this interface is about to transition into the down state.
2002	ifLinkUp	This trap indicates that the SNMP agent has detected that the ifOperStatus of this interface has transitioned out of the down state.
2003	framNakMsg	This trap is sent when a message from the SCP to the FramePlus network module does not succeed. The request that has failed is encoded within the trap as the messageType, and the reason for failure is encoded within the errorCode. See Table A.3 for the message types and the message requests. See Table A.4 for the error codes and their meanings.

Table A.3 defines the message types and the message requests for Trap 2003.

**Table A.3 - Message Type Encodings for Trap 2003** 

MessageType	Message Request		
0	Configure network module		
1	Configure port		
2	Send egress Fdl		
4	Network module reset		
5	Port reset		
6	Configure a channel/service		
7	Request channel/service configuration		
8	Request port configuration		
9	Request connection configuration		
10	Channel/service reset		
17	Configure LEDs		
28	Request application revision		
29	Create a FUNI service		
30	Create a Frame Relay service		
31 Request to change service admin status 32 Request to enable/disable service statistics 33 Request to enable/disable network module statistics			
		34	Request to enable/disable network module admin status
		35	Request to enable/disable port admin status
36 Delete a service			
37 Create a Frame Relay connection			
38	Create a FUNI connection		
40	Request to enable/disable connection admin status		
41 Request to enable/disable service egress rate enforcement			

Message TypeMessage Request42Request to enable/disable connection ingress rate enforcement43Delete a connection47Configure port to be timing source

**Table A.3 - Message Type Encodings for Trap 2003 (Continued)** 

Table A.4 defines the error codes for Trap 2003 and their meanings.

**Table A.4 -** Error Codes for Trap 2003

Error Code	Error Code Meaning		
1	Invalid or out of range parameter		
2	System disabled		
3	Memory allocation failure		
4	Timer allocation failure		
5	Circuits on link exceeded		
6	Invalid CIR configuration		
7	No CCB for link		
8	Unknown link		
9	Unknown DLCI		
0A	Link number out of range		
0B	DLCI number out of range		
0C	Link already exists		
0D	Circuit on link already exists		
0E	Link contains circuits		
0F	Out of memory		
10	Out of timers		
11	Out of timers		
12	Length error		
13	Duplicate connection ID		

**Table A.4 -** Error Codes for Trap 2003 (Continued)

Error Code	Error Code Meaning			
14	Invalid parameter			
15	Duplicate DLCI			
16	Memory allocation failure			
17	Service not created			
18	Connection does not exist			
19	Invalid Service			
1A	Connection ID unavailable			
1B	Invalid Connection ID			
1C	Unexpected message			
1D	Object already configured			
21	Operation failed			
22	Physical channel configuration failed			
23	Physical connection configuration failed			
24	Physical channel teardown failed			
25	Physical connection teardown failed			
26	Port configuration failed			
27	EPD/PPD configuration failed			
28	Unknown SDPM message received			
29	Netmod not configured			
2A	SDPM message size error			
2B	Port not configured			
2C	Allocation of buffer failed			

### A.2.1 Adding SNMP Trap Destinations

To create one or more SNMP trap destinations on a *ForeRunner* switch, log in to AMI and open a session on the switch. Enter the following parameters:

```
configuration snmp trap destinations new <ipaddress>
```

The <ipaddress> variable indicates the IP address of the SNMP trap destination that is to be created. Repeat this for as many SNMP trap destinations as needed. Traps are active as soon as they are set.

### A.2.2 Displaying SNMP Trap Destinations

To list all of the current SNMP trap destinations, log in to AMI and open a session on the switch. The SNMP traps supported by this switch are detailed in the FORE-Switch-MIB. Enter the following parameters:

configuration snmp trap destinations show

The switch responds with a list similar to the following:

Trap	Destination
1	192.88.243.18
2	198.29.16.14
3	198.29.16.18

If no trap destinations have been configured, then the following is displayed:

```
myswitch::configuration snmp trap destinations> show
No trap information is available
```

### A.2.3 Removing SNMP Trap Destinations

To delete one or more SNMP trap destinations for a *ForeRunner* switch, log in to AMI and open a session on the switch. Prior to deleting any trap that may need to be recreated later, as a precaution, a recommended practice is to list all trap destinations using AMI and either copy the screen or write down the destinations. To delete a trap, enter the following parameters:

configuration snmp trap destinations delete <trap>

The *<trap>* variable indicates the index number of the SNMP trap destination that is to be removed. Repeat this for as many SNMP trap destinations as needed.



For more information about the SNMP trap destination commands, see Part 2 of the *AMI Configuration Commands Reference Manual*.

### APPENDIX B

## **Configuring Circuit Emulation Services**

FORE Systems' Circuit Emulation Services (CES) Network Modules (NMCE-6/DS1A and NMCE-6/E1A) provide adaptation from time-division multiplexed (TDM) equipment (i.e., PBXs, WAN multiplexers, channel banks, video codecs, etc.) and traffic to ATM. Both modules provide structured and unstructured services, with a maximum of 127 connections supported on each module.

- All six ports of the DS1 CES Network Module may support fractional DS1 services (n x 56 Kbps/n x 64 Kbps) where 1 to 24 contiguous or non-contiguous DS0 channels are mapped to a single ATM VCC not to exceed 127 total connections.
- All six ports of the E1 CES Network Module may support fractional E1 services (n x 56 Kbps/n x 64 Kbps) where 1 to 31 contiguous or non-contiguous DS0 channels are mapped to a single ATM VCC not to exceed 127 total connections.

Structured services provide digital access and cross-connect system (DACS) connectivity where n x 64 Kbps and n x 56 Kbps digital signal level zero (DS0) channels are adapted to ATM cells and mapped to unique ATM virtual connections (VCCs).

Unstructured services provide support and maintenance of a single full bandwidth 1.544 Mbps (DS1) or 2.048 Mbps (E1) clear channel across a single ATM virtual connection.

Configurations of both the DS1 and the E1 version of the CES network module are detailed in this chapter. This chapter assumes the proper configuration of general switch parameters and of CES specific parameters, such as timing distribution/recovery, additional CES alarms, CES ports, etc.

### **B.1 Configuring CES Connections**

The ces commands let you create and delete CES connections, as well as display the status of existing connections. You can display the list of available subcommands by typing? at the ces level.

### **B.1.1 Creating a New CES Connection**

To create a new CES connection, you must set several parameters. Enter the following to create a new CES connection:

```
myswitch::configuration ces> new <port> <timeslots>
```

The CES new command can also be used as shown below. When the following parameters are used, by default, an appropriate entry is made in the UPC table and a bidirectional PVC is created with the proper UPC index.

```
or: new <port> <timeslots> -oport <oport> -ots <ots>
or: new <port> <timeslots> -oport <oport> -ovpi <ovpi> -ovci <ovci>
```

The following advanced options can be used when creating CES connections:

```
advanced options:
[-srts (on|off)] [-fupc <index>] [-bupc <index>]
[-cas (basic|cas)] [-partialfill <partialfill>] [-reassCDVT <cdvt>]
[-bufSize <bufSize>] [-integ <integ>]
```



SRTS is only available on unstructured connections, which are created by specifying all for the <timeslots> parameter.

The **-cas** and **-partialfill** options are not applicable to unstructured mode.

Structured mode is selected by indicating the exact timeslots to be used. For example, timeslots 1, 2, and 3 would be entered as 1-3, timeslots 2, 4, and 6 would be entered as 2:4:6, and combinations such as 1-4:9-11:12 are allowed.

The parameters for new are defined as follows:

Parameter	Description
port	The port on which the CES connection is to be created.
timeslots	Indicates which timeslots (1-24 for DS1, 1-31 for E1) are being configured for a particular PVC. all indicates unstructured service. The time slot assignments may be either contiguous or non-contiguous DS0s.
oport	The output port of the CES connection, which can be a CES port or an ATM port.
ovpi	The output Virtual Path Identifier (VPI) of the CES connection when the output port is not a CES port.
ovci	The output Virtual Channel Identifier (VCI) of the CES connection when the output port is not a CES port.
ots	The output timeslots of the CES connection when the output port is a CES port.
srts	Indicates whether Synchronous Residual Time Stamp (SRTS) clock recovery is to be enabled on this connection. on indicates that SRTS is enabled, off indicates that SRTS is disabled. The default is off.
-fupc <index></index>	The UPC contract type to be used in the ingress direction of the connection. (See Part 2 of the AMI Configuration Commands Reference Manual for more information about UPC contracts.)
-bupc <index></index>	The UPC contract type to be used in the egress direction of the connection. (See Part 2 of the AMI Configuration Commands Reference Manual for more information about UPC contracts.)
cas	Indicates whether Channel Associated Signalling (CAS) is to be used on the connection. basic indicates that CAS will not be used, cas indicates that CAS will be used. The default is basic.
partialfill	Indicates how many of the available 47 payload bytes in each cell are used before they are deemed "full" and ready for transmission across the ATM network (i.e., how much of the ATM cell contains data and how much is filler). The range for this parameter is 5 to 47. The default value is 47, for 47 bytes of data. partialfill is used to minimize network transmission latency and is useful especially with time-sensitive, robbed-bit signalling sources.

### Configuring Circuit Emulation Services

Parameter	Description
-reassCDVT <cdvt></cdvt>	The Cell Delay Variation Tolerance for cells being received by the segmentation and reassembly (SAR) engine. The range for this parameter is 100 to 24000 (in $\mu$ s), and the default is 2000 (i.e., 2 ms).
bufSize	The amount of reassembly buffer space allocated for the connection. The default is 512 bytes per timeslot.
integ	The amount of time allocated to re-establish the connection before, while, or after the call is established, or in the case of interruption. The default is 2500 ms.

### **B.1.2 Displaying CES Connections**

To display the current CES connections, enter the following:

myswitch:	:configuration	ces>	show
-----------	----------------	------	------

CES			Input				Output				
	Service	State	Port	Timeslots	VPI	VCI	Type	Port	TimeSlots	VPI	VCI
	24	down	1A1	1	0	129	-	-	_	-	-
	31	down	1A1	2-3	0	130	spvc	1D3	_	0	32
	32	down	1A1	4-5	0	131	pvc	1D4	_	0	150
	33	down	1A1	6-7	0	132	spvc	1D3	_	0	35

The fields in this display are defined as follows:

Field	Description
CES Service	The identification number (assigned by the switch) of this CES connection.
State	Indicates whether the CES connection is enabled (up) or disabled (down).
Input Port	The incoming port on which the CES connection exists.
Timeslots	Indicates which timeslots (1-24 for DS1, 1-31 for E1) are configured for the input port. all indicates unstructured service.
Input VPI	The incoming VPI value of the connection.
Input VCI	The incoming VCI value of the connection.
Туре	The type of ATM connection (i.e., PVC or SPVC) that is associated with the CES connection.
Output Port	The outgoing port on which the CES connection exists.
Timeslots	Indicates which timeslots (1-24 for DS1, 1-31 for E1) are configured for the output port. all indicates unstructured service.
Output VPI	The outgoing VPI value of the connection.
Output VCI	The outgoing VCI value of the connection.

To display the advanced settings of the current CES connections, enter the following:

myswitch::configuration ces> show advanced

			Service	Clock		Partial	Max		Integ.
Service	MapVPI	MapVCI	Type	Mode	Cas	Fill	BufSize	CDVT	Period
2024	0	129	structured	synch	basic	0	256	900	2500

#### The fields in this display are defined as follows:

Field	Description
CES Service	The identification number (assigned by the switch) of this CES connection.
MapVPI	The incoming VPI value of the connection.
MapVCI	The incoming VCI value of the connection.
Service Type	Shows if this connection uses structured or unstructured service.
Clock Mode	Synch means that the connection is in synchronous mode (either structured or unstructured). SRTS means that the connection is in asynchronous (unstructured) mode. (Synchronous Residual Time Stamp (SRTS) clock recovery is enabled on this connection.)
Cas	basic indicates that Channel Associated Signalling (CAS) will not be used, cas indicates that CAS will be used.
Partial Fill	Indicates how many of the available 47 payload bytes in each cell are used before they are deemed "full" and ready for transmission across the ATM network (i.e., how much of the ATM cell contains data and how much is filler). The range for this parameter is 5 to 47. The default value is 47, for 47 bytes of data. partialfill is used to minimize network transmission latency and is useful especially with time-sensitive, robbed-bit signalling sources.
Max BufSize	The amount of reassembly buffer space allocated for the connection. The default is 512 bytes per timeslot.
CDVT	The Cell Delay Variation Tolerance for cells being received by the segmentation and reassembly (SAR) engine. The range for this parameter is 100 to 24000 (in $\mu$ s), and the default is 2000 (i.e., 2 ms).
Integ. Period	The amount of time allocated to re-establish the connection before, while, or after the call is established, or in the case of interruption. The default is $2500\mu s$ .

### APPENDIX C

### Converting from FT-PNNI to PNNI

This appendix discusses the conversion of both non-hierarchical and hierarchical FT-PNNI networks to ATM Forum PNNI (hereafter referred to as PNNI) routing. It is assumed that you are familiar with the fundamentals of FT-PNNI and PNNI routing and familiar with FORE's implementation of PNNI as described in Chapter 6 of this manual.

The first section discusses PNNI routing in networks that contain ASX-1000 switches. The various procedures for converting your network are described in the later sections. If you do not have any ASX-1000 switches in your network, you can skip to Section C.2 or Section C.3 to learn how to convert your network from FT-PNNI to PNNI. If you do have ASX-1000 switches in your network, it is recommended that you read Section C.1 first to avoid potential configuration problems when migrating your network.

- Section C.1 ASX-1000 Routing Configuration Issues
- Section C.2 Migration of a Non-Hierarchical FT-PNNI Network
- Section C.3 Migration of a Hierarchical FT-PNNI Network

### C.1 ASX-1000 Routing Configuration Issues

This section discusses routing in hierarchical networks in which there are ASX-1000 switches present in your network. Because each one of the four fabrics appears as a single node in routing and because you cannot take a two-hop path to go from one switch fabric to another switch fabric in the same ASX-1000, there are certain FT-PNNI and PNNI node configurations that should be avoided. These special configurations are addressed individually and alternate configurations are suggested.

### C.1.1 ASX-1000s in FT-PNNI Peer Groups

Figure C.1 shows an invalid configuration of an ASX-1000 divided between two FT-PNNI peer groups called A and B. The ASX-1000's four FT-PNNI nodes are A.1 and A.2 in peer group A and B.1 and B.2 in peer group B. From peer group A, the intra-peer group links of peer group B are not visible. So, the link B.1 to B.2 is not visible in peer group A. Similarly, non-border nodes B.3 and B.4 belonging to peer group B are not visible in peer group A.

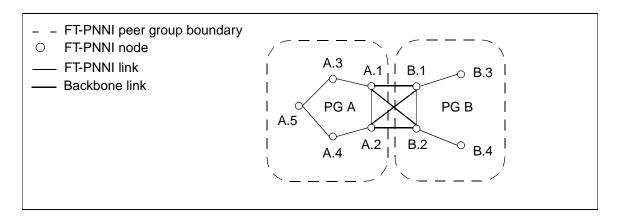


Figure C.1 - Invalid Configuration of ASX-1000 Split between Two FT-PNNI Peer Groups

Suppose you want to route from A.4 to B.3. After the initial step of prefix matching, the Peer Group Summary Node (PGSN) of peer group B (not shown in the figure) is chosen as the destination for this path computation. Since B.1 and B.2 both advertise links to the PGSN and they are equidistant from A.4, A.4 may decide to construct the Destination Transit List (DTL) as A.4 to A.2 to B.2 to B's PGSN. When the setup of this call proceeds into peer group B and reaches node B.2, the only way to route to B.3 is through B.1. But since you cannot take a two-hop path across the ASX-1000, this call setup will fail.

This problem can be avoided by not breaking up an ASX-1000 between multiple FT-PNNI peer groups. For example, in the network in Figure C.2, the entire ASX-1000 could have been made part of peer group B with only A.3, A.4, and A.5 in peer group A. A.3 and A.4 would be the border nodes in peer group A connected to A.1 and A.2 (re-numbered to have B as their peer group ID) in peer group B.

#### C.1.2 ASX-1000s in PNNI Areas

The problem described in Section C.1.1 can also occur in a network of two PNNI peer groups (areas). For example, in the poorly-constructed network shown in Figure C.2, the fabric on the top left (fabric 1) is configured as a split switch with two nodes, one in peer group A and the other in peer group B. All of the other fabrics have one node each. The node on the bottom left fabric (fabric 2) is in peer group A and the nodes in the other two fabrics (3 and 4) are in peer group B. The links between the fabric 2 and the fabrics 3 and 4 do not come up because the peer group IDs of the PNNI nodes connected these links do not match.

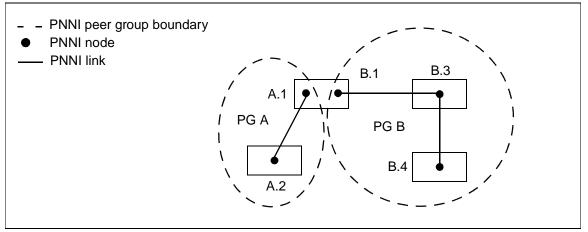


Figure C.2 - Invalid Configuration of ASX-1000 Split between Two PNNI Peer Groups

The problem in the network shown above occurs when computing a path from A.2 to B.3. In this case, routing is forced to use a two-hop path across the backplane, thereby failing to set up a call. To avoid this problem, be sure that you configure all fabrics in an ASX-1000 as part of the same peer group.

### C.1.3 Multiple Gateways in an ASX-1000

Another potential problem can occur when an ASX-1000 in a FT-PNNI area has gateways (split-switches with one FT-PNNI and one PNNI node) to a PNNI area. The ASX-1000 can only have, at most, one of its fabrics as a gateway between the FT-PNNI area and the PNNI area. This is because if two or more of the fabrics have gateways to the PNNI area, then the link connecting any two of the gateways will come up as a PNNI link in the PNNI area. This will divide up the backplane links between two different peer groups. This invalid configuration is shown in Figure C.3.

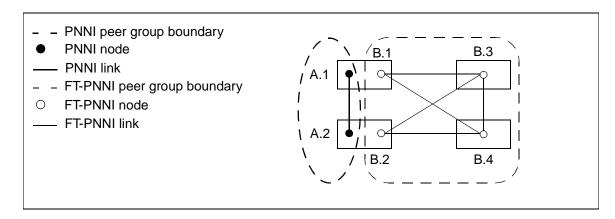


Figure C.3 - Multiple Fabrics of an ASX-1000 Incorrectly Configured as Gateways

In the above figure, the ASX-1000 shown is predominantly part of the FT-PNNI peer group B. However, the ASX-1000 has gateways configured to the PNNI peer group A. Since fabric 1 and fabric 2 are both configured as gateways, the link between them comes up as a PNNI link in peer group A. Because of this, the backplane links of the ASX-1000 are divided between peer group A and peer group B. This may cause two-hop paths to be taken across the backplane links and is an invalid configuration. To avoid this problem, be sure that you configure either only one of the fabrics on an ASX-1000 as a gateway between a single FT-PNNI area and a single PNNI area, or all of the fabrics on an ASX-1000 as a gateway between the same FT-PNNI area and the same PNNI area.

### C.1.4 Migrating from FT-PNNI to PNNI Routing

The migration of FT-PNNI networks to PNNI networks is discussed in the following sections. It is possible that while migrating a network containing ASX-1000 switches, invalid configurations similar to the ones described in this section may occur during the intermediate steps. To avoid these potential problems, migrate all of the fabrics in an ASX-1000 at once, keeping them in the same peer groups or areas.

# C.2 Migration of a Non-Hierarchical FT-PNNI Network

This section discusses the conversion of a non-hierarchical FT-PNNI network to a non-hierarchical (single peer group) PNNI network. It is assumed that all switches in your network are FORE switches.

# **C.2.1 Migration Overview**

The following basic steps are involved in converting your network. Each of these steps is described in detail in the following section. It is recommended that you read the entire section before attempting to change over your network.

- 1. Upgrade each switch in the network to *ForeThought* 5.2.x.
- Choose a switch that is at the edge of the network. Convert it to a gateway switch by changing the default protocol of this switch's default domain to gateway and reboot the switch. Upon rebooting, this switch will come up with an FT-PNNI node and a PNNI node. The PNNI node is isolated at this point and does not have any links attached to it.
- 3. Choose a switch adjacent to the gateway switch, convert it to a gateway switch by changing its default protocol for the default domain to gateway, and reboot this switch. It will come up with a FT-PNNI node and a PNNI node. The link between the two gateway switches is attached on either end to the respective PNNI nodes. This link is now the first PNNI link in the network. This single link constitutes a PNNI area connected by the two gateway switches to the FT-PNNI area. The FT-PNNI area now contains all the links in the network except the one link between the two gateways. (It is assumed that there is only one link between the two gateway switches. If there are multiple links, then all of them will become PNNI links after this step).
- 4. Choose another switch that is adjacent to either one of the two switches already converted to gateway and repeat step 3. As more and more switches are converted to be gateways, the PNNI area becomes progressively larger and the FT-PNNI area becomes smaller.
- 5. For each gateway, when the last link gets converted from FT-PNNI to PNNI (i.e., when the last FT-PNNI switch directly connected to this switch becomes PNNI), modify the default protocol of the default domain to PNNI and reboot the switch. By doing this, the FT-PNNI node on this gateway, which no longer has any link, will be deactivated.

# **C.2.2 Detailed Migration Example**

This section gives a detailed example of how to change the example non-hierarchical FT-PNNI network shown in Figure C.4 to a PNNI network.

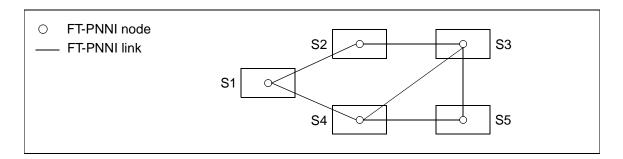


Figure C.4 - A Non-Hierarchical FT-PNNI Network

This network has five switches named S1 through S5. The following steps are involved in changing over this network.

1. Upgrade all of the switches to *ForeThought* 5.2.x using the following AMI command:

oper upgrade <remotehost>:<full path to remotefile>

For more information about upgrading your switches, see Chapter 4 of the Installation and Maintenance manual for your switch.

2. Convert S5 to a gateway switch by modifying the default protocol of the default domain in S5 to gateway using the following AMI command:

conf atmroute domain modify <domain ID> gateway

Reboot switch S5. Upon the reboot, S5 will come up with two nodes: one FT-PNNI node and one PNNI node. By default, these nodes are in areas 4 and 5, respectively, and levels 4 and 5, respectively. At this point, the nodes in the network are divided between the two areas with the PNNI node being the only node in area 5 as shown in Figure C.5.

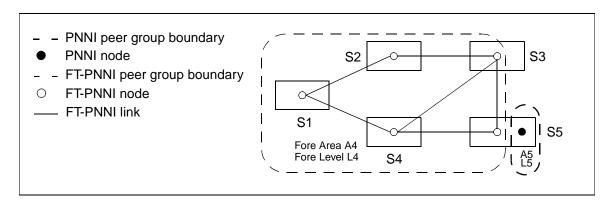


Figure C.5 - S5 Changed to a Gateway Switch

3. Convert S3 to a gateway switch by modifying the default protocol of the default domain in S3 to gateway using the following AMI command:

conf atmroute domain modify <domain ID> gateway

Reboot switch S3. Upon the reboot, S3 will come up with two nodes: one FT-PNNI node in area 4 and one PNNI node in area 5. The link between the two gateway switches S5 and S3 will come up as a PNNI link as shown in Figure C.6. This is the first PNNI link in the network. All of the other links are still FT-PNNI links.

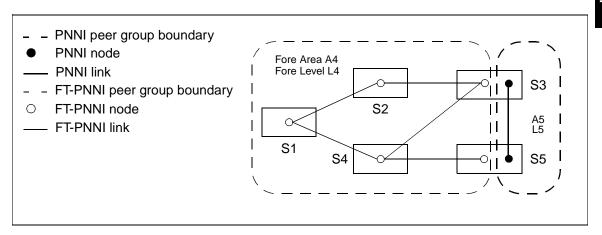


Figure C.6 - S3 Changed to a Gateway Switch

4. Convert S4 to a gateway switch by modifying the default protocol of the default domain in S4 to gateway using the following AMI command:

#### conf atmroute domain modify <domain ID> gateway

Reboot switch S4. Upon the reboot, S4 will come up with a FT-PNNI node in area 4 and a PNNI node in area 5. The link between S4 and S5 will come up as a PNNI link, and the link between S3 and S4 will come up as a PNNI link, since they are both gateway switches.

5. At this point, S5 no longer has any links attached to its FT-PNNI node, so it does not need to be a gateway switch anymore. Modify the default protocol of S5 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch S5. The state of the network at this point is shown in Figure C.7.

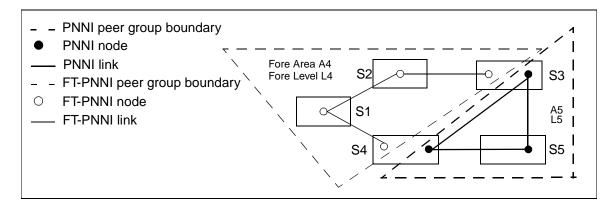


Figure C.7 - S4 Changed to a Gateway Switch and S5 to PNNI

6. Convert S2 to a gateway switch by modifying the default protocol of the default domain in S2 to gateway using the following AMI command:

conf atmroute domain modify <domain ID> gateway

Reboot switch S2. Upon the reboot, S2 will come up with a FT-PNNI node in area 4 and a PNNI node in area 5. The only FT-PNNI links left in the network are between S1 and S2 and between S1 and S4.

7. At this point, S3 no longer has any links attached to its FT-PNNI node, so it does not need to be a gateway switch anymore. Modify the default protocol of S3 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch S3.

8. Since S1 is the last switch in the network to have its default protocol modified, it can be changed directly to a PNNI switch rather than to a gateway switch first and then to a PNNI switch. So, modify the default protocol of S1 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch S1. Now all of the links have been switched over to PNNI.

9. Convert both S2 and S4 to PNNI by modifying the default protocol of both S2 and S4 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switches S2 and S4. The final state of the network upon completion of the conversion is shown in Figure C.8.

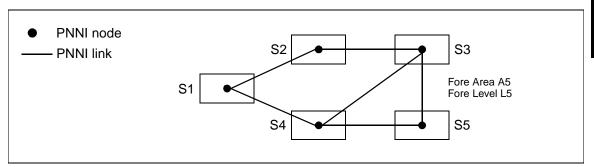


Figure C.8 - A Completely Converted PNNI Network

# C.3 Migration of a Hierarchical FT-PNNI Network

This section discusses the migration of a hierarchical FT-PNNI network to a FORE hierarchical PNNI network. It is assumed that all switches in your network are FORE switches. If your network contains a contiguous backbone, use the instructions found in Section C.3.1. If your network does not contain a contiguous backbone, use the instructions found in Section C.3.2.

# C.3.1 Migration of a Hierarchical FT-PNNI Network with a Contiguous Backbone

There are two different ways to change a FT-PNNI network with a contiguous backbone. If you want to convert your network starting with the backbone, use the method described in Section C.3.1.1. If you want to convert your network starting with the peer groups, use the method described in Section C.3.1.2. You can use either method.

## C.3.1.1 Migration Starting with the Backbone

This section gives a detailed example of how to convert a hierarchical FT-PNNI network with a contiguous backbone as shown in Figure C.9 to a PNNI network. In this figure, only the border nodes are shown. The individual nodes within each peer group are not shown.

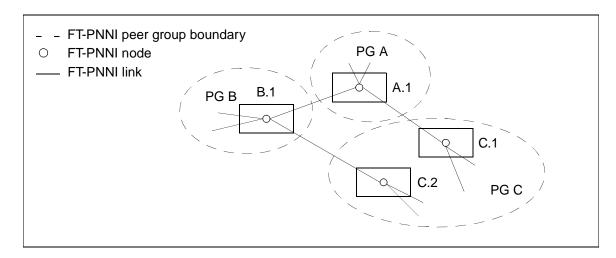


Figure C.9 - Hierarchical FT-PNNI Network with 3 Peer Groups and a Contiguous Backbone

#### C.3.1.1.1 Migration Overview

The following basic steps are involved in the migration. Each of these steps is described in detail in the following sections. It is recommended that you read the entire section before attempting to change over your network.

- 1. Upgrade each switch in the network to *ForeThought* 5.2x.
- 2. Convert the backbone of the FT-PNNI network.
- 3. Convert the individual peer groups of the FT-PNNI network.

### C.3.1.1.1.1 Upgrade the Switches

Upgrade all of the switches to *ForeThought* 5.2.x using the following AMI command:

```
oper upgrade <remotehost>:<full path to remotefile>
```

For more information about upgrading your switches, see Chapter 4 of the Installation and Maintenance manual for your switch.

After the upgrade, there is only one FT-PNNI area with area ID 4 and level 4. The three FT-PNNI peer groups are contained within this one area.

#### C.3.1.1.1.2 Convert the Backbone

C.1 will be the first switch to be converted to a gateway between the existing FT-PNNI area and the PNNI area to be created. Since the backbone is being migrated to PNNI first, the backbone PNNI area should have a higher level (numerically smaller) than the FT-PNNI area. The level of the backbone will be 2.

1. Change the peer group ID of the PNNI node on C.1 to D (so that it has a new peer group ID which is different from the peer group IDs of the three existing FT-PNNI peer groups: A, B, and C). Although the PNNI node in C.1 will not be activated until the default protocol gets changed to gateway, it is possible to change the level of the PNNI node while it is still down. So, change the level of the PNNI node at switch C.1 to 2. Use the following AMI command to change both:

conf atmroute pnni node modify <index> -pgid d -forelevel 2

2. Modify the default protocol of the default domain in C.1 to gateway using the following AMI command:

conf atmroute domain modify <domain ID> gateway

Reboot switch C.1. Upon the reboot, C.1 will come up with two nodes: one FT-PNNI node and one PNNI node. There are now two areas: one FT-PNNI area 4 at level 4, and one PNNI area 5 at level 2.

3. A.1 will be the second switch to be converted to a gateway between the existing FT-PNNI area and the PNNI area to be created. Change the peer group ID of its PNNI node (currently down) to D, modify the level of its PNNI node to 2, and modify the default protocol of the default domain in A.1 to gateway using the following AMI commands:

conf atmroute pnni node modify <index> -pgid d -forelevel 2

conf atmroute domain modify <domain ID> gateway

Reboot switch A.1. Upon the reboot, the link between the PNNI nodes on C.1 and A.1 now becomes the first PNNI link in the network. The state of the network at this stage is shown in Figure C.10.

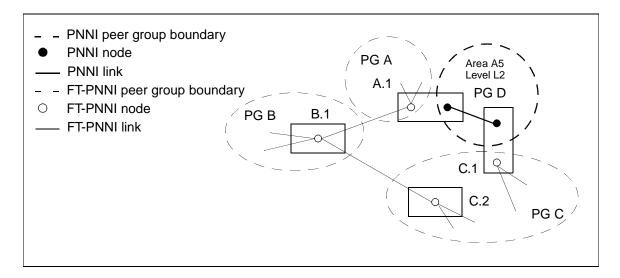


Figure C.10 - C.1 and A.1 as Gateway Switches

4. B.1 will be the next switch to be converted to a gateway switch. Administer the node down. Change the peer group ID of its PNNI node to D, modify the level of its PNNI node to 2, modify the default protocol of the default domain in B.1 to gateway, then administer the node up using the following AMI commands:

conf atmroute pnni node admin <index> down

conf atmroute pnni node modify <index> -pgid d -forelevel 2

conf atmroute domain modify <domain ID> gateway

conf atmroute pnni node admin <index> up

Reboot switch B.1. Upon the reboot, the link between the A.1 and B.1 now becomes a PNNI link.

5. At this point, the FT-PNNI peer group A is now severed from the rest of the FT-PNNI area. Since there are now two areas with area ID 4, one of them has to be renamed. Change the area ID of the FT-PNNI node on A.1 to 2 using the following AMI command:

#### conf atmroute ftpnni forearea 2

Reboot A.1 so the change takes effect. Figure C.11 shows the state of the network at this point.

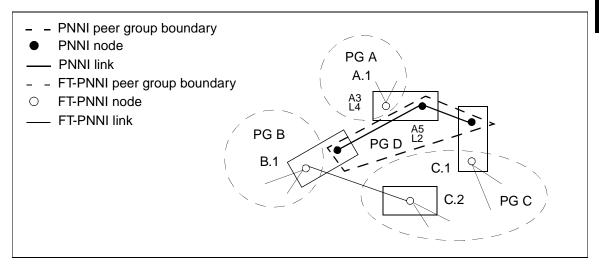


Figure C.11 - Peer Group Severed from the Rest of the FT-PNNI Area

6. C.2 is the final backbone switch to be converted to a gateway switch. Change the peer group ID of its PNNI node (currently down) to D, modify the level of its PNNI node to 2, and modify the default protocol of the default domain in C.2 to gateway using the following AMI commands:

conf atmroute pnni node modify <index> -pgid d -forelevel 2

conf atmroute domain modify <domain ID> gateway

Reboot switch C.2. Upon the reboot, all of the backbone links of the network are converted to PNNI.

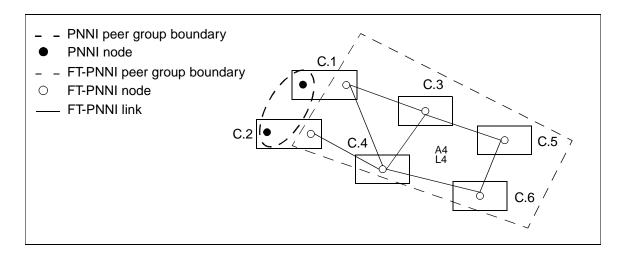
7. At this point, the FT-PNNI peer group B is an area by itself. So, the area ID of the FT-PNNI node on B.1 needs to be changed to 3 using the following AMI command:

#### conf atmroute ftpnni forearea 3

Reboot switch B.1. The FT-PNNI peer group C is also an area by itself, but it can be allowed to retain its original area ID of 4. Upon the reboot of switch B.1, the migration of the backbone is complete.

#### C.3.1.1.1.3 Convert the Individual Peer Groups

Now each individual peer group needs to be converted to PNNI. In this case, peer group C is being used as an example. The other peer groups would be converted using the same steps outlined here. Figure C.12 shows an example configuration of the FT-PNNI nodes within peer group C that need to be converted.



**Figure C.12 -** Peer Group C before the Conversion to PNNI

C.6 will be the first switch to be converted to a gateway between the existing FT-PNNI area and the PNNI area. This area should be at a lower level. The level of this peer group will be 6.

1. Modify the level of the PNNI node (currently down) in C.6 to 2, change the area ID of its PNNI node to 6, and modify the default protocol of the default domain in C.6 to gateway using the following AMI commands:

conf atmroute pnni node modify <index> -forelevel 6 -forearea 6

conf atmroute domain modify <domain ID> gateway

Reboot switch C.6. Upon the reboot, the first PNNI node is created in area 6. At this point, this node has no links attached to it.

2. Convert C.5 to a gateway switch by changing the level of its PNNI node to 6, changing the area ID of C.5 to 6, and modifying the default protocol of the default domain in C.5 to gateway using the following AMI commands:

conf atmroute pnni node modify <index> -forelevel 6 -forearea 6

conf atmroute domain modify <domain ID> gateway

Reboot switch C.6. Upon the reboot, the first PNNI link in area 6 comes up between switches C.6 and C.5.

3. Convert C.4 to a gateway switch by changing the level of its PNNI node to 6, changing the area ID of C.4 to 6, and modifying the default protocol of the default domain in C.4 to gateway using the following AMI commands:

conf atmroute pnni node modify <index> -forelevel 6 -forearea 6

conf atmroute domain modify <domain ID> gateway

Reboot switch C.4. Upon the reboot, the link between C.4 and C.6 comes up as PNNI in area 6. The links between C.4 and C.2 and between C.4 and C.1 will attempt to come up as PNNI (because these links are between two gateway switches), but they will not become operational (because the PNNI links will not reach the two-way-inside hello state).

4. Since C.6 does not need to be a gateway switch anymore, change the default protocol of the default domain in C.6 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.6.

5. Since C.3 is the last FT-PNNI switch, convert it directly to a PNNI switch by changing the level of its PNNI node to 6, changing the area ID of C.3 to 6, and modifying the default protocol of the default domain in C.3 to pnni using the following AMI commands:

conf atmroute pnni node modify <index> -forelevel 6 -forearea 6

conf atmroute domain modify <domain ID> pnni

Reboot switch C.3. Upon the reboot, the link between C.3 and C.1 will attempt to become PNNI, but will not become operational. At this point, the peer group is temporarily separated from the backbone and other peer groups in the network as shown in Figure C.13. It will be restored to full connectivity when C.1 and C.2 have their second PNNI nodes created.

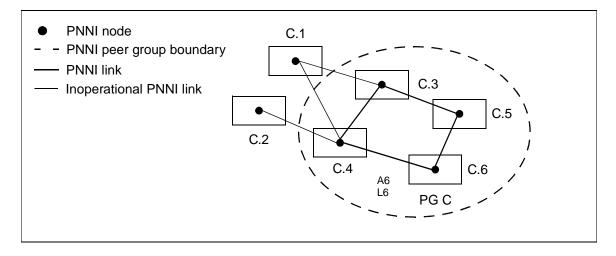


Figure C.13 - C.1 and C.2 Not Part of Peer Group C

6. Since C.5 and C.4 do not need to be gateway switches anymore, change the default protocol of the default domain in C.5 and in C.4 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.5 and switch C.4.

7. Change the default protocol of the default domain in C.1 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.1.

8. Create a second PNNI node in C.1 with a node index of 2, the peer group ID set to C, area 6, and level 6 using the following AMI commands:

Reboot switch C.1. Then, modify the interfaces on C.1 corresponding to the links to C.4 and C.3 and attach them to node 2 on C.1. Use the following AMI command for each link:

conf atmroute pnni interface modify <port> <vpi> -nodeix 2

This brings the links to C.3 and C.4 back to being operational.

9. Change the default protocol of the default domain in C.2 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.2.

10. Create a second PNNI node in C.2 with a node index of 2, the peer group ID set to C, area 6, and level 6 using the following AMI commands:

Reboot switch C.2.

11. Modify the interfaces on C.2 corresponding to the link to C.4 and attach it to node 2 on C.2. Use the following AMI command for each link:

conf atmroute pnni interface modify <port> <vpi> -nodeix 2

Reboot C.2. This brings the link between C.2 and C.4 back to being operational.

12. Upon reboot of C.1, C.3 and C.4 will no longer need to be gateway switches. So, modify the default protocol to PNNI using the following AMI command on each switch:

conf atmroute domain modify <domain ID> pnni

Reboot both C.3 and C.4.

This completes the migration of peer group C to PNNI. Peer groups A and B can be migrated to PNNI in a similar way using the steps found in Section C.3.1.1.1.3. Once they are changed over to PNNI, the entire conversion is complete.

After the migration, the network has four areas at two discrete levels. Each area is a PNNI peer group. The final state of the network upon completion of the conversion is shown in Figure C.14.

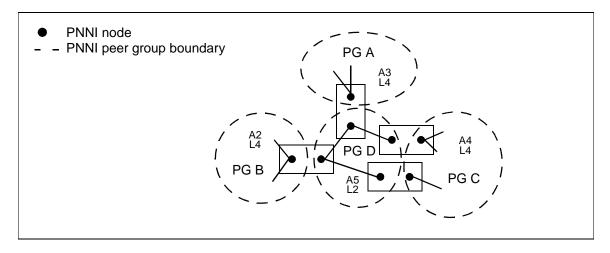


Figure C.14 - A Completely Converted PNNI Network

## C.3.1.2 Migration Starting with the Peer Groups

This section gives a detailed example of how to migrate a hierarchical FT-PNNI network with a contiguous backbone as shown in Figure C.9 to a PNNI network. In that figure, only the border nodes are shown. The individual nodes within each peer group are not shown.

#### C.3.1.2.1 Overview of the Migration

The following basic steps are involved in the migration. Each of these steps is described in detail in the following sections. It is recommended that you read the entire section before attempting to change over your network.

- 1. Upgrade each switch in the network to *ForeThought* 5.2x.
- 2. Convert the individual peer groups of the FT-PNNI network.
- 3. Convert the backbone of the FT-PNNI network.

#### C.3.1.2.1.1 Upgrade the Switches

Upgrade all of the switches from to *ForeThought 5.2.x* using the following AMI command:

```
oper upgrade <remotehost>:<full path to remotefile>
```

For more information about upgrading your switches, see Chapter 4 of the Installation and Maintenance manual for your switch.

#### C.3.1.2.1.2 Convert Peer Group C

1. Peer group C is the first to be converted. Start with switch C.6. Modify the default protocol of the default domain in C.6 to gateway using the following AMI command:

```
conf atmroute domain modify <domain ID> gateway
```

Reboot the switch. This creates the first PNNI node in the network with a default area 5 and level 5 as shown in Figure C.15.

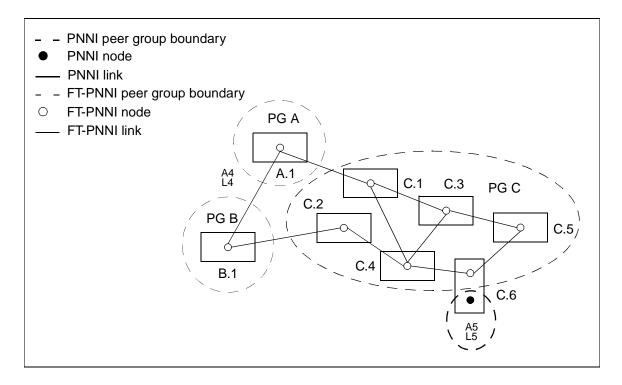


Figure C.15 - C.6 as a Gateway

2. Modify the default protocol of the default domain in C.5 to gateway using the following AMI command:

```
conf atmroute domain modify <domain ID> gateway
```

Reboot switch C.5. This creates the second PNNI node in the network and the first PNNI link in area 5 between C.6 and C.5.

3. Modify the default protocol of the default domain in C.3 to gateway using the following AMI command:

```
conf atmroute domain modify <domain ID> gateway
```

Reboot switch C.3. This creates the second PNNI link in the network between C.5 and C.3.

4. At this point, C.5 does not have any more FT-PNNI links left. So, modify the default protocol of the default domain in C.5 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.5.

5. C.4 is the next switch to be converted. Modify the default protocol of the default domain in C.4 to gateway using the following AMI command:

conf atmroute domain modify <domain ID> gateway

Reboot switch C.4.

6. Since C.6 has now lost its last FT-PNNI link, change the default protocol of the default domain in C.6 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.6.

7. C.1 is the next switch to be changed to a gateway. Since the FT-PNNI peer group of C will become partitioned when C.1 becomes a gateway switch, it is a good idea to change the peer group ID of the FT-PNNI node on C.1 so that other peer groups will not attempt to use peer group C as a transit peer group while computing interpeer group routes. Change the peer group ID of the FT-PNNI node on C.1 to E (something other than C) using the following AMI commands:

conf atmroute ftpnni pgmask <mask>
conf atmroute ftpnni prefix prefix>

Answer no to the question of whether or not you want to the switch to be rebooted.

8. Modify the default protocol of the default domain in C.1 to gateway using the following AMI command:

conf atmroute domain modify <domain ID> gateway

Now reboot switch C.1. Two more PNNI links between C.1 and C.3 and between C.1 and C.4 are created.

9. Since C.3 no longer needs to be a gateway switch, modify the default protocol of the default domain in C.3 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.3.

10. C.2 is the last switch in the peer group to become a gateway. Modify the default protocol of the default domain in C.2 to gateway using the following AMI command:

conf atmroute domain modify <domain ID> gateway

Upon the reboot of C.2, the last FT-PNNI link (between C.2 and C.4) in the peer group becomes a PNNI link.

11. Modify the default protocol of the default domain in C.4 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.4.

This completes the migration of peer group C to PNNI. The current state of the network is shown in Figure C.16.

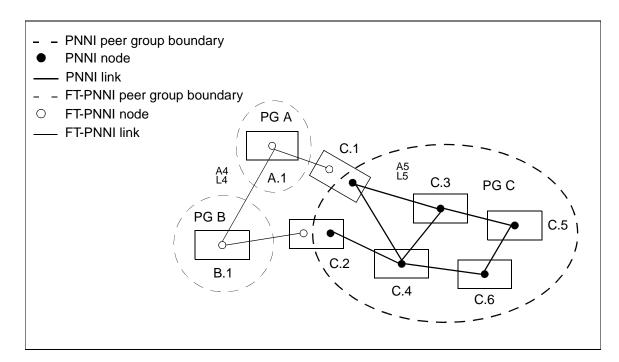


Figure C.16 - Peer Group C Fully Migrated to PNNI

#### C.3.1.2.1.3 Convert Peer Group A

Peer group A is chosen next to be migrated to PNNI. Step 7 from the conversion of peer group C is unnecessary in peer group A's case because there is only one border node in this peer group, and, therefore, the peer group does not get partitioned during the migration. Also, ensure that the area ID of the PNNI nodes in peer group A is different from the area ID of peer group C. Make 6 the area ID of peer group A and make the level 5 (the same level as peer group C).

Migrate all of the nodes in peer group A using the same method that was used to migrate peer group C in Section C.3.1.2.1.2.

#### C.3.1.2.1.4 Convert the Backbone

Once all of the nodes in peer group A have been modified, the link between A.1 and C.1 will become an inoperational PNNI link because it is between two gateway switches and because the peer group IDs of the PNNI nodes on A.1 and C.1 do not match. However, connectivity between the peer groups is still possible through the FT-PNNI backbone.

1. Modify the default protocol of the default domain in C.1 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.1.

2. Create a second PNNI node in C.1 with a node index of 2, the peer group ID set to D, area 3, and level 2 using the following AMI commands:

conf atmroute pnni node modify 2 -pgid d -forelevel 2 -forearea 3

conf atmroute domain modify <domain ID> pnni

Reboot switch C.1. This is the first node in the backbone area.

3. Modify the interface on C.1 corresponding to the link to A.1 and attach it to the newly-created node 2 on C.1 using the following AMI command:

conf atmroute pnni interface modify <port> <vpi> -nodeix 2

4. Change the default protocol of the default domain in A.1 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch A.1. The FT-PNNI link between A.1 and B.1 is lost at this stage and connectivity between peer group A and the rest of the network is severed.

5. Create a second PNNI node in A.1 with a node index of 2, the peer group ID set to D, area 3, and level 2 using the following AMI commands:

conf atmroute pnni node modify 2 -pgid d -forelevel 2 -forearea 3

conf atmroute domain modify <domain ID> pnni

Reboot switch A.1.

6. Modify the interface on A.1 corresponding to the link to C.1 and attach it to the newly-created node 2 on A.1 using the following AMI command:

conf atmroute pnni interface modify <port> <vpi> -nodeix 2

At this point, the connectivity between peer groups A and C should be restored, but peer groups A and B still remain disconnected. The current state of the network is shown in Figure C.17.

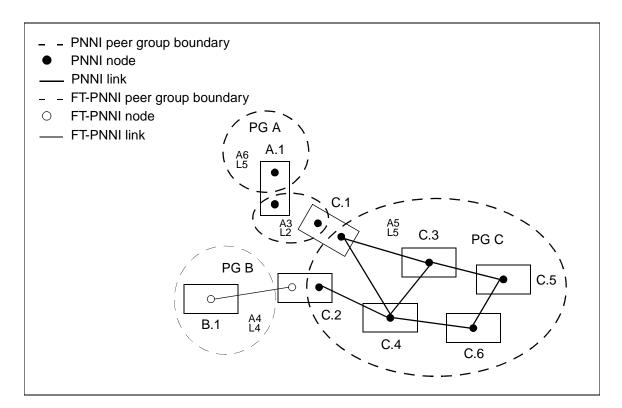


Figure C.17 - Peer Group B Disconnected from Peer Group A

To quickly restore the connectivity, the rest of the backbone should be migrated to PNNI.

7. Change the peer group ID to D, the level to 2, and area to 3 on the PNNI node in B.1, and change the default protocol of the default domain in B.1 to gateway using the following AMI commands:

Reboot switch B.1. The link between A.1 and B.1 should now be restored as a PNNI link and connectivity resumes between A and B.

8. C.2 no longer needs to be a gateway switch. Modify the default protocol of the default domain in C.2 to pnni using the following AMI command:

conf atmroute domain modify <domain ID> pnni

Reboot switch C.2. Upon the reboot, the backbone and peer groups A and C are now running PNNI and only peer group B needs to be migrated. The current state of the network is shown in Figure C.18.

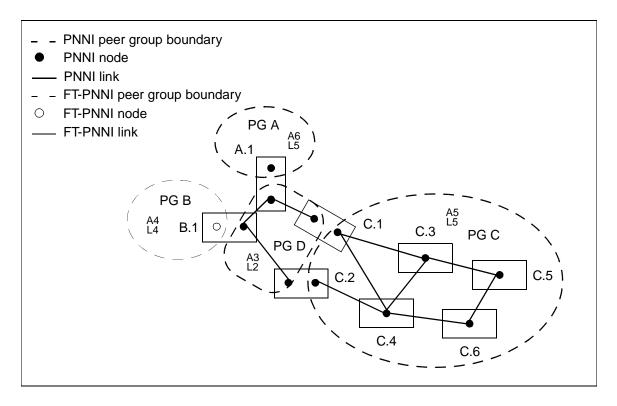


Figure C.18 - A Migrated Backbone

#### C.3.1.2.1.5 Convert Peer Group B

Migrate all of the nodes in peer group B using the same method that was used to migrate peer group C in Section C.3.1.2.1.2. Once peer group B has been converted, the migration to PNNI is complete.

# C.3.2 Migration of a Hierarchical FT-PNNI Network with a Non-Contiguous Backbone

This section gives a detailed example of how to migrate a hierarchical FT-PNNI network with a non-contiguous backbone as shown in Figure C.19 to a PNNI network. In this figure, only the border nodes are shown. The individual nodes within each peer group are not shown.

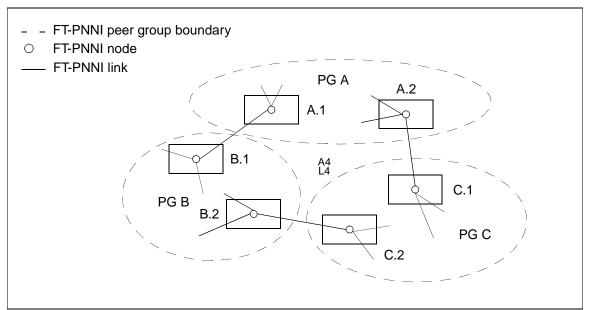


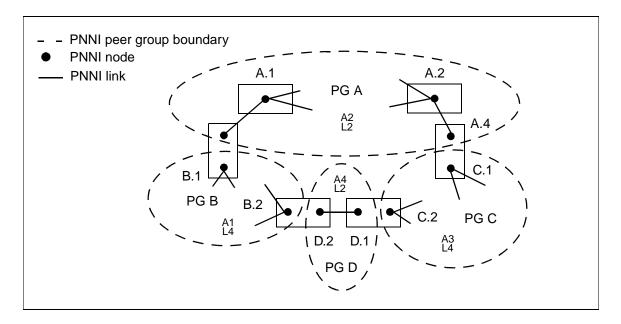
Figure C.19 - Hierarchical FT-PNNI Network with 3 Peer Groups and a Non-contiguous Backbone

Upon the complete migration of this network to PNNI, the network is modified as follows:

- Make peer group A the backbone area in a split-switch based hierarchical network.
- Peer group A would become a PNNI area with an area ID of 2 and level 2 with a single PNNI peer group in the new network.
- Peer group B would become a PNNI area with an area ID 1 and level 4 (a lower level than 2) with a single PNNI peer group.
- Peer group C would become a PNNI area with an area ID 3 and level 4 (a lower level than 2) with a single PNNI peer group.

• The link between B.2 and C.2 in the FT-PNNI network can be placed in a new area (with peer group ID D), which is a higher level area connecting peer groups B and C and acts as a back door entry between the two peer groups. This provides the redundancy that this link offered in the original FT-PNNI network.

Figure C.20 shows the state of the network on the completion of the migration.



**Figure C.20 -** Hierarchical PNNI Network after Migration



# Configuring *FramePlus* Modules

FORE Systems' *FramePlus* network modules (NMFR-4/DS1A and NMFR-4/E1A) are four-port interfaces that support interworking between ATM and Frame Relay, or ATM and Frame-based UNI (FUNI).

All four ports of the DS1 *FramePlus* Network Module can support fractional DS1 services (n x 64 Kbps) where 1 to 24 contiguous or non-contiguous DS0 channels are mapped to a single ATM port.

All four ports of the E1 *FramePlus* Network Module can support fractional E1 services (n x 64 Kbps) where 1 to 31 contiguous or non-contiguous DS0 channels are mapped to a single ATM port.

This appendix provides an overview of a *FramePlus* network module and provides configuration examples. Topics that are covered include the following:

- Section D.1 Frame Relay Overview
- Section D.2 Configuring the Module Level
- Section D.3 Profiles
- Section D.4 Services
- Section D.5 Configuring Frame Relay
- Section D.6 Configuring FUNI
- Section D.7 Upgrading the FramePlus Network Module Software

# **D.1 Frame Relay Overview**

Serving as an interface between user and network equipment, Frame Relay provides a means for statistically multiplexing many logical data channels (or virtual circuits) over a single physical transmission link. Frame Relay is ideal for supporting multiple data streams because it provides flexible and efficient utilization of available bandwidth.

# **D.1.1** Interworking Function (IWF)

When connecting an ATM Customer Premise Equipment (CPE) and a Frame Relay CPE, each device has no knowledge that the distant device supports a different protocol. The *FramePlus* network module performs a service interworking function (IWF) to terminate the respective Frame Relay and ATM services and to perform appropriate interworking conversions before transmitting the CPE information. IWF uses two different modes of operation to address interoperability between upper layer user protocol encapsulation techniques on the *FramePlus* network module: translation mode and transparent mode.

#### **D.1.1.1 Translation Mode**

When the two CPEs support different encapsulation protocols, a translation, or re-encapsulation, must take place in the CPE payload portion of the cells that are being transmitted. Since Frame Relay supports RFC 1490 encapsulation and ATM supports RFC 1483 encapsulation, the information in the frames is converted to cells and vice versa.

## D.1.1.2 Transparent Mode

When both CPEs support the same encapsulation protocol, there is no need for any translation or re-encapsulation of the CPE payload portion of the cells that are being transmitted. Therefore, the traffic is sent "transparently." When the transparent method of IWF is used, there is a mapping that occurs of each Frame Relay DLCI to an ATM VPI, VCI combination.

One example would be if proprietary encapsulation techniques are being used. A second example would be if there is FR/CPE equipment at both edges of the network with ATM providing the transport, there would be no need for any translation.

# **D.2 Configuring the Module Level**

Before configuring any services or PVCs, you should divide the buffer space between the high and low priority buffers because the PVCs are going to go through these buffers. Cells are read from the high priority buffer first. Then, after the high priority buffer is emptied, cells are read from the low priority buffer.

Then you should set the EPD/PPD thresholds for the high and low priority buffers. First, you need to partition the buffer space.

# D.2.1 Dividing the Buffer Space

The FramePlus network module has two priority buffers: high and low. The total, combined buffering space available is 32,768 cells. The amount of buffering space is split between the two buffers using conf module fram setmem. There are four fixed configuration models for partitioning the buffer sizes as shown in Table A.1. The default model is highzero.

Model	High Priority Buffer	Low Priority Buffer
highzero	0 cells	32,768 cells
high1quarter	8,192 cells	24,576 cells
high2quarter	16,384 cells	16,384 cells
high3quarter	24,576 cells	16,384 cells

Table D.1 - Buffer Models

If you choose highzero, you want all of the buffering space to go in the low priority buffer and none in the high priority buffer. If you choose high1quarter, then one fourth of the buffering space goes in the high priority buffer and three fourths goes in the low priority buffer, and so on.

Before you can change the buffer allocation, you must first take the network module out of service by administering it down as follows:

```
myswitch::configuration module> admin 1d down
Disabling the network module will destroy all
existing connections going through it.
Disable the network module [n]? y

myswitch::configuration module> fram
myswitch::configuration module fram> set 1d high2quarter
```

The network module comes back up after you make the change so you do not need to use the conf module admin <module> up command.

# **D.2.2 Setting the Thresholds**

After the buffer space is configured, set the various discard thresholds under conf module fram.



It is important that you follow the <u>order</u> shown in this section for configuring the thresholds. The calculation of each threshold is <u>dependent</u> on the previous one.

For this buffer configuration example, assume you have set the buffers to use high2quarter, meaning that both the high and low priority buffers have 16, 384 cells as shown in Figure D.1.

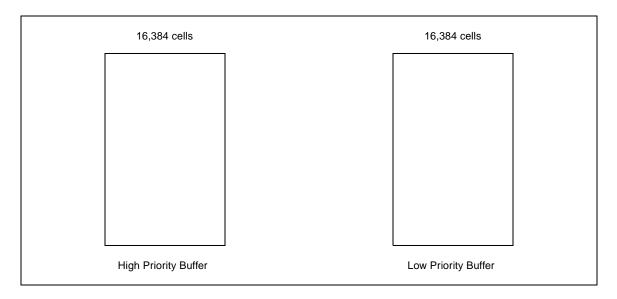


Figure D.1 - Buffer Sizes Configured Using the setmem Command



The number of cells shown in this section for buffer sizes are approximations.

## D.2.2.1 Noting the CLP0PPD Threshold

Each buffer has four different thresholds. The first is for Partial Packet Discard (PPD) on cells with CLP=0. This threshold is automatically set for you as 87.5% of the size of each buffer and cannot be changed. (It is not displayed in AMI currently, but you need to know this value because the other thresholds are calculated on the <u>remaining</u> buffer size.)

Therefore, in this example, since the total buffer size you set is 16,384 cells then the threshold for CLPOPPD is 87.5% of each buffer as shown in Figure D.2:

16,384 cells \* 0.875 percent = 14,336 cells

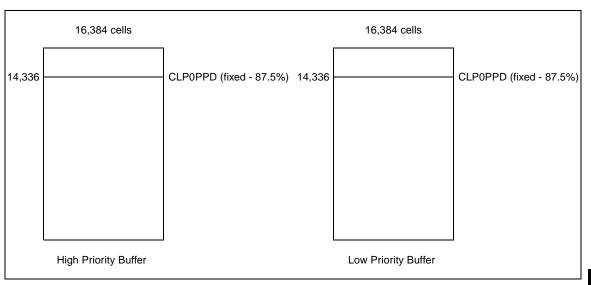


Figure D.2 - CLP0PPD Automatically Calculated

## D.2.2.2 Configuring the CLP1EPD Threshold

Next, calculate the CLP1EPD threshold for each buffer. You have four choices: 25, 37, 50, or 62 percent. For example, you could set the threshold to 50 percent of the <u>remaining</u> buffer size of each buffer, which is 14,336, as shown by the arrow in Figure D.3:

14,336 cells \* 0.5 = 7,168 cells

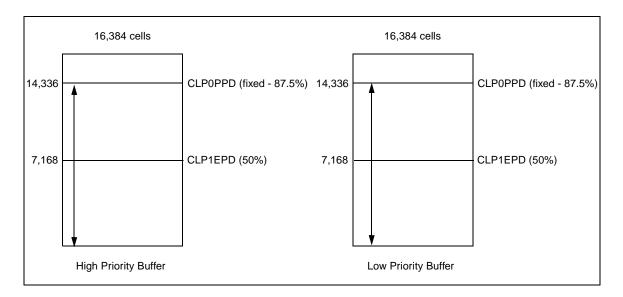


Figure D.3 - Calculated CLP1EPD

## D.2.2.3 Configuring the CLP0EPD Threshold

Now, calculate the CLP0EPD threshold for each buffer. You have four choices: 50, 62, 75, or 87 percent. For example, you could set the threshold to 50 percent of the <u>remaining</u> buffer size of each buffer, which is 7,168, as shown by the arrow in Figure D.4. Then, add the base of 7,168 back in.

7,168 cells \* 0.5 = 3,584 cells + 7,168 cells = 10,752 cells

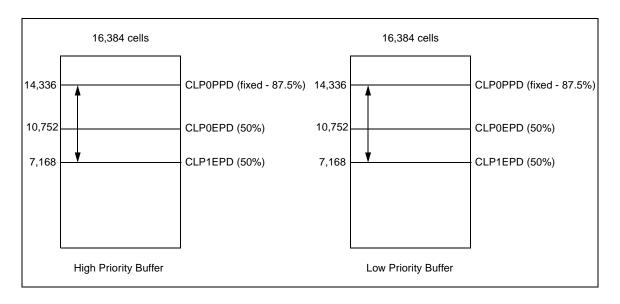


Figure D.4 - Calculated CLP0EPD

## D.2.2.4 Configuring the CLP1PPD Threshold

Finally, calculate the CLP1PPD threshold for each buffer. You have four choices: 50, 62, 75, or 87 percent. For example, you could set it to 50 percent of the <u>remaining</u> buffer size, which is 3,584, as shown by the arrow in Figure D.5. Then, add the base of 7,168 back in as shown:

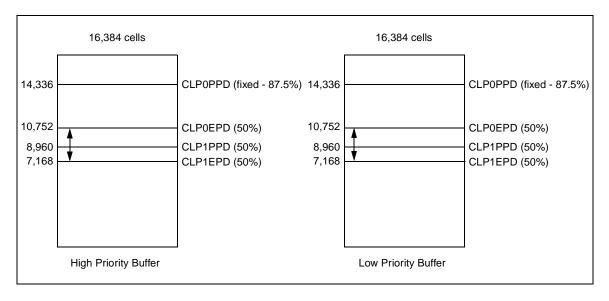


Figure D.5 - Calculated CLP1PPD

# **D.3 Profiles**

A profile lets you define a set of information that can be applied to a particular service or PVC using a single index number, much in the same way that you can define a set of traffic management information in a UPC contract and then apply it to a PVC using a single index number.

For most profile types, several default profiles are provided for your convenience. Only one default is provided for the funi profile. These default profiles reflect what typical networks may need. You cannot delete the default profiles. If a service or connection is created without specifying any profile, the first default profile in the list (index 0) is used.

You can also create new profiles and delete existing profiles, but you cannot modify existing ones. Each profile has an associated reference count, which identifies the number of services or connections that are currently using that profile. You cannot delete any profile that is currently in use (has a non-zero reference count in the Ref field under the profile's **show** command).

There are six different profiles that are supported by the *FramePlus* network module: FRF.8, EPD/PPD, FRRATE, LMI, service, and FUNI. The FRF.8 and FRRATE profiles can only be applied to Frame Relay PVCs. The FUNI profiles can only be applied to FUNI PVCs. The EPD/PPD profiles can be applied to either Frame Relay or FUNI PVCs. The LMI profiles can only be applied to Frame Relay services. The service profiles can be applied to either Frame Relay or FUNI services. All profiles can be configured under the <code>confprofile</code> menu. Please see Part 2 of the *AMI Configuration Commands Reference Manual* for detailed information about the AMI parameters. Each of these profiles is discussed in the following sections.

## D.3.1 EPD/PPD Profile

The Early Packet Discard/Partial Packet Discard (EPD/PPD) profile lets you determine how selective cell discard is performed. The profile then can be applied on a per-VC basis to a Frame Relay PVC using the <code>-epdppd</code> <code><index></code> option under <code>conf fratm pvc new</code> or to a FUNI PVC using the <code>-epdppd</code> <code><index></code> option under <code>conf funi pvc new</code>. The EPD/PPD profile lets you specify which priority buffer (high or low) will handle the traffic of the given PVC. The different thresholds can be enabled or disabled per PVC. These thresholds can be configured under <code>conf module fram</code>.

## D.3.2 FRF.8 Profile

FRF.8 is the Frame Relay Forum's FR/ATM PVC Service Interworking Implementation Agreement. The FRF.8 profile allows you to define the interworking mappings between Frame Relay and ATM. The profile then can be applied on a per-VC basis to a Frame Relay PVC using the -frf8 <index> option under conf fratm pvc new.

The FRF.8 profile lets you decide if the Discard Eligibility (DE) field in every frame is mapped to the Cell Loss Priority (CLP) field in every ATM cell. It lets you define if the Forward Explicit Congestion Notification (FECN) field in every frame is mapped to the EFCI field of every ATM cell. It also lets you define the possible encapsulation translation protocols that can be used.

# D.3.3 Frame Relay Rate Profile

The Frame Relay rate profile allows you to manage and define rate enforcement and rate adaptation characteristics that can be applied on a per-VC basis to a Frame Relay PVC using the -frrate <index> option under conf fratm pvc new.

You can explicitly enable rate enforcement in both the ingress and the egress direction on a FramePlus network module. Ingress rate enforcement is configured using the parameters inBc, inBe, and inCir. Egress rate enforcement is configured using the parameters outBc, outBe, and outCir. The rate adaptation parameters are minBc and cmPeriod. If no rate enforcement parameters are specified, the default values are applied.

The egress rate enforcement also provides an option by which frames on a per-VC basis can be buffered as opposed to being discarded when thresholds have been crossed. This buffering option can increase goodput. The selection of buffering is configured at a service level, but only takes effect on a connection if egress rate enforcement has been enabled. The amount of buffering assigned to a connection is finite and is determined internally by the system. If the buffering has been selected, then the system holds traffic in excess of Bc+Be for a period Tc (= Bc/CIR), before trying to resend the buffered traffic. If the buffer becomes full, excess traffic is discarded. The rate adaptation mechanism is activated for a PVC after an ingress frame, addressed to the PVC, has been received with BECN set.

The Frame Relay rate profile performs the exact same function that UPC contracts do for ATM. If you want to use the overbooking option in the service profile, then you need to apply a Frame Relay rate profile. Bc is the committed burst, Be is the excess burst, and CIR is the committed information rate. CIR is derived as follows:

$$CIR = Bc / time period$$

For example, if the Bc is 10 and the time period is 125 microseconds, then the CIR is  $10 \times 8192$  or 81,920 bps.

EIR is the excess information rate and is derived as follows:

For example, if the Be is 10 and the time period is 125 microseconds, then the EIR is  $10 \times 8192$  or 81,920 bps.

CIR is very similar to VBR SCR and EIR is very similar to VBR PCR.

## D.3.4 LMI Profile

The Link Management Interface (LMI) profile lets you define the version of LMI that is supported by a particular Frame Relay service. The profile can be applied on a per-service basis using the <code>-lmi</code> <code><index></code> option under <code>conffratmnew</code>. If you do not want to use LMI, then enter the <code>none</code> option for <code>-flavor</code> under <code>conffratmnew</code> and use this profile when creating your services.

LMI serves a function that is similar to ILMI for ATM. LMI signalling occurs on DLCI 1023. LMI timers poll the network and report if services are up or down. This information can be optionally translated into ATM OAM F5 cells so that when a DLCI goes down, it triggers an F5 cell for the corresponding PVC. This profile lets you specify the values of the various LMI timers. Please see Part 2 of the AMI Configuration Commands Reference Manual for the default values and the ranges for the timers.

### D.3.5 Service Profile

The generic service profile lets you specify service attributes that are common to Frame Relay interworking services and FUNI services. The profile then can be applied on a per-service basis to a Frame Relay service using the -service <index> option under conf fratm new or to a FUNI service using the -service <index> option under conf funi new.

This profile is important when using multiple timeslots on a port because you must ensure that the access rate of the referenced service profile index is large enough to support them. For example, the default access rate is 64 Kbps. If you leave the rate at 64, you can only use a single timeslot. If you set the rate to 128 Kbps, you can use two timeslots. If you set the rate 1536 Kbps, you can use 24 timeslots, etc.

This profile also lets you define the maximum frame size, the maximum number of DLCIs/PVCs supported, and input and output bandwidth overbooking for a given service.

## D.3.6 FUNI Profile

The FUNI profile lets you determine the VPI, VCI range to use for FUNI connections. The profile then can be applied on a per-service basis to a FUNI service using the <code>-funi <index></code> option under <code>conf funi new</code>.

## **D.4 Services**

A service is a grouping of timeslots on a port. In this respect, a service is similar to an ATM PVP. Multiple DLCIs (connections) can exist within each service.

# **D.5 Configuring Frame Relay**

To configure Frame Relay on a *FramePlus* network module, these steps must be performed in the following order:

- 1. Choose or create any profiles that you wish to use.
- 2. Create each service that you need and associate the profile(s) with it.
- 3. Create each PVC and associate the profile(s) and service(s) with it.

These steps are described in detail in the following sections. There are also steps described for configuring SPANS SPVCs and PNNI SPVCs on a *FramePlus* network module.



A FramePlus network module can run either a Frame Relay over ATM or Frame-based User to Network Interface (FUNI) application. Because Frame Relay is the default, it is assumed that you do not need to change the application from FUNI to Frame Relay at this point. If you need to change the application, see the instructions in Section D.6.1. The Appln field under conf mod display shows which application is currently running (either fratm for Frame Relay or funi for FUNI).

# **D.5.1** Choosing Frame Relay Profiles

First, you need to decide which profiles you want to use. For most profile types, several default profiles are provided for your convenience. If a service or connection is created without specifying any profile, the first default profile in the list (index 0) is used.

To display the default profiles, use the show command under each profile; e.g., conf profile epdppd show. If these defaults do not match your network's needs, you can create your own using the new command under each profile; e.g., conf profile epdppd new. Create the new profiles for each profile type that you are going to use. Once you have determined which default profiles to use, you can move on to Section D.5.2.

## D.5.2 Creating the Services for Frame Relay

Create any Frame Relay services that you want and apply the LMI and service profiles. For example:

```
myswitch::configuration fratm> new <port> <timeslots> [-lmi <index>]
[-service <index>]
[-egress_re (enabled|disabled)] [-status (enabled|disabled)]
[-name <name>]
myswitch::configuration fratm> new 4a1 1 -name service_a
The newly created service id is 4A1:00
myswitch::configuration fratm> new 4a1 2 -lmi 1 -name service_b
The newly created service id is 4A1:01
myswitch::configuration fratm> new 4a1 3 -name service_c
The newly created service id is 4A1:02
myswitch::configuration fratm> show
Searching for FR-ATM services
SvcId TimeSlt
                Admin EgressRE Lmi Serv Traps Stats
                                                         Name
4A1:00 1 up disabled 0 0 enabled disabled service_a
4A1:01 2
                up disabled 1 0 enabled disabled service_b
                 up disabled 0 0 enabled disabled service_c
4A1:02 3
```

As you can see in the Stats field, the collection of statistics is disabled, by default, at the service level. It is enabled, by default, at the module level. If you want to collect statistics, you must enable that functionality here at the service level as well:

```
myswitch::configuration fratm> stats <serviceid> (enabled | disabled)

myswitch::configuration fratm> stats 4A1:00 enabled

myswitch::configuration fratm> stats 4A1:01 enabled

myswitch::configuration fratm> stats 4A1:02 enabled

myswitch::configuration fratm> show

Searching for FR-ATM services

SvcId TimeSlt Admin EgressRE Lmi Serv Traps Stats Name

4A1:00 1 up disabled 0 0 enabled enabled service_a

4A1:01 2 up disabled 1 0 enabled enabled service_b

4A1:02 3 up disabled 0 0 enabled enabled service_c
```

### D.5.3 Creating Frame Relay PVCs

Now that the profiles and services have been created, you can create your PVCs and apply the EPD/PPD, FRF.8, and FRRATE profiles to the PVCs. For example:

```
myswitch::configuration fratm pvc> new <serviceid> <dlci> [-oport <oport>]
        [-ovpi <ovpi>]
        [-ovci <ovci>] [-faupc <index>] [-afupc <index>] [-epdppd <index>]
        [-status (enabled|disabled)] [-frrate <index>] [-frf8 <index>]
        [-name <name>]
```

note: if oport/ovpi/ovci aren't specified, a dangling FRATM PVC will be created.



Currently, the *FramePlus* network module does not permit the establishment of Frame-to-Frame connections across a single fabric. This restriction prevents connections between *FramePlus* ports located on the same network module, or *FramePlus* ports located across separate modules, but attached to the same switch fabric.

```
myswitch::configuration fratm pvc> new 4A1:00 100 -oport 4c1 -ovpi 0 -ovci 100
-epdppd 1 -name pvc_a
```

myswitch::configuration fratm pvc> new 4a1:01 101 -oport 4c1 -ovpi 0 -ovci 101 -name pvc\_b

myswitch::configuration fratm pvc> new 4a1:02 102 -oport 4c1 -ovpi 0 -ovci 102 -name pvc\_c

myswitch::configuration fratm pvc> show

		Input			Outp	ıt							
SvcId	d dlci	Port	VPI	VCI	Port	VPI	VCI	State	Epd	Frr	Frf8	IngRE	Name
4A1:0	00 101	4A1	16	32	4C1	0	101	up	1	0	0	dis	pvc_a
4A1:0	00 101	4C1	0	101	4A1	16	32	up	1	0	0	dis	pvc_a
4A1:0	100	4A1	0	32	4C1	0	100	up	0	0	0	dis	pvc_b
4A1:0	100	4C1	0	100	4A1	0	32	up	0	0	0	dis	pvc_b
4A1:0	102	4A1	32	32	4C1	0	102	up	0	0	0	dis	pvc_c
4A1:0	102	4C1	0	102	4A1	32	32	up	0	0	0	dis	pvc_c

## D.5.4 Configuring Frame Relay SPVCs

You can also configure SPVCs (Smart Permanent Virtual Circuits) on a *FramePlus* network module. An SPVC is a connection that spans multiple switch fabrics and looks like a PVC at the local and remote endpoints with an SVC in the middle. If a link carrying an SPVC goes down and there is an alternate route, then the end switch fabrics of the SPVC automatically reroute the SPVC around the failed link.

However, the procedure for creating both SPANS SPVCs and PNNI SPVCs on a *FramePlus* network module is slightly different than it is for other network modules. Section D.5.4.1 describes how to create a Frame Relay SPANS SPVC. Section D.5.4.2 describes how to create a Frame Relay PNNI SPVC.

#### D.5.4.1 Creating a Frame Relay SPANS SPVC

To create a Frame Relay SPANS SPVC, perform the following steps:

1. Configure the ingress Frame Relay PVC without specifying the -oport, -ovpi, and -ovci parameters. For example:

```
myswitch::configuration fratm pvc> new 4A1:00 100 -epdppd 1 -name spvc_a
```

2. Display the Frame Relay PVC so you can see what Input Port, Input VPI, and Input VCI values were assigned to the PVC. For example:

```
myswitch::configuration fratm pvc> show

Input Output

SvcId dlci Port VPI VCI Port VPI VCI State Epd Frr Frf8 IngRE Name

4A1:00 100 4A1 0 32 up 1 0 0 dis spvc_a
```

3. Configure a new SPANS SPVC where the *<port>*, *<vpi>*, and *<vci>* values are the Input Port, Input VPI, and Input VCI values that are displayed in step 2. The dest parameters identify the endpoint of the connection. For example:

```
myswitch::> open 198.29.22.46 private
Opening a session for "198.29.22.46", please wait...
Connected to "198.29.22.46" (asx200bx).

*fishtank::> localhost

myswitch::configuration spvc spans> new 4a1 0 32 198.29.22.46 1a2 0 100
```

This creates a SPANS SPVC where the ATM portion of the interworking connection is supported via the SPVC logic.

#### D.5.4.2 Creating a Frame Relay PNNI SPVC

To create a Frame Relay PNNI SPVC, perform the following steps:

1. Configure the ingress Frame Relay PVC without specifying the -oport, -ovpi, and -ovci parameters. For example:

```
myswitch::configuration fratm pvc> new 4A1:00 100 -epdppd 1 -name spvc_a
```

2. Display the Frame Relay PVC so you can see what Input Port, Input VPI, and Input VCI values were assigned to the PVC. For example:

```
myswitch::configuration fratm pvc> show

Input
Output
SvcId dlci Port VPI VCI Port VPI VCI State Epd Frr Frf8 IngRE Name
4A1:00 100 4A1 0 32 up 1 0 0 dis spvc_a
```

3. Configure a new PNNI SPVC where the *<port>*, *<vpi>*, and *<vci>* values are the Input Port, Input VPI, and Input VCI values that are displayed in step 2. The dest parameters identify the endpoint of the connection. For example:

```
myswitch::configuration spvc pnni> new 4a1 0 32 47.0005.80.ffe100.0000.f2lb.19cd:ldl -destvpi 0 -destvci 100 -name spvc_a
```



Usually, when creating a PNNI SPVC, the **-destvpi** and **-destvci** parameters are optional (i.e., if these values are not supplied the terminating switch uses the same VPI and VCI values as the originating switch). However, when at least one endpoint of the PNNI SPVC is a port on a *FramePlus* network module, you must specify these parameters.

This creates a bi-directional PNNI SPVC from the source to the destination where the ATM portion of the interworking connection is supported via the SPVC logic.

## **D.6 Configuring FUNI**

To configure FUNI on a *FramePlus* network module, these steps must be performed in the following order:

- 1. Change the application key, if necessary.
- 2. Choose or create any profiles that you wish to use.
- 3. Create each service that you need and associate the profile(s) with it.
- 4. Create each PVC and associate the profile(s) and service(s) with it.

These steps are described in detail in the following sections. There are also steps described for configuring PNNI SPVCs on a *FramePlus* network module.

## D.6.1 Changing the Application Key

This command lets you configure a *FramePlus* network module to run either Frame Relay over ATM or to run frame-based User to Network Interface (FUNI). The network module can only run one application or the other at a time. However, you can have some network modules running FUNI and some running Frame Relay in the same switch fabric.

FramePlus network modules run Frame Relay by default. To run FUNI services on the network module, please contact FORE's Technical Assistance Center for a valid FUNI key. To change from FUNI back to Frame Relay, the key is fratm170358. The Appln field under conf module display shows you which application is currently running.

The application key should be specified only if you want to reconfigure the network module to run a different type of application. Before changing the application, you must administer the network module down as follows:

```
myswitch::configuration module> admin 4a down
Disabling the network module will destroy all
existing connections going through it.
Disable the network module [n]? y
```

Now you can change the application. When you change the application, the switch <u>deletes</u> all existing services and PVCs that use a different application, and removes them from the CDB (i.e., if you are changing from Frame Relay to FUNI, the switch deletes existing Frame Relay information, and vice versa). The switch warns you as follows:

```
myswitch::configuration module fram> app 4a XXXXXXXXXXC Changing application may cause deletion of service/connection on switch and CDB. Proceed [n]? y
```

Then, you need to administer the network module up as follows:

```
myswitch::configuration module> admin 4a up

myswitch::configuration module> display

Module Appln Appl Boot Stats Oam Operational Product
swRel swRel monitor monitor state number

4A funi 1.0.0 1.0.0 enabled enabled appluprunning NMFR-4/DS1A
```



If you attempt to display information before administering the network module up, a message is displayed that no configuration information is available for this network module.

## D.6.2 Creating the Profiles for FUNI

First, you need to decide which profiles you want to use. A default profile is provided for your convenience. If a service or connection is created without specifying any profile, the default profile in the list is used.

To display the default profile, use the conf profile funi show command. If this default does not match your network's needs, you can create your own using the new command under each profile; e.g., conf profile funi new. The values that you choose for [-minVci <vci>] and [-maxVci <vci>] determine the range of VCIs that you can use for <fvci> under conf funi pvc new. (See Part 2 of the AMI Configuration Commands Reference Manual for more information about the parameters.) Once you have determined which profile(s) to use, you can move on to Section D.6.3.

## **D.6.3 Creating FUNI Services**

Create any FUNI services that you want and apply the FUNI and service profiles. For example:

```
myswitch::configuration funi > new <port > <timeslots > [-funi <index >]
[-service <index>]
[-status (enabled|disabled)] [-name <name>]
myswitch::configuration funi > new 4al 1 -funi 1 -name service_a
The newly created service id is 4A1:00
myswitch::configuration funi> new 4a1 2 -service 1 -name service_b
The newly created service id is 4A1:01
myswitch::configuration funi> new 4a1 3 -name service_c
The newly created service id is 4A1:02
myswitch::configuration funi> show
Searching for FUNI services
SvcId TimeSlot Admin Funi Serv Signal
                                           Traps Stats
                                                            Name
 4A1:00 1
                           1 0 nonexist enabled disabled service_a
                  up
4A1:01 2
                           0 1 nonexist enabled disabled service_b
                  up
4A1:02 3
                  uр
                           0
                                0 nonexist enabled disabled service_c
```

As you can see in the Stats field, the collection of statistics is disabled, by default, at the service level. It is enabled, by default, at the module level. If you want to collect statistics, you must enable that functionality here at the service level as well:

```
myswitch::configuration funi> stats <serviceid> (enabled | disabled)
myswitch::configuration funi> stats 4A1:00 enabled
myswitch::configuration funi> stats 4A1:01 enabled
myswitch::configuration funi> stats 4A1:02 enabled
myswitch::configuration module fram> stats <module> (enabled|disabled)
myswitch::configuration module fram> stats 4a enabled
myswitch::configuration funi> show
Searching for FUNI services
SvcId TimeSlot Admin Funi Serv Signal
                                          Traps Stats
                                                           Name
4A1:00 1
                          1 0 nonexist enabled enabled service_a
                 up
                          0 1 nonexist enabled enabled service_b
4A1:01 2
                 up
                         0 0 nonexist enabled enabled service_c
4A1:02 3
                 up
```

### D.6.4 Creating FUNI PVCs

Now that the profiles and services have been created, you can create your FUNI PVCs and apply the EPD/PPD profiles to the PVCs. For example:

```
myswitch::configuration funi pvc> new <serviceid> <fvpi> <fvci> [-oport <oport>]
[-ovpi <ovpi>] [-ovci <ovci>] [-faupc <index>] [-afupc <index>] [-epdppd <index>]
[-status (enabled | disabled)] [-name <name>]
```

note: if oport/ovpi/ovci aren't specified, a dangling FRATM PVC will be created.



The default minimum VCI is 32 and the default maximum VCI is 63. If you are trying to create a PVC on a FUNI service that uses the default FUNI service profile, you are limited to this VCI range when specifying the *<fvci>* parameter.



Currently, the *FramePlus* network module does not permit the establishment of Frame-to-Frame connections across a single fabric. This restriction prevents connections between FramePlus ports located on the same network module, or *FramePlus* ports located across separate modules, but attached to the same switch fabric.

```
myswitch::configuration funi pvc> new 4A1:00 0 40 -oport 4c1 -fvpi 0 -fvci 40
-name pvc_a
```

myswitch::configuration funi pvc> new 4A1:01 0 41 -oport 4c1 -fvpi 0 -fvci 41
-name pvc\_b

myswitch::configuration funi pvc> new 4A1:02 0 42 -oport 4c1 -fvpi 0 -fvci 42 -name pvc\_c

myswitch::configuration funi pvc> show

	FUNI	FUNI	Input		(	Output					
SvcId	VPI	VCI	Port	VPI	VCI	Port	VPI	VCI	Status	Eppd	Name
4A1:02	0	42	4A1	32	42	4C1	0	42	up	0	pvc_c
4A1:02	0	42	4C1	0	42	4A1	32	42	up	0	pvc_c
4A1:00	0	40	4A1	0	40	4C1	0	40	up	0	pvc_a
4A1:00	0	40	4C1	0	40	4A1	0	40	up	0	pvc_a
4A1:01	0	41	4A1	16	41	4C1	0	41	up	0	pvc_b
4A1:01	0	41	4C1	0	41	4A1	16	41	up	0	pvc_b

## D.6.5 Configuring FUNI SPVCs

You can also configure SPVCs (Smart Permanent Virtual Circuits) on a *FramePlus* network module. An SPVC is a connection that spans multiple switch fabrics and looks like a PVC at the local and remote endpoints with an SVC in the middle. If a link carrying an SPVC goes down and there is an alternate route, then the end switch fabrics of the SPVC automatically reroute the SPVC around the failed link.

However, the procedure for creating both SPANS SPVCs and PNNI SPVCs on a *FramePlus* network module is slightly different than it is for other network modules. Section D.6.5.1 describes how to create a FUNI SPANS SPVC. Section D.6.5.2 describes how to create a FUNI PNNI SPVC.

#### D.6.5.1 Creating a FUNI SPANS SPVC

To create a FUNI SPANS SPVC, perform the following steps:

1. Configure the ingress FUNI PVC without specifying the -oport, -ovpi, and -ovci parameters. For example:

```
myswitch::configuration funi pvc> new 4A1:00 0 40 -epdppd 1 -name spvc_a
```

2. Display the FUNI PVC so you can see what Input Port, Input VPI, and Input VCI values were assigned to the PVC. For example:

3. Configure a new SPANS SPVC where the *<port>*, *<vpi>*, and *<vci>* values are the Input Port, Input VPI, and Input VCI values that are displayed in step 2. The dest parameters identify the endpoint of the connection. For example:

```
myswitch::> open 198.29.22.46 private
Opening a session for "198.29.22.46", please wait...
Connected to "198.29.22.46" (asx200bx).

*fishtank::> localhost

myswitch::configuration spvc spans> new 4a1 0 40 198.29.22.46 1a2 0 100
```

This creates a SPANS SPVC where the ATM portion of the interworking connection is supported via the SPVC logic.

#### D.6.5.2 Creating a FUNI PNNI SPVC

To create a FUNI PNNI SPVC, perform the following steps:

1. Configure the ingress FUNI PVC without specifying the -oport, -ovpi, and -ovci parameters. For example:

```
myswitch::configuration funi pvc> new 4A1:00 0 40 -epdppd 1 -name spvc_a
```

2. Display the FUNI PVC so you can see what iport, ivpi, and ivci values were assigned to the PVC. For example:

3. Configure a new PNNI SPVC where the *<port>*, *<vpi>*, and *<vci>* values are the iport, ivpi, and ivci values that are displayed in step 2. The dest parameters identify the endpoint of the connection. For example:

```
myswitch::configuration spvc pnni> new 4c1 0 100
47.0005.80.ffe100.0000.f21b.19cd:4a1 -destvpi 0 -destvci 40 -name spvc_a
```



Usually, when creating a PNNI SPVC, the -destvpi and -destvci parameters are optional (i.e., if these values are not supplied the terminating switch uses the same VPI and VCI values as the originating switch). However, when at least one endpoint of the PNNI SPVC is a port on a *FramePlus* network module, you must specify these parameters.

This creates a bi-directional PNNI SPVC from the source to the destination where the ATM portion of the interworking connection is supported via the SPVC logic.

# D.7 Upgrading the FramePlus Network Module Software

Because the *FramePlus* network module has an on-board i960 processor, you can upgrade the the network module application software using TFTP. The method for upgrading the network module software is similar to that for upgrading the switch software.

myswitch::configuration module fram> upgrade <module> <remotehost>:<fullpath to
remotefile>

#### These parameters are defined as follows:

Parameter	Description					
module	The FramePlus network module on which you want to upgrade the software.					
remotehost	The IP address of the remote host on which the upgrade file resides.					
full path to remotefile	The full path name of the upgrade file.					

#### For example:

```
myswitch::configuration module fram> upgrade 1A 124.11.1.13:/net/mercFCS01
WARNING: Do not remove the netmod while upgrading
{File download to netmod 1A successful}
Transfer successful.
Reset the network module [n]? y
```

To display the current revision number of the application software, use the following command:

Since this command uses TFTP as the transfer protocol, the remote host on which the upgrade file resides must be a tftpboot server. If you are unsure of how to configure the bootp server and the tftpboot server properly, see Chapter 4 of the Installation and Maintenance manual for your switch.

Configuring FramePlus Modules

# **Acronyms**

The networking terms in the following list are defined in the Glossary of this manual. Glossary items are listed alphabetically according to the full term.

AAL ATM Adaptation Layer
ABR Available Bit Rate

**ACM** Address Complete Message

ACR Allowable Cell Rate

**ADPCM** Adaptive Differential Pulse Code Modulation

AHFG ATM-attached Host Functional Group

AIMUX ATM Inverse Multiplexing
AIS Alarm Indication Signal
AMI Alternate Mark Inversion
AMI ATM Management Interface

ANSI American National Standards Institute
APCM Adaptive Pulse Code Modulation
API Application Program Interface

APP Application Program

APS Automatic Protection Switching
ARP Address Resolution Protocol

**ASCII** American Standard Code for Information Interchange

**ATDM** Asynchronous Time Division Multiplexing

ATM Asynchronous Transfer Mode
AUI Attachment User Interface
BBZS Bipolar 8 Zero Substitution

**BCOB** Broadband Connection Oriented Bearer

BCOB-A Bearer Class A
BCOB-C Bearer Class C
BCOB-X Bearer Class X

**BECN** Backward Explicit Congestion Notification

**BER** Bit Error Rate

BES Bursty Errored SecondsBGP Border Gateway ProtocolB-ISDN Inter-Carrier Interface.

BIP Bit Interleaved Parity

**B-ISDN** Broadband Integrated Services Digital Network

B-ISUP Broadband ISDN User's Part

#### Acronyms

Building Integrated Timing Supply

BNC Bayonet-Neill-Concelman
BPDU Bridge Protocol Data Unit

bps Bits per SecondBPV Bipolar Violation

B-TE Broadband Terminal Equipment
BUS Broadcast and Unknown Server
CAC Connection Admission Control
CAS Channel Associated Signaling

CBDS Connectionless Broadband Data Service

**CBR** Constant Bit Rate

**CCITT** International Telephone and Telegraph Consultative Committee

**CCS** Common Channel Signaling

CDV Cell Delay Variation
CE Connection Endpoint

CEI Connection Endpoint Identifier
CES Circuit Emulation Service
CGA Carrier Group Alarm

CIP Carrier Identification Parameter
CIR Committed Information Rate

CLIP Classical IP
CLP Cell Loss Priority
CLR Cell Loss Ratio-1-15
CLS Connectionless service

**CMIP** Common Management Interface Protocol

**CMR** Cell Misinsertion Rate

**CPE** Customer Premise Equipment

CRA Cell Rate Adaptation
CRC Cyclic Redundancy Check

CRS Cell Relay Service
CS Controlled Slip, or

CSU Channel Service Unit
CTD Cell Transfer Delay
CTS Clear To Send

DACS Digital Access and Cross-Connect System
DARPA Defense Advanced Research Projects Agency

**DCC** Data Country Code

DCE Data Communications Equipment
DCS Digital Cross-connect System
DES Destination End Station

**DFA** DXI Frame Address

**DLCI** Data Link Connection Identifier

**DNS** Domain Naming System

**DSn** Digital Standard n (n=0, 1, 1C, 2, and 3)

**DSR** Data Set Ready

DTE Data Terminal Equipment
DTR Data Terminal Ready

**EEPROM** Electrically Erasable Programmable Read Only Memory

**EFCI** Explicit Forward Congestion Indication

**EGP** Exterior Gateway Protocol

**EIA** Electronics Industries Association

**EISA** Extended Industry Standard Architecture

EMI Emulated Local Area Network Electromagnetic Interference

**EPROM** Erasable Programmable Read Only Memory

EQL Equalization

ER Explicit Rate

ES End System, or

Errored Second

**ESF** Extended Super Frame **ESI** End System Identifier

**EXZ** Excessive Zeroes (Error Event)

FC Face Contact

**FCC** Federal Communications Commission

**FCS** Frame Check Sequence

FDDI Fiber Distributed Data Interface
FDM Frequency Division Multiplexing

FEE Far End Block Error
FEC Forward Error Correction

**FECN** Forward Explicit Congestion Notification

FERF Far End Receive Failure
FIFO First-In, First-Out
FRS Frame-Relay Service
FTP File Transfer Protocol
FT-PNNI ForeThought PNNI
FUNI Frame-Based UNI

GCAC Generic Connection Admission Control

GCRA Generic Cell Rate Algorithm

GFC Generic Flow Control HDB3 High Density Bipolar

**HDLC** High Level Data Link Control

**HEC** Header Error Control

**HIPPI** High Performance Parallel Interface

**HSSI** High-Speed Serial Interface

ICMP Internet Control Message Protocol

#### Acronyms

**IDU** Interface Data Unit

**IEEE** Institute of Electrical and Electronics Engineers

IETF Internet Engineering Task Force
ILMI Interim Local Management Interface

IP Internet Protocol

IPX Internetwork Packet Exchange

IS Intermediate system

ISDN Integrated Services Digital Network
ISO International Standards Organization

ITU-T International Telecommunication Union Telecommunication

IWF Interworking Function IXC Interexchange Carriers

JPEG Joint Photographic Experts Group

Kbps Kilobits per second
LAN Local Area Network
LANE LAN Emulation

LAPB Link Access Procedure, Balanced LATA Local Access and Transport Area

LINE Build Out
LINE Code Violations

**LE\_ARP** LAN Emulation Address Resolution Protocol

**LEC** LAN Emulation Client

**LECS** LAN Emulation Configuration Server

LES LAN Emulation Server
LLC Logical Link Control
LOF Loss Of Frame
LOP Loss Of Pointer
LOS Loss Of Signal
LSB Least Significant Bit
MAC Media Access Control

MAC Media Access Control

MAN Metropolitan Area Network

MAU Media Attachment Unit

MBS Maximum Burst Size

MCDV Maximum Cell Delay Variance
MCLR Maximum Cell Loss Ratio

MCR Minimum Cell Rate

MCTDMaximum Cell Transfer DelayMIBManagement Information BaseMICMedia Interface Connector

MID Message Identifier

MMF Multimode Fiber Optic Cable
MPEG Motion Picture Experts Group
MPOA Multiprotocol over ATM

MSB Most Significant Bit

MTU Maximum Transmission Unit
NM Network Management Entity
NML Network Management Layer
NMS Network Management Station

NNI Network-to-Network Interface or Network Node Interface

**NPC** Network Parameter Control

NRZ Non Return to Zero

NRZI Non Return to Zero Inverted
NSAP Network Service Access Point
NTSC National TV Standards Committee
OAM Operation and Maintenance Cell

OC-n Optical Carrier level-n
OID Object Identifier
OOF Out-of-Frame

OSI Open Systems Interconnection
OSPF Open Shortest Path First Protocol
OUI Organizationally Unique Identifier
PAD Packet Assembler Disassembler

PAL Phase Alternate Line
PBX Private Branch Exchange

**PCI** Peripheral Component Interconnect

**PCM** Pulse Code Modulation

PCR Peak Cell Rate

PDN Public Data Network
PDU Protocol Data Unit
PHY Physical Layer

ping Packet Internet Groper

**PLCP** Physical Layer Convergence Protocol

PLP Packet Level Protocol
PM Physical Medium

PMD Physical Medium Dependent

PNNI Private Network Node Interface or Private Network-to-Network Interface

**PPP** Point-to-Point Protocol

**PROM** Programmable Read-Only Memory

PRS Primary Reference Source
PSN Packet Switched Network

PT Payload Type

PVC Permanent Virtual Circuit (or Channel)
PVCC Permanent Virtual Channel Connection
PVPC Permanent Virtual Path Connection

QD Queuing Delay
QoS Quality of Service

#### Acronyms

RD Routing Domain
RFCs Requests For Comment
RFI Radio Frequency Interference
RIP Routing Information Protocol
RISC Reduced Instruction Set Computer

RTS Request To Send
SA Source Address
SA Source MAC Address
SAP Service Access Point

SAR Segmentation And Reassembly

SC Structured Cabling, or

Structured Connectors, or

Stick and Click

**SCR** Sustainable Cell Rate

SCSI Small Computer Systems Interface
SDLC Synchronous Data Link Control

**SDU** Service Data Unit

SEAL Simple and Efficient Adaptation Layer
SECAM Systeme En Coleur Avec Memoire

SEL Selector

**SES** Severely Errored Seconds

**SF** Super Frame

**SGMP** Simple Gateway Management Protocol

**SIR** Sustained Information Rate

**SLIP** Serial Line IP

SMDS Switched Multimegabit Data Service

**SMF** Single Mode Fiber

SMTP Simple Mail Transfer Protocol
SNA Systems Network Architecture
SNAP SubNetwork Access Protocol
SNI Subscriber Network Interface

**SNMP** Simple Network Management Protocol

**SONET** Synchronous Optical Network

**SPANS** Simple Protocol for ATM Network Signalling

**SPARC** Scalable Processor Architecture Reduced instruction set Computer

**SPE** Synchronous Payload Envelope

SPVC Smart PVC

SS7 Signaling System No. 7

**SSCOP** Service Specific Connection Oriented Protocol

SSCS Service Specific Convergence Sublayer

ST Straight Tip, or

Stick and Turn

**STM** Synchronous Transfer Mode

STP Shielded Twisted Pair, Spanning Tree Protocol

STS Synchronous Transport Signal

SVC Switched Virtual Circuit (or Channel)
SVCC Switched Virtual Channel Connection
SVPC Switched Virtual Path Connection

TAXI Transparent Asynchronous Transmitter/Receiver Interface

TC Transmission Convergence
TCP Transmission Control Protocol

TCP/IP Transmission Control Protocol/Internet Protocol

TCR Tagged Cell Rate

TCS Transmission Convergence Sublayer

**TDM** Time Division Multiplexing

TE Terminal Equipment

**TFTP** Trivial File Transfer Protocol

TM Traffic Management
UAS Unavailable Seconds
UBR Unspecified Bit Rate
UDP User Datagram Protocol
UNI User-to-Network Interface
UPC Usage Parameter Control

**UTOPIA** Universal Test & Operations Interface for ATM

**UTP** Unshielded Twisted Pair

VBR Variable Bit Rate

VC Virtual Channel (or Circuit)
VCC Virtual Channel Connection
VCI Virtual Channel Identifier
VCL Virtual Channel Link
VINES Virtual Network Software
VLAN Virtual Local Area Network

**VP** Virtual Path

VPC Virtual Path Connection
VPDN Virtual Private Data Network

VPI Virtual Path Identifier
VPL Virtual Path Link
VPN Virtual Private Network
VPT Virtual Path Terminator

**VS/VD** Virtual Source/Virtual Destination

VT Virtual Tributary WAN Wide-Area Network

**ZBTSI** Zero Byte Time Slot Interchange

Acronyms

# **Glossary**

**10Base-T -** a 10 Mbps baseband Ethernet specification utilizing twisted-pair cabling (Category 3, 4, or 5). 10BaseT, which is part of the IEEE 802.3 specification, has a distance limit of approximately 100 meters per segment.

**802.1d Spanning Tree Bridging -** the IEEE standard for bridging; a MAC layer standard for transparently connecting two or more LANs (often called subnetworks) that are running the same protocols and cabling. This arrangement creates an extended network, in which any two workstations on the linked LANs can share data.

**802.3 Ethernet -** the IEEE standard for Ethernet; a physical-layer standard that uses the CSMA/CD access method on a bus-topology LAN.

**802.5 Token Ring -** the IEEE physical-layer standard that uses the token-passing access method on a ring-topology LAN.

**AAL Connection -** an association established by the AAL between two or more next higher layer entities.

**Adapter -** A fitting that supplies a passage between two sets of equipment when they cannot be directly interconnected.

**Adaptive Differential Pulse Code Modulation (ADPCM)** - A technique that allows analog voice signals to be carried on a 32K bps digital channel. Sampling is done at 8Hz with 4 bits used to describe the difference between adjacent samples.

**Adaptive Pulse Code Modulation (APCM) -** A technique that effectively reduces occupied bandwidth per active speaker by reducing sampling rates during periods of overflow peak traffic.

**Address -** A unique identity of each network station on a LAN or WAN.

**Address Complete Message (ACM) -** A B-ISUP call control message from the receiving exchange to sending exchange indicating the completion of address information.

**Address Mask** - a bit mask used to identify which bits in an address (usually an IP address) are network significant, subnet significant, and host significant portions of the complete address. This mask is also known as the subnet mask because the subnetwork portion of the address can be determined by comparing the binary version of the mask to an IP address in that subnet. The mask holds the same number of bits as the protocol address it references.

**Address Prefix -** A string of 0 or more bits up to a maximum of 152 bits that is the lead portion of one or more ATM addresses.

**Address Resolution -** The procedure by which a client associates a LAN destination with the ATM address of another client or the BUS.

**Address Resolution Protocol (ARP) -** a method used to resolve higher level protocol addressing (such as IP) into the appropriate header data required for ATM; i.e., port, VPI, and VCI; also defines the AAL type to be used.

**Agent -** a component of network- and desktop-management software, such as SNMP, that gathers information from MIBs.

**alarm -** an unsolicited message from a device, typically indicating a problem with the system that requires attention.

Alarm Indication Signal (AIS) - In T1, an all ones condition used to alert a receiver that its incoming signal (or frame) has been lost. The loss of signal or frame is detected at the receiving end, and the failed signal is replaced by all the ones condition which the receiver interprets as an AIS. The normal response to this is AIS is for the receiving end to generate a yellow alarm signal as part of its transmission towards the faulty end. (The AIS itself is sometimes called a Blue Signal).

**A-Law -** The PCM coding and companding standard used in Europe.

Allowable Cell Rate (ACR) - parameter defined by the ATM Forum for ATM traffic management. ACR varies between the MCR and the PCR, and is dynamically controlled using congestion control mechanisms.

**Alternate Mark Inversion (AMI) -** A line coding format used on T1 facilities that transmits ones by alternate positive and negative pulses.

**Alternate Routing -** A mechanism that supports the use of a new path after an attempt to set up a connection along a previously selected path fails.

American National Standards Institute (ANSI) - a private organization that coordinates the setting and approval of some U.S. standards. It also represents the United States to the International Standards Organization.

American Standard Code for Information Interchange (ASCII) - a standard character set that (typically) assigns a 7-bit sequence to each letter, number, and selected control characters.

**AppleTalk** - a networking protocol developed by Apple Computer for communication between Apple's products and other computers. Independent of the network layer, AppleTalk runs on LocalTalk, EtherTalk and TokenTalk.

**Application Layer -** Layer seven of the ISO reference model; provides the end-user interface.

**Application Program (APP) -** a complete, self-contained program that performs a specific function directly for the user.

**Application Program Interface (API) -** a language format that defines how a program can be made to interact with another program, service, or other software; it allows users to develop custom interfaces with FORE products.

**Assigned Cell -** a cell that provides a service to an upper layer entity or ATM Layer Management entity (ATMM-entity).

**asxmon -** a FORE program that repeatedly displays the state of the switch and its active ports.

**Asynchronous Time Division Multiplexing (ATDM) -** a multiplexing technique in which a transmission capability is organized into a priori, unassigned time slots. The time slots are assigned to cells upon request of each application's instantaneous real need.

**Asynchronous Transfer Mode (ATM) -** a transfer mode in which the information is organized into cells. It is asynchronous in the sense that the recurrence of cells containing information from an individual user is not necessarily periodic.

**ATM Adaptation Layer (AAL) -** the AAL divides user information into segments suitable for packaging into a series of ATM cells. AAL layer types are used as follows:

**AAL-1** - constant bit rate, time-dependent traffic such as voice and video

AAL-2 - still undefined; a placeholder for variable bit rate video transmission

**AAL-3/4** - variable bit rate, delay-tolerant data traffic requiring some sequencing and/or error detection support (originally two AAL types, connection-oriented and connectionless, which have been combined)

**AAL-5 -** variable bit rate, delay-tolerant, connection-oriented data traffic requiring minimal sequencing or error detection support

ATM Address - Defined in the UNI Specification as 3 formats, each having 20 bytes in length.

**ATM Forum -** an international non-profit organization formed with the objective of accelerating the use of ATM products and services through a rapid convergence of interoperability specifications. In addition, the Forum promotes industry cooperation and awareness.

**ATM Inverse Multiplexing (AIMUX) -** A device that allows multiple T1 or E1 communications facilities to be combined into a single broadband facility for the transmission of ATM cells.

**ATM Layer link -** a section of an ATM Layer connection between two adjacent active ATM Layer entities (ATM-entities).

**ATM Link -** a virtual path link (VPL) or a virtual channel link (VCL).

**ATM Management Interface (AMI) -** the user interface to FORE Systems' *ForeThought* switch control software (SCS). AMI lets users monitor and change various operating configurations of FORE Systems switches and network module hardware and software, IP connectivity, and SNMP network management.

**ATM Peer-to-Peer Connection -** a virtual channel connection (VCC) or a virtual path connection (VPC) directly established, such as workstation-to-workstation. This setup is not commonly used in networks.

**ATM Traffic Descriptor -** a generic list of parameters that can be used to capture the intrinsic traffic characteristics of a requested ATM connection.

**ATM User-to-User Connection -** an association established by the ATM Layer to support communication between two or more ATM service users (i.e., between two or more next higher layer entities or between two or more ATM entities). The communication over an ATM Layer connection may be either bidirectional or unidirectional. The same Virtual Channel Identifier (VCI) is used for both directions of a connection at an interface.

**atmarp -** a FORE program that shows and manipulates ATM ARP entries maintained by the given device driver. This is also used to establish PVC connections.

**ATM-attached Host Functional Group (AHFG) -** The group of functions performed by an ATM-attached host that is participating in the MPOA service.

atmconfig - a FORE program used to enable or disable SPANS signalling.

**atmstat -** a FORE program that shows statistics gathered about a given adapter card by the device driver. These statistics include ATM layer and ATM adaptation layer cell and error counts. This can also be used to query other hosts via SNMP.

Attachment User Interface (AUI) - IEEE 802.3 interface between a media attachment unit (MAU) and a network interface card (NIC). The term AUI can also refer to the rear panel port to which an AUI cable might attach.

**Auto-logout -** a feature that automatically logs out a user if there has been no user interface activity for a specified length of time.

**Automatic Protection Switching (APS) -** Equipment installed in communications systems to detect circuit failures and automatically switch to redundant, standby equipment.

**Available Bit Rate (ABR) -** a type of traffic for which the ATM network attempts to meet that traffic's bandwidth requirements. It does not guarantee a specific amount of bandwidth and the end station must retransmit any information that did not reach the far end.

**Backbone** - the main connectivity device of a distributed system. All systems that have connectivity to the backbone connect to each other, but systems can set up private arrangements with each other to bypass the backbone to improve cost, performance, or security.

**Backplane -** High-speed communications line to which individual components are connected.

**Backward Explicit Congestion Notification (BECN) -** A Resource Management cell type generated by the network or the destination, indicating congestion or approaching congestion for traffic flowing in the direction opposite that of the BECN cell.

**Bandwidth -** usually identifies the capacity or amount of data that can be sent through a given circuit; may be user-specified in a PVC.

**Baud -** unit of signalling speed, equal to the number of discrete conditions or signal events per second. If each signal event represents only one bit, the baud rate is the same as bps; if each signal event represents more than one bit (such as a dibit), the baud rate is smaller than bps.

**Bayonet-Neill-Concelman (BNC) -** a bayonet-locking connector used to terminate coaxial cables. BNC is also referred to as Bayonet Network Connector.

**Bipolar 8 Zero Substitution (B8ZS)** - a technique used to satisfy the ones density requirements of digital T-carrier facilities in the public network while allowing 64 Kbps clear channel data. Strings of eight consecutive zeroes are replaced by an eight-bit code representing two intentional bipolar pulse code violations (000V10V1).

**Bipolar Violation (BPV) -** an error event on a line in which the normal pattern of alternating high (one) and low (zero) signals is disrupted. A bipolar violation is noted when two high signals occur without an intervening low signal, or vice versa.

**B-ISDN Inter-Carrier Interface (B-ICI) -** An ATM Forum defined specification for the interface between public ATM networks to support user services across multiple public carriers.

Bit Error Rate (BER) - A measure of transmission quality, generally shown as a negative exponent, (e.g.,  $10^{-7}$  which means 1 out of  $10^7$  bits [1 out of 10,000,000 bits] are in error).

Bit Interleaved Parity (BIP) - an error-detection technique in which character bit patterns are forced into parity, so that the total number of one bits is always odd or always even. This is accomplished by the addition of a one or zero bit to each byte, as the byte is transmitted; at the other end of the transmission, the receiving device verifies the parity (odd or even) and the accuracy of the transmission.

Bit Robbing - The use of the least significant bit per channel in every sixth frame for signaling.

**Bit Stuffing -** A process in bit-oriented protocols where a zero is inserted into a string of ones by the sender to prevent the receiver from interpreting valid user data (the string of ones) as control characters (a Flag character for instance).

**Border Gateway Protocol (BGP) -** used by gateways in an internet connecting autonomous networks. It is derived from experiences learned using the EGP.

**bps** - bits per second

**Bridge -** a device that expands a Local Area Network by forwarding frames between data link layers associated with two separate cables, usually carrying a common protocol. Bridges can usually be made to filter certain packets (to forward only certain traffic).

**Bridge Protocol Data Unit (BPDU) -** A message type used by bridges to exchange management and control information.

**Broadband -** a service or system requiring transmission channels capable of supporting rates greater than the Integrated Services Digital Network (ISDN) primary rate.

**Broadband Access -** an ISDN access capable of supporting one or more broadband services.

**Broadband Connection Oriented Bearer (BCOB) -** Information in the SETUP message that indicates the type of service requested by the calling user.

**BCOB-A (Bearer Class A) -** Indicated by ATM end user in SETUP message for connection-oriented, constant bit rate service. The network may perform internetworking based on AAL information element (IE).

**BCOB-C (Bearer Class C) -** Indicated by ATM end user in SETUP message for connection-oriented, variable bit rate service. The network may perform internetworking based on AAL information element (IE).

**BCOB-X (Bearer Class X) -** Indicated by ATM end user in SETUP message for ATM transport service where AAL, traffic type and timing requirements are transparent to the network.

**Broadband Integrated Services Digital Network (B-ISDN) -** a common digital network suitable for voice, video, and high-speed data services running at rates beginning at 155 Mbps.

**Broadband ISDN User's Part (B-ISUP) -** A protocol used to establish, maintain and release broadband switched network connections across an SS7/ATM network.

**Broadband Terminal Equipment (B-TE) -** An equipment category for B-ISDN which includes terminal adapters and terminals.

**Broadcast** - Data transmission to all addresses or functions.

Broadcast and Unknown Server (BUS) - in an emulated LAN, the BUS is responsible for accepting broadcast, multicast, and unknown unicast packets from the LECs to the broadcast MAC address (FFFFFFFFFF) via dedicated point-to-point connections, and forwarding the packets to all of the members of the ELAN using a single point-to-multipoint connection.

**Brouter (bridging/router) -** a device that routes some protocols and bridges others based on configuration information.

**Buffer -** A data storage medium used to compensate of a difference in rate of data flow or time of occurrence of events when transmitting data from one device to another.

**Building Integrated Timing Supply (BITS)** - a master timing supply for an entire building, which is a master clock and its ancillary equipment. The BITS supplies DS1 and/or composite clock timing references for synchronization to all other clocks and timing sources in that building.

**Bursty Errored Seconds (BES)** - a BES contains more than 1 and fewer than 320 path coding violation error events, and no severely errored frame or AIS defects. Controlled slips are not included in determining BESs.

**Bursty Second -** a second during which there were at least the set number of BES threshold event errors but fewer than the set number of SES threshold event errors.

Byte - A computer-readable group of bits (normally 8 bits in length).

**Call -** an association between two or more users or between a user and a network entity that is established by the use of network capabilities. This association may have zero or more connections.

**Carrier -** a company, such as any of the "baby Bell" companies, that provide network communications services, either within a local area or between local areas.

Carrier Group Alarm (CGA) - A service alarm generated by a channel bank when an out-of-frame (OOF) condition exists for some predetermined length of time (generally 300 milliseconds to 2.5 seconds). The alarm causes the calls using a trunk to be dropped and trunk conditioning to be applied.

**Carrier Identification Parameter (CIP) -** A 3 or 4 digit code in the initial address message identifying the carrier to be used for the connection.

**cchan -** a FORE program that manages virtual channels on a *ForeRunner* switch running asxd.

**Cell -** an ATM Layer protocol data unit (PDU). The basic unit of information transported in ATM technology, each 53-byte cell contains a 5-byte header and a 48-byte payload.

**Cell Delay Variation (CDV) -** a quantification of cell clumping for a connection. The cell clumping CDV (yk) is defined as the difference between a cell's expected reference arrival time (ck) and its actual arrival time (ak). The expected reference arrival time (ck) of cell k of a specific connection is max. T is the reciprocal of the negotiated peak cell rate.

**Cell Delineation** - the protocol for recognizing the beginning and end of ATM cells within the raw serial bit stream.

**Cell Header -** ATM Layer protocol control information.

**Cell Loss Priority (CLP) -** the last bit of byte four in an ATM cell header; indicates the eligibility of the cell for discard by the network under congested conditions. If the bit is set to 1, the cell may be discarded by the network depending on traffic conditions.

**Cell Loss Ratio -** In a network, cell loss ratio is (1-x/y), where y is the number of cells that arrive in an interval at an ingress of the network; and x is the number of these y cells that leave at the egress of the network element.

**Cell Loss Ratio (CLR) -** CLR is a negotiated QoS parameter and acceptable values are network specific. The objective is to minimize CLR provided the end-system adapts the traffic to the changing ATM layer transfer characteristics. The Cell Loss Ratio is defined for a connection as: Lost Cells/Total Transmitted Cells. The CLR parameter is the value of CLR that the network agrees to offer as an objective over the lifetime of the connection. It is expressed as an order of magnitude, having a range of 10-1 to 10-15 and unspecified.

**Cell Misinsertion Rate (CMR) -** the ratio of cells received at an endpoint that were not originally transmitted by the source end in relation to the total number of cells properly transmitted.

**Cell Rate Adaptation (CRA)** - a function performed by a protocol module in which empty cells (known as unassigned cells) are added to the output stream. This is because there always must be a fixed number of cells in the output direction; when there are not enough cells to transmit, unassigned cells are added to the output data stream.

**Cell Relay Service (CRS)** - a carrier service which supports the receipt and transmission of ATM cells between end users in compliance with ATM standards and implementation specifications.

**Cell Transfer Delay -** the transit delay of an ATM cell successfully passed between two designated boundaries. See CTD.

**Cell Transfer Delay (CTD) -** This is defined as the elapsed time between a cell exit event at the measurement point 1 (e.g., at the source UNI) and the corresponding cell entry event at the measurement point 2 (e.g., the destination UNI) for a particular connection. The cell transfer delay between two measurement points is the sum of the total inter-ATM node transmission delay and the total ATM node processing delay.

**Channel -** A path or circuit along which information flows.

**Channel Associated Signaling (CAS) -** a form of circuit state signaling in which the circuit state is indicated by one or more bits of signaling status sent repetitively and associated with that specific circuit.

**Channel Bank -** A device that multiplexes many slow speed voice or data conversations onto high speed link and controls the flow.

**Channel Service Unit (CSU) -** An interface for digital leased lines which performs loopback testing and line conditioning.

**Channelization -** capability of transmitting independent signals together over a cable while still maintaining their separate identity for later separation.

**Circuit -** A communications link between points.

**Circuit Emulation Service (CES) -** The ATM Forum circuit emulation service interoperability specification specifies interoperability agreements for supporting Constant Bit Rate (CBR) traffic over ATM networks that comply with the other ATM Forum interoperability agreements. Specifically, this specification supports emulation of existing TDM circuits over ATM networks.

Classical IP (CLIP) - IP over ATM which conforms to RFC 1577.

**Clear to Send (CTS) -** and RS-232 modem interface control signal (sent from the modem to the DTE on pin 5) which indicates that the attached DTE may begin transmitting; issuance in response to the DTE's RTS.

**Clocking -** Regularly timed impulses.

**Closed User Group -** A subgroup of network users that can be its own entity; any member of the subgroup can only communicate with other members of that subgroup.

**Coaxial Cable -** Coax is a type of electrical communications medium used in the LAN environment. This cable consists of an outer conductor concentric to an inner conductor, separated from each other by insulating material, and covered by some protective outer material. This medium offers large bandwidth, supporting high data rates with high immunity to electrical interference and a low incidence of errors. Coax is subject to distance limitations and is relatively expensive and difficult to install.

**Cold Start Trap -** an SNMP trap which is sent after a power-cycle (see *trap*).

**Collision -** Overlapping transmissions that occur when two or more nodes on a LAN attempt to transmit at or about the same time.

**Committed Information Rate (CIR) -** CIR is the information transfer rate which a network offering Frame Relay Services (FRS) is committed to transfer under normal conditions. The rate is averaged over a minimum increment of time.

**Common Channel Signaling (CCS) -** A form signaling in which a group of circuits share a signaling channel. Refer to SS7.

**Common Management Interface Protocol (CMIP) -** An ITU-TSS standard for the message formats and procedures used to exchange management information in order to operate, administer maintain and provision a network.

**Concatenation -** The connection of transmission channels similar to a chain.

**Concentrator -** a communications device that offers the ability to concentrate many lower-speed channels into and out of one or more high-speed channels.

**Configuration -** The phase in which the LE Client discovers the LE Service.

**Congestion Management -** traffic management feature that helps ensure reasonable service for VBR connections in an ATM network, based on a priority, sustained cell rate (SCR), and peak cell rate (PCR). During times of congestion, bandwidth is reduced to the SCR, based on the priority of the connection.

**Connection -** the concatenation of ATM Layer links in order to provide an end-to-end information transfer capability to access points.

**Connection Admission Control (CAC) -** the procedure used to decide if a request for an ATM connection can be accepted based on the attributes of both the requested connection and the existing connections.

Connection Endpoint (CE) - a terminator at one end of a layer connection within a SAP.

**Connection Endpoint Identifier (CEI) -** an identifier of a CE that can be used to identify the connection at a SAP.

**Connectionless Broadband Data Service (CBDS) -** A connectionless service similar to Bellcore's SMDS defined by European Telecommunications Standards Institute (ETSI).

**Connectionless Service -** a type of service in which no pre-determined path or link has been established for transfer of information, supported by AAL 4.

**Connectionless Service (CLS) -** A service which allows the transfer of information among service subscribers without the need for end-to- end establishment procedures.

**Connection-Oriented Service -** a type of service in which information always traverses the same pre-established path or link between two points, supported by AAL 3.

**Constant Bit Rate (CBR) -** a type of traffic that requires a continuous, specific amount of bandwidth over the ATM network (e.g., digital information such as video and digitized voice).

**Controlled Slip (CS)** - a situation in which one frame's worth of data is either lost or replicated. A controlled slip typically occurs when the sending device and receiving device are not using the same clock.

**Convergence Sublayer (CS)** - a portion of the AAL. Data is passed first to the CS where it is divided into rational, fixed-length packets or PDUs (Protocol Data Units). For example, AAL 4 processes user data into blocks that are a maximum of 64 kbytes long.

**Corresponding Entities -** peer entities with a lower layer connection among them.

cpath - a FORE program used to manage virtual paths on a ForeRunner switch running asxd.

**cport -** a FORE program that monitors and changes the state of ports on a *ForeRunner* switch running asxd.

**Cross Connection -** a mapping between two channels or paths at a network device.

**Customer Premise Equipment (CPE) -** equipment that is on the customer side of the point of demarcation, as opposed to equipment that is on a carrier side. See also point of demarcation.

**Cut Through** - Establishment of a complete path for signaling and/or audio communications.

**Cyclic Redundancy Check (CRC)** - an error detection scheme in which a number is derived from the data that will be transmitted. By recalculating the CRC at the remote end and comparing it to the value originally transmitted, the receiving node can detect errors.

D3/D4 - Refers to compliance with AT&T TR (Technical Reference) 62411 definitions for coding, supervision, and alarm support. D3/D4 compatibility ensures support of digital PBXes, M24 services, Megacom services, and Mode 3 D3/D4 channel banks at DS-1 level.

**D4 Channelization -** refers to compliance with AT&T Technical Reference 62411 regarding DS1 frame layout (the sequential assignment of channels and time slot numbers within the DS1).

**D4** Framed/Framing Format - in T1, a 193-bit frame format in which the 193rd bit is used for framing and signaling information (the frame/framing bit). To be considered in support of D4 Framing, a device must be able to synchronize and frame-up on the 193rd bit.

**Data Communications Equipment (DCE) -** a definition in the RS232C standard that describes the functions of the signals and the physical characteristics of an interface for a communication device such as a modem.

**Data Country Code (DCC) -** This specifies the country in which an address is registered. The codes are given in ISO 3166. The length of this field is two octets. The digits of the data country code are encoded in Binary Coded Decimal (BCD) syntax. The codes will be left justified and padded on the right with the hexadecimal value "F" to fill the two octets.

**Data Link** - Communications connection used to transmit data from a source to a destination.

**Data Link Connection Identifier (DLCI) -** connection identifier associated with frame relay packets that serves the same functions as, and translates directly to, the VPI/VCI on an ATM cell.

**Data Link Layer -** Layer 2 of the OSI model, responsible for encoding data and passing it to the physical medium. The IEEE divides this layer into the LLC (Logical Link Control) and MAC (Media Access Control) sublayers.

**Data Set Ready (DSR) -** an RS-232 modem interface control signal (sent from the modem to the DTE on pin 6) which indicates that the modem is connected to the telephone circuit. Usually a prerequisite to the DTE issuing RTS.

**Data Terminal Equipment (DTE) -** generally user devices, such as terminals and computers, that connect to data circuit-terminating equipment. They either generate or receive the data carried by the network.

**Data Terminal Ready (DTR) -** an RS232 modem interface control signal (sent from the DTE to the modem on pin 20) which indicates that the DTE is ready for data transmission and which requests that the modem be connected to the telephone circuit.

**Datagram -** a packet of information used in a connectionless network service that is routed to its destination using an address included in the datagram's header.

**DECnet -** Digital Equipment Corporation's proprietary LAN.

**Defense Advanced Research Projects Agency (DARPA) -** the US government agency that funded the ARPANET.

**Demultiplexing -** a function performed by a layer entity that identifies and separates SDUs from a single connection to more than one connection (see *multiplexing*).

**Destination End Station (DES) -** An ATM termination point which is the destination for ATM messages of a connection and is used as a reference point for ABR services. See SES.

**Digital Access and Cross-Connect System (DACS) -** Digital switching system for routing T1 lines, and DS-0 portions of lines, among multiple T1 ports.

**Digital Cross-connect System (DCS)** - an electronic patch panel used to route digital signals in a central office.

**Digital Standard n (0, 1, 1C, 2, and 3) (DSn) -** a method defining the rate and format of digital hierarchy, with asynchronous data rates defined as follows:

DS0	64kb/s	1 voice channel
DS1	1.544Mb/s	24 DS0s
DS1C	3.152 Mb/s	2 DS1s
DS2	6.312 Mb/s	4 DS1s
DS3	44.736 Mb/s	28 DS1s

Synchronous data rates (SONET) are defined as:

STS-1/OC-1	51.84 Mb/s	28 DS1s or 1 DS3
STS-3/OC-3	155.52 Mb/s	3 STS-1s byte interleaved
STS-3c/OC-3c	155.52 Mb/s	Concatenated, indivisible payload
STS-12/OC-12	622.08 Mb/s	12 STS-1s, 4 STS-3cs, or any mixture
STS-12c/OC-12c	622.08 Mb/s	Concatenated, indivisible payload
STS-48/OC-48	2488.32 Mb/s	48 STS-1s, 16 STS-3cs, or any mixture

**DIP (Dual In-line Package) Switch -** a device that has two parallel rows of contacts that let the user switch electrical current through a pair of those contacts to on or off. They are used to reconfigure components and peripherals.

**Domain Name Server -** a computer that converts names to their corresponding Internet numbers. It allows users to telnet or FTP to the name instead of the number.

**Domain Naming System (DNS) -** the distributed name and address mechanism used in the Internet.

**Duplex -** Two way communication.

**DXI -** a generic phrase used in the full names of several protocols, all commonly used to allow a pair of DCE and DTE devices to share the implementation of a particular WAN protocol. The protocols define the packet formats used to transport data between DCE and DTE devices.

**DXI Frame Address (DFA) -** a connection identifier associated with ATM DXI packets that serves the same functions as, and translates directly to, the VPI/VCI on an ATM cell.

**Dynamic Allocation -** A technique in which the resources assigned for program execution are determined by criteria applied at the moment of need.

**E.164** - A public network addressing standard utilizing up to a maximum of 15 digits. ATM uses E.164 addressing for public network addressing.

**E1 -** Wide-area digital transmission scheme used predominantly in Europe that carries data at a rate of 2.048 Mbps. E1 lines can be leased for private use from common carriers.

**E3 -** Wide-area digital transmission scheme used predominantly in Europe that carries data at a rate of 34.368 Mbps. E3 lines can be leased for private use from common carriers.

**Edge Device -** A physical device which is capable of forwarding packets between legacy interworking interfaces (e.g., Ethernet, Token Ring, etc.) and ATM interfaces based on data-link and network layer information but which does not participate in the running of any network layer routing protocol. An Edge Device obtains forwarding descriptions using the route distribution protocol.

**elarp -** a FORE program that shows and manipulates MAC and ATM address mappings for LAN Emulation Clients (LECs).

**elconfig** - a FORE program that shows and modifies LEC configuration. Lets the user set the NSAP address of the LAN Emulation Configuration Server, display the list of Emulated LANs configured in the LECS for this host, display the list of ELANs locally configured along with the membership state of each, and locally administer ELAN membership.

**Electrically Erasable Programmable Read Only Memory (EEPROM) -** an EPROM that can be cleared with electrical signals rather than the traditional ultraviolet light.

**Electromagnetic Interference (EMI) -** signals generated and radiated by an electronic device that cause interference with radio communications, among other effects.

**Electronics Industries Association (EIA) -** a USA trade organization that issues its own standards and contributes to ANSI; developed RS-232. Membership includes USA manufacturers.

**Embedded SNMP Agent -** an SNMP agent can come in two forms: embedded or proxy. An embedded SNMP agent is integrated into the physical hardware and software of the unit.

**Emulated Local Area Network (ELAN) -** A logical network initiated by using the mechanisms defined by LAN Emulation. This could include ATM and legacy attached end stations.

**End System (ES)** - a system where an ATM connection is terminated or initiated (an originating end system initiates the connection).

**End System Identifier (ESI) -** This identifier distinguishes multiple nodes at the same level in case the lower level peer group is partitioned.

**End-to-End Connection -** when used in reference to an ATM network, a connection that travels through an ATM network, passing through various ATM devices and with endpoints at the termination of the ATM network.

**Enterprise -** Terminology generally referring to customers with multiple, non-contiguous geographic locations.

**Equalization (EQL)** - the process of compensating for line distortions.

**Erasable Programmable Read Only Memory (EPROM) -** A PROM which may be erased and rewritten to perform new or different functions (normally done with a PROM burner).

**Errored Second (ES)** - a second during which at least one code violation occurred.

**Ethernet -** a 10-Mbps, coaxial standard for LANs in which all nodes connect to the cable where they contend for access.

**Excessive Zeroes (EXZ) Error Event -** An Excessive Zeroes error event for an AMI-coded signal is the occurrence of more than fifteen contiguous zeroes. For a B8ZS coded signal, the defect occurs when more than seven contiguous zeroes are detected.

**Explicit Forward Congestion Indication (EFCI)** - the second bit of the payload type field in the header of an ATM cell, the EFCI bit indicates network congestion to receiving hosts. On a congested switch, the EFCI bit is set to "1" by the transmitting network module when a certain number of cells have accumulated in the network module's shared memory buffer. When a cell is received that has its EFCI bit set to "1," the receiving host notifies the sending host, which should then reduce its transmission rate.

**Explicit Rate (ER) -** The Explicit Rate is an RM-cell field used to limit the source ACR to a specific value. It is initially set by the source to a requested rate (such as PCR). It may be subsequently reduced by any network element in the path to a value that the element can sustain. ER is formatted as a rate.

**Extended Industry Standard Architecture (EISA) -** bus architecture for desktop computers that provides a 32-bit data passage and maintains compatibility with the ISA or AT architecture.

**Extended Super Frame (ESF)** - a T1 framing format that utilizes the 193rd bit as a framing bit, but whose Superframe is made up of 24 frames instead of 12 as in D4 format. ESF also provides CRC error detection and maintenance data link functions.

**Exterior Gateway Protocol (EGP) -** used by gateways in an internet, connecting autonomous networks.

**Fairness** - related to Generic Flow Control, fairness is defined as meeting all of the agreed quality of service requirements by controlling the order of service for all active connections.

**Far End Block Error (FEBE)** - an error detected by extracting the 4-bit FEBE field from the path status byte (G1). The legal range for the 4-bit field is between 0000 and 1000, representing zero to eight errors. Any other value is interpreted as zero errors.

**Far End Receive Failure (FERF)** - a line error asserted when a 110 binary pattern is detected in bits 6, 7, 8 of the K2 byte for five consecutive frames. A line FERF is removed when any pattern other than 110 is detected in these bits for five consecutive frames.

**Far-End** - in a relationship between two devices in a circuit, the far-end device is the one that is remote.

**Face Contact (FC) -** Designation for fiber optic connector designed by Nippon Telegraph and Telephone which features a movable anti-rotation key allowing good repeatable performance despite numerous mating. Normally referred to as Fiber Connector, FC actually stands for Face Contact and sometimes linked with PC (Point Contact), designated as FC or FC-PC.

**FCC Part 68 -** The FCC rules regulating the direct connection of non-telephone company provided equipment to the public telephone network.

**Federal Communications Commission (FCC) -** a board of commissioners appointed by the President under the Communications Act of 1934, with the authority to regulate all interstate telecommunications originating in the United States, including transmission over phone lines.

**Fiber Distributed Data Interface (FDDI) -** high-speed data network that uses fiber-optic as the physical medium. Operates in similar manner to Ethernet or Token Ring, only faster.

**File Transfer Protocol (FTP) -** a TCP/IP protocol that lets a user on one computer access, and transfer data to and from, another computer over a network. ftp is usually the name of the program the user invokes to accomplish this task.

First-In, First-Out (FIFO) - method of coordinating the sequential flow of data through a buffer.

**Flag -** a bit pattern of six binary "1"s bounded by a binary "0" at each end (forms a 0111 1110 or Hex "7E"). It is used to mark the beginning and/or end of a frame.

**Flow Control -** The way in which information is controlled in a network to prevent loss of data when the receiving buffer is near its capacity.

**ForeThought PNNI (FT-PNNI) -** a FORE Systems routing and signalling protocol that uses private ATM (NSAP) addresses; a precursor to ATM Forum PNNI (see PNNI).

**Forward Error Correction (FEC)** - A technique used by a receiver for correcting errors incurred in transmission over a communications channel without requiring retransmission of any information by the transmitter; typically involves a convolution of the transmitted bits and the appending of extra bits by both the receiver and transmitter using a common algorithm.

**Forward Explicit Congestion Notification (FECN)** - Bit set by a Frame Relay network to inform data terminal equipment (DTE) receiving the frame that congestion was experienced in the path from source to destination. DTE receiving frames with the FECN bit set can request that higher-level protocols take flow control action as appropriate.

Fractional T1 - the use of bandwidth in 64Kbps increments up to 1.544Mbps from a T1 facility.

**Frame -** a variable length group of data bits with a specific format containing flags at the beginning and end to provide demarcation.

**Frame Check Sequence (FCS) -** In bit-oriented protocols, a 16-bit field that contains transmission error checking information, usually appended to the end of the frame.

**Frame Relay -** a fast packet switching protocol based on the LAPD protocol of ISDN that performs routing and transfer with less overhead processing than X.25.

Frame Synchronization Error - an error in which one or more time slot framing bits are in error.

**Frame-Based UNI (FUNI) -** An ATM switch-based interface which accepts frame-based ATM traffic and converts it into cells.

**Frame-Relay Service (FRS) -** A connection oriented service that is capable of carrying up to 4096 bytes per frame.

**Framing -** a protocol that separates incoming bits into identifiable groups so that the receiving multiplexer recognizes the grouping.

**Frequency Division Multiplexing (FDM) -** a method of dividing an available frequency range into parts with each having enough bandwidth to carry one channel.

**Gbps** - gigabits per second (billion)

**Generic Cell Rate Algorithm (GCRA) -** an algorithm which is employed in traffic policing and is part of the user/network service contract. The GCRA is a scheduling algorithm which ensures that cells are marked as conforming when they arrive when expected or later than expected and non-conforming when they arrive sooner than expected.

**Generic Connection Admission Control (GCAC) -** This is a process to determine if a link has potentially enough resources to support a connection.

Generic Flow Control (GFC) - the first four bits of the first byte in an ATM cell header. Used to control the flow of traffic across the User-to-Network Interface (UNI), and thus into the network. Exact mechanisms for flow control are still under investigation and no explicit definition for this field exists at this time. (This field is used only at the UNI; for NNI-NNI use (between network nodes), these four bits provide additional network address capacity, and are appended to the VPI field.)

**GIO** - a proprietary bus architecture used in certain Silicon Graphics, Inc. workstations.

**Header -** protocol control information located at the beginning of a protocol data unit.

**Header Error Control (HEC) -** a CRC code located in the last byte of an ATM cell header that is used for checking cell header integrity only.

**High Density Bipolar (HDB3) -** A bipolar coding method that does not allow more than 3 consecutive zeroes.

**High Level Data Link Control (HDLC) -** An ITU-TSS link layer protocol standard for point-to-point and multi-point communications.

**High Performance Parallel Interface (HIPPI) -** ANSI standard that extends the computer bus over fairly short distances at speeds of 800 and 1600 Mbps.

**High-Speed Serial Interface (HSSI) -** a serial communications connection that operates at speeds of up to 1.544 Mbps.

Host - In a network, the primary or controlling computer in a multiple computer installation.

**HPUX -** the Hewlett-Packard version of UNIX.

**Hub** - a device that connects several other devices, usually in a star topology.

**I/O Module -** FORE's interface cards for the LAX-20 LAN Access Switch, designed to connect Ethernet, Token Ring, and FDDI LANs to *ForeRunner* ATM networks.

**Institute of Electrical and Electronics Engineers (IEEE) -** the world's largest technical professional society. Based in the U.S., the IEEE sponsors technical conferences, symposia & local meetings worldwide, publishes nearly 25% of the world's technical papers in electrical, electronics & computer engineering, provides educational programs for members, and promotes standardization.

**IEEE 802 -** Standards for the interconnection of LAN computer equipment. Deals with the Data Link Layers of the ISO Reference Model for OSI.

**IEEE 802.1 -** Defines the high-level network interfaces such as architecture, internetworking and network management.

**IEEE 802.2 -** Defines the Logical Link Control interface between the Data Link and Network Layers.

**IEEE 802.3 -** Defines CSMA/CD (Ethernet).

**IEEE 802.4 -** Defines the token-passing bus.

**IEEE 802.5 -** Defines the Token Ring access methodology. This standard incorporates IBM's Token Ring specifications.

**IEEE 802.6 -** Defines Metropolitan Area Networks.

**IEEE 802.7 -** The broadband technical advisory group.

**IEEE 802.9 -** Defines integrated data and voice networks.

**Integrated Services Digital Network (ISDN) -** an emerging technology that is beginning to be offered by the telephone carriers of the world. ISDN combines voice and digital network services into a single medium or wire.

**Interexchange Carriers (IXC) -** Long-distance communications companies that provide service between Local Access Transport Areas (LATAs).

**Interface Data -** the unit of information transferred to/from the upper layer in a single interaction across a SAP. Each Interface Data Unit (IDU) controls interface information and may also contain the whole or part of the SDU.

**Interface Data Unit (IDU) -** The unit of information transferred to/from the upper layer in a single interaction across the SAP. Each IDU contains interface control information and may also contain the whole or part of the SDU.

**Interim Local Management Interface (ILMI) -** the standard that specifies the use of the Simple Network Management Protocol (SNMP) and an ATM management information base (MIB) to provide network status and configuration information.

**Intermediate System (IS)** - a system that provides forwarding functions or relaying functions or both for a specific ATM connection. OAM cells may be generated and received.

**International Standards Organization (ISO)** - a voluntary, non treaty organization founded in 1946 that is responsible for creating international standards in many areas, including computers and communications.

**International Telephone and Telegraph Consultative Committee (CCITT) -** the international standards body for telecommunications.

**Internet -** (note the capital "I") the largest internet in the world including large national backbone nets and many regional and local networks worldwide. The Internet uses the TCP/IP suite. Networks with only e-mail connectivity are not considered on the Internet.

**internet -** while an internet is a network, the term "internet" is usually used to refer to a collection of networks interconnected with routers.

Internet Addresses - the numbers used to identify hosts on an internet network. Internet host numbers are divided into two parts; the first is the network number and the second, or local, part is a host number on that particular network. There are also three classes of networks in the Internet, based on the number of hosts on a given network. Large networks are classified as Class A, having addresses in the range 1-126 and having a maximum of 16,387,064 hosts. Medium networks are classified as Class B, with addresses in the range 128-191 and with a maximum of 64,516 hosts. Small networks are classified as Class C, having addresses in the range 192-254 with a maximum of 254 hosts. Addresses are given as dotted decimal numbers in the following format:

## nnn.nnn.nnn.nnn

In a Class A network, the first of the numbers is the network number, the last three numbers are the local host address.

In a Class B network, the first two numbers are the network, the last two are the local host address.

In a Class C network, the first three numbers are the network address, the last number is the local host address.

The following table summarizes the classes and sizes:

Class	First #	Max# Hosts
A	1-126	16,387,064
В	129-191	64,516
С	192-223	254

## Glossary

Network mask values are used to identify the network portion and the host portion of the address. Default network masks are as follows:

Class A - 255.0.0.0

Class B - 255.255.0.0

Class C - 255,255,255.0

Subnet masking is used when a portion of the host ID is used to identify a subnetwork. For example, if a portion of a Class B network address is used for a subnetwork, the mask could be set as 255.255.255.0. This would allow the third byte to be used as a subnetwork address. All hosts on the network would still use the IP address to get on the Internet.

**Internet Control Message Protocol (ICMP)** - the protocol that handles errors and control messages at the IP layer. ICMP is actually a part of the IP protocol layer. It can generate error messages, test packets, and informational messages related to IP.

**Internet Engineering Task Force (IETF) -** a large, open, international community of network designers, operators, vendors and researchers whose purpose is to coordinate the operation, management and evolution of the Internet to resolve short- and mid-range protocol and architectural issues.

**Internet Protocol (IP) -** a connectionless, best-effort packet switching protocol that offers a common layer over dissimilar networks.

**Internetwork Packet Exchange (IPX) Protocol -** a NetWare protocol similar to the Xerox Network Systems (XNS) protocol that provides datagram delivery of messages.

**Interoperability -** The ability of software and hardware on multiple machines, from multiple vendors, to communicate.

**Interworking Function (IWF) -** provides a means for two different technologies to interoperate.

**IP Address -** a unique 32-bit integer used to identify a device in an IP network. You will most commonly see IP addresses written in "dot" notation (e.g., 192.228.32.14).

**IP Netmask -** a 32-bit pattern that is combined with an IP address to determine which bits of an IP address denote the network number and which denote the host number. Netmasks are useful for sub-dividing IP networks. IP netmasks are written in "dot" notation (e.g., 255.255.0.0).

**ISA Bus -** a bus standard developed by IBM for expansion cards in the first IBM PC. The original bus supported a data path only 8 bits wide. IBM subsequently developed a 16-bit version for its AT class computers. The 16-bit AT ISA bus supports both 8- and 16-bit cards. The 8-bit bus is commonly called the PC/XT bus, and the 16-bit bus is called the AT bus.

**Isochronous -** signals carrying embedded timing information or signals that are dependent on uniform timing; usually associated with voice and/or video transmission.

International Telecommunications Union Telecommunications (ITU-T) - an international body of member countries whose task is to define recommendations and standards relating to the international telecommunications industry. The fundamental standards for ATM have been defined and published by the ITU-T (Previously CCITT).

**J2 -** Wide-area digital transmission scheme used predominantly in Japan that carries data at a rate of 6.312 Mbps.

**Jitter -** analog communication line distortion caused by variations of a signal from its reference timing position.

**Joint Photographic Experts Group (JPEG) -** An ISO Standards group that defines how to compress still pictures.

**Jumper -** a patch cable or wire used to establish a circuit, often temporarily, for testing or diagnostics; also, the devices, shorting blocks, used to connect adjacent exposed pins on a printed circuit board that control the functionality of the card.

**Kbps** - kilobits per second (thousand)

**LAN Access Concentrator -** a LAN access device that allows a shared transmission medium to accommodate more data sources than there are channels currently available within the transmission medium.

**LAN Emulation Address Resolution Protocol (LE\_ARP) -** A message issued by a LE client to solicit the ATM address of another function.

**LAN Emulation Client (LEC) -** the component in an end system that performs data forwarding, address resolution, and other control functions when communicating with other components within an ELAN.

**LAN Emulation Configuration Server (LECS) -** the LECS is responsible for the initial configuration of LECs. It provides information about available ELANs that a LEC may join, together with the addresses of the LES and BUS associated with each ELAN.

**LAN Emulation Server (LES) -** the LES implements the control coordination function for an ELAN by registering and resolving MAC addresses to ATM addresses.

**LAN Emulation (LANE) -** technology that allows an ATM network to function as a LAN backbone. The ATM network must provide multicast and broadcast support, address mapping (MAC-to-ATM), SVC management, and a usable packet format. LANE also defines Ethernet and Token Ring ELANs.

lane - a program that provides control over the execution of the LAN Emulation Server (LES), Broadcast/Unknown Server (BUS), and LAN Emulation Configuration Server (LECS) on the local host.

**Latency -** The time interval between a network station seeking access to a transmission channel and that access being granted or received.

Layer Entity - an active layer within an element.

**Layer Function -** a part of the activity of the layer entities.

**Layer Service -** a capability of a layer and the layers beneath it that is provided to the upper layer entities at the boundary between that layer and the next higher layer.

**Layer User Data -** the information transferred between corresponding entities on behalf of the upper layer or layer management entities for which they are providing services.

**le -** a FORE program that implements both the LAN Emulation Server (LES) and the Broadcast/Unknown Server (BUS).

**Leaky Bucket -** informal cell policing term for the Generic Cell Rate Algorithm which in effect receives cells into a bucket and leaks them out at the specified or contracted rate (i.e., PCR).

**Least Significant Bit (LSB) -** lowest order bit in the binary representation of a numerical value.

**lecs -** a FORE program that implements the assignment of individual LECs to different emulated LANs.

**leq -** a FORE program that provides information about an ELAN. This information is obtained from the LES, and includes MAC addresses registered on the ELAN together with their corresponding ATM addresses.

**Line Build Out (LBO)** - Because T1 circuits require the last span to lose 15-22.5 dB, a selectable output attenuation is generally required of DTE equipment (typical selections include 0.0, 7.5 and 15 dB of loss at 772 KHz).

**Line Code Violations (LCV) -** Error Event. A Line Coding Violation (LCV) is the occurrence of either a Bipolar Violation (BPV) or Excessive Zeroes (EXZ) Error Event.

**Link -** An entity that defines a topological relationship (including available transport capacity) between two nodes in different subnetworks. Multiple links may exist between a pair of subnetworks. Synonymous with logical link.

**Link Access Procedure, Balanced (LAPB) -** Data link protocol in the X.25 protocol stack. LAPB is a bit-oriented protocol derived from HDLC. See also HDLC and X.25.

Link Down Trap - an SNMP trap, sent when an interface changes from a normal state to an error state, or is disconnected.

Link Layer - layer in the OSI model regarding transmission of data between network nodes.

**Link Up Trap -** an SNMP trap, sent when an interface changes from an error condition to a normal state.

**Load Sharing -** Two or more computers in a system that share the load during peak hours. During periods of non peak hours, one computer can manage the entire load with the other acting as a backup.

**Local Access and Transport Area (LATA) -** Geographic boundaries of the local telephone network, specified by the FCC, in which a single LEC may perform its operations. Communications outside or between LATAs are provided by IXCs.

**Local Area Network (LAN) -** a data network intended to serve an area of only a few square kilometers or less. Because the network is known to cover only a small area, optimizations can be made in the network signal protocols that permit higher data rates.

**Logical Link Control (LLC) -** protocol developed by the IEEE 802 committee for data-link-layer transmission control; the upper sublayer of the IEEE Layer 2 (OSI) protocol that complements the MAC protocol; IEEE standard 802.2; includes end-system addressing and error checking.

**Loopback** - a troubleshooting technique that returns a transmitted signal to its source so that the signal can be analyzed for errors. Typically, a loopback is set at various points in a line until the section of the line that is causing the problem is discovered.

**looptest -** program that tests an interface for basic cell reception and transmission functionality, usually used for diagnostic purposes to determine if an interface is functioning properly.

Loss Of Frame (LOF) - a type of transmission error that may occur in wide-area carrier lines.

**Loss Of Pointer (LOP)** - a type of transmission error that may occur in wide-area carrier lines.

**Loss Of Signal (LOS)** - a type of transmission error that may occur in wide-area carrier lines, or a condition declared when the DTE senses a loss of a DS1 signal from the CPE for more the 150 milliseconds (the DTE generally responds with an all ones "Blue or AIS" signal).

**Management Information Base (MIB) -** the set of parameters that an SNMP management station can query or set in the SNMP agent of a networked device (e.g., router).

**Maximum Burst Size (MBS)** - the Burst Tolerance (BT) is conveyed through the MBS which is coded as a number of cells. The BT together with the SCR and the GCRA determine the MBS that may be transmitted at the peak rate and still be in conformance with the GCRA.

**Maximum Burst Tolerance -** the largest burst of data that a network device is guaranteed to handle without discarding cells or packets. Bursts of data larger than the maximum burst size may be subject to discard.

**Maximum Cell Delay Variance (MCDV) -** This is the maximum two-point CDV objective across a link or node for the specified service category.

**Maximum Cell Loss Ratio (MCLR) -** This is the maximum ratio of the number of cells that do not make it across the link or node to the total number of cells arriving at the link or node.

**Maximum Cell Transfer Delay (MCTD) -** This is the sum of the fixed delay component across the link or node and MCDV.

**Maximum Transmission Unit (MTU) -** the largest unit of data that can be sent over a type of physical medium.

Mbps - megabits per second (million)

**Media Access Control (MAC) -** a media-specific access control protocol within IEEE 802 specifications; currently includes variations for Token Ring, token bus, and CSMA/CD; the lower sublayer of the IEEE's link layer (OSI), which complements the Logical Link Control (LLC).

**Media Attachment Unit (MAU)** - device used in Ethernet and IEEE 802.3 networks that provides the interface between the AUI port of a station and the common medium of the Ethernet. The MAU, which can be built into a station or can be a separate device, performs physical layer functions including conversion of the digital data from the Ethernet interface, collision detection, and injection of bits onto the network.

Media Interface Connector (MIC) - fiber optic connector that joins fiber to the FDDI controller.

**Message Identifier (MID) -** message identifier used to associate ATM cells that carry segments from the same higher layer packet.

**Metasignalling -** an ATM Layer Management (LM) process that manages different types of signalling and possibly semipermanent virtual channels (VCs), including the assignment, removal, and checking of VCs.

**Metasignalling VCs -** the standardized VCs that convey metasignalling information across a User-to-Network Interface (UNI).

**Metropolitan Area Network (MAN) -** network designed to carry data over an area larger than a campus such as an entire city and its outlying area.

**MicroChannel -** a proprietary 16- or 32-bit bus developed by IBM for its PS/2 computers' internal expansion cards; also offered by others.

**Minimum Cell Rate (MCR) -** parameter defined by the ATM Forum for ATM traffic management, defined only for ABR transmissions and specifying the minimum value for the ACR.

Most Significant Bit (MSB) - highest order bit in the binary representation of a numerical value.

**Motion Picture Experts Group (MPEG) -** ISO group dealing with video and audio compression techniques and mechanisms for multiplexing and synchronizing various media streams.

**MPOA Client -** A device which implements the client side of one or more of the MPOA protocols, (i.e., is a SCP client and/or an RDP client. An MPOA Client is either an Edge Device Functional Group (EDFG) or a Host Behavior Functional Group (HBFG).

MPOA Server - An MPOA Server is any one of an ICFG or RSFG.

**MPOA Service Area -** The collection of server functions and their clients. A collection of physical devices consisting of an MPOA server plus the set of clients served by that server.

**MPOA Target -** A set of protocol address, path attributes, (e.g., internetwork layer QoS, other information derivable from received packet) describing the intended destination and its path attributes that MPOA devices may use as lookup keys.

**Mu-Law -** The PCM coding and companding standard used in Japan and North America.

**Multicasting -** The ability to broadcast messages to one node or a select group of nodes.

Multi-homed - a device having both an ATM and another network connection, like Ethernet.

**Multimode Fiber Optic Cable (MMF)** - fiber optic cable in which the signal or light propagates in multiple modes or paths. Since these paths may have varying lengths, a transmitted pulse of light may be received at different times and smeared to the point that pulses may interfere with surrounding pulses. This may cause the signal to be difficult or impossible to receive. This pulse dispersion sometimes limits the distance over which a MMF link can operate.

**Multiplexing -** a function within a layer that interleaves the information from multiple connections into one connection (see demultiplexing).

**Multipoint Access -** user access in which more than one terminal equipment (TE) is supported by a single network termination.

**Multipoint-to-Multipoint Connection -** a collection of associated ATM VC or VP links, and their associated endpoint nodes, with the following properties:

- 1. All N nodes in the connection, called Endpoints, serve as a Root Node in a Point-to-Multipoint connection to all of the (N-1) remaining endpoints.
- 2. Each of the endpoints can send information directly to any other endpoint, but the receiving endpoint cannot distinguish which of the endpoints is sending information without additional (e.g., higher layer) information.

**Multipoint-to-Point Connection -** a Point-to-Multipoint Connection may have zero bandwidth from the Root Node to the Leaf Nodes, and non-zero return bandwidth from the Leaf Nodes to the Root Node. Such a connection is also known as a Multipoint-to-Point Connection.

**Multiprotocol over ATM (MPOA) -** An effort taking place in the ATM Forum to standardize protocols for the purpose of running multiple network layer protocols over ATM.

Narrowband Channel - sub-voicegrade channel with a speed range of 100 to 200 bps.

National TV Standards Committee (NTSC) - Started in the US in 1953 from a specification laid down by the National Television Standards Committee. It takes the B-Y and R-Y color difference signals, attenuates them to I and Q, then modulates them using double-sideband suppressed subcarrier at 3.58MHz. The carrier reference is sent to the receiver as a burst during the back porch. An industry group that defines how television signals are encoded and transmitted in the US. (See also PAL, SECAM for non-U.S. countries).

**Near-End -** in a relationship between two devices in a circuit, the near-end device is the one that is local.

**Network Layer** - Layer three In the OSI model, the layer that is responsible for routing data across the network.

**Network Management Entity (NM) -** body of software in a switching system that provides the ability to manage the PNNI protocol. NM interacts with the PNNI protocol through the MIB.

**Network Management Layer (NML) -** an abstraction of the functions provided by systems which manage network elements on a collective basis, providing end-to-end network monitoring.

**Network Management Station (NMS) -** system responsible for managing a network or portion of a network by talking to network management agents, which reside in the managed nodes.

**Network Module -** ATM port interface cards which may be individually added to or removed from any *ForeRunner* ATM switch to provide a diverse choice of connection alternatives.

**Network Parameter Control (NPC)** - Defined as the set of actions taken by the network to monitor and control traffic from the NNI. Its main purpose is to protect network resources from malicious as well as unintentional misbehavior which can affect the QoS of other already established connections by detecting violations of negotiated parameters and taking appropriate actions. Refer to UPC.

**Network Redundancy -** Duplicated network equipment and/or data which can provide a backup in case of network failures.

**Network Service Access Point (NSAP) -** OSI generic standard for a network address consisting of 20 octets. ATM has specified E.164 for public network addressing and the NSAP address structure for private network addresses.

**Network-to-Network Interface or Network Node Interface (NNI) -** the interface between two public network pieces of equipment.

**Node -** A computer or other device when considered as part of a network.

Non Return to Zero (NRZ) - a binary encoding scheme in which ones and zeroes are represented by opposite and alternating high and low voltages and where there is no return to a zero (reference) voltage between encoded bits.

**Non Return to Zero Inverted (NRZI) -** A binary encoding scheme that inverts the signal on a "1" and leaves the signal unchanged for a "0". (Also called transition encoding.)

Nonvolatile Storage - Memory storage that does not lose its contents when power is turned off.

**NuBus -** a high-speed bus used in Macintosh computers, structured so users can put a card into any slot on the board without creating conflict over the priority between those cards.

nx64K - This refers to a circuit bandwidth or speed provided by the aggregation of nx64 kbps channels (where n= integer > 1). The 64K or DS0 channel is the basic rate provided by the T Carrier systems.

**Nyquist Theorem -** In communications theory, a formula stating that two samples per cycle is sufficient to characterize a bandwidth limited analog signal; in other words, the sampling rate must be twice the highest frequency component of the signal (i.e., sample 4 KHz analog voice channels 8000 times per second).

**Object Identifier (OID) -** the address of a MIB variable.

Octet - a grouping of 8 bits; similar, but not identical to, a byte.

One's Density - The requirement for digital transmission lines in the public switched telephone network that eight consecutive "0"s cannot be in a digital data stream; exists because repeaters and clocking devices within the network will lose timing after receiving eight "0"s in a row; a number of techniques are used to insert a "1" after every seventh-consecutive "0" (see Bit Stuffing).

**Open Shortest Path First (OSPF) Protocol -** a routing algorithm for IP that incorporates least-cost, equal-cost, and load balancing.

**Open Systems Interconnection (OSI) -** the 7-layer suite of protocols designed by ISO committees to be the international standard computer network architecture.

**OpenView -** Hewlett-Packard's network management software.

**Operation and Maintenance (OAM) Cell -** a cell that contains ATM LM information. It does not form part of the upper layer information transfer.

**Optical Carrier level-n (OC-n) -** The optical counterpart of STS-n (the basic rate of 51.84 Mbps on which SONET is based is referred to as OC-1 or STS-1).

Organizationally Unique Identifier (OUI) - Part of RFC 1483. A three-octet field in the SubNetwork Attachment Point (SNAP) header, identifying an organization which administers the meaning of the following two octet Protocol Identifier (PID) field in the SNAP header. Together they identify a distinct routed or bridged protocol.

**Out-of-Band Management -** refers to switch configuration via the serial port or over Ethernet, not ATM.

Out-of-Frame (OOF) - a signal condition and alarm in which some or all framing bits are lost.

**Packet -** An arbitrary collection of data grouped and transmitted with its user identification over a shared facility.

**Packet Assembler Disassembler (PAD) -** interface device that buffers data sent to/from character mode devices, and assembles and disassembles the packets needed for X.25 operation.

**Packet Internet Groper (ping) -** a program used to test reachability of destinations by sending them an ICMP echo request and waiting for a reply.

**Packet Level Protocol (PLP)** - Network layer protocol in the X.25 protocol stack. Sometimes called X.25 Level 3 or X.25 Protocol.

**Packet Switched Network (PSN) -** a network designed to carry data in the form of packets. The packet and its format is internal to that network.

**Packet Switching -** a communications paradigm in which packets (messages) are individually routed between hosts with no previously established communications path.

**Payload Scrambling -** a technique that eliminates certain bit patterns that may occur within an ATM cell payload that could be misinterpreted by certain sensitive transmission equipment as an alarm condition.

**Payload Type (PT)** - bits 2...4 in the fourth byte of an ATM cell header. The PT indicates the type of information carried by the cell. At this time, values 0...3 are used to identify various types of user data, values 4 and 5 indicate management information, and values 6 and 7 are reserved for future use.

**Peak Cell Rate -** at the PHY Layer SAP of a point-to-point VCC, the Peak Cell Rate is the inverse of the minimum inter-arrival time T0 of the request to send an ATM-SDU.

**Peak Cell Rate (PCR) -** parameter defined by the ATM Forum for ATM traffic management. In CBR transmissions, PCR determines how often data samples are sent. In ABR transmissions, PCR determines the maximum value of the ACR.

**Peer Entities -** entities within the same layer.

Peripheral Component Interconnect (PCI) - a local-bus standard created by Intel.

**Permanent Virtual Channel Connection (PVCC)** - A Virtual Channel Connection (VCC) is an ATM connection where switching is performed on the VPI/VCI fields of each cell. A Permanent VCC is one which is provisioned through some network management function and left up indefinitely.

**Permanent Virtual Circuit (or Channel) (PVC) -** a circuit or channel through an ATM network provisioned by a carrier between two endpoints; used for dedicated long-term information transport between locations.

**Permanent Virtual Path Connection (PVPC) -** A Virtual Path Connection (VPC) is an ATM connection where switching is performed on the VPI field only of each cell. A PVPC is one which is provisioned through some network management function and left up indefinitely.

Phase Alternate Line (PAL) - Largely a German/British development in the late 60s, used in the UK and much of Europe. The B-Y and R-Y signals are weighted to U and V, then modulated onto a double-sideband suppressed subcarrier at 4.43MHz. The V (R-Y) signal's phase is turned through 180 degrees on each alternate line. This gets rid of NTSC's hue changes with phase errors at the expense of de-saturation. The carrier reference is sent as a burst in the back porch. The phase of the burst is alternated every line to convey the phase switching of the V signal. The burst's average phase is -V. (see NTSC for U.S.).

**Physical Layer (PHY) -** the actual cards, wires, and/or fiber-optic cabling used to connect computers, routers, and switches.

**Physical Layer Connection -** an association established by the PHY between two or more ATM-entities. A PHY connection consists of the concatenation of PHY links in order to provide an end-to-end transfer capability to PHY SAPs.

**Physical Layer Convergence Protocol (PLCP) -** a framing protocol that runs on top of the T1 or E1 framing protocol.

**Physical Medium (PM) -** Refers to the actual physical interfaces. Several interfaces are defined including STS-1, STS-3c, STS-12c, STM-1, STM-4, DS1, E1, DS2, E3, DS3, E4, FDDI-based, Fiber Channel-based, and STP. These range in speeds from 1.544Mbps through 622.08 Mbps.

**Physical Medium Dependent (PMD) -** a sublayer concerned with the bit transfer between two network nodes. It deals with wave shapes, timing recovery, line coding, and electro-optic conversions for fiber based links.

**Plesiochronous -** two signals are plesiochronous if their corresponding significant instants occur at nominally the same rate, with variations in rate constrained to specified limits.

**Point of Demarcation -** the dividing line between a carrier and the customer premise that is governed by strict standards that define the characteristics of the equipment on each side of the demarcation. Equipment on one side of the point of demarcation is the responsibility of the customer. Equipment on the other side of the point of demarcation is the responsibility of the carrier.

**Point-to-Multipoint Connection -** a collection of associated ATM VC or VP links, with associated endpoint nodes, with the following properties:

- 1. One ATM link, called the Root Link, serves as the root in a simple tree topology. When the Root node sends information, all of the remaining nodes on the connection, called Leaf nodes, receive copies of the information.
- 2. Each of the Leaf Nodes on the connection can send information directly to the Root Node. The Root Node cannot distinguish which Leaf is sending information without additional (higher layer) information. (See the following note for Phase 1.)
- 3. The Leaf Nodes cannot communicate directly to each other with this connection type.

Note: Phase 1 signalling does not support traffic sent from a Leaf to the Root.

**Point-to-Point Connection -** a connection with only two endpoints.

**Point-to-Point Protocol (PPP) -** Provides a method for transmitting packets over serial point-to-point links.

**Policing -** the function that ensures that a network device does not accept traffic that exceeds the configured bandwidth of a connection.

**Port Identifier -** The identifier assigned by a logical node to represent the point of attachment of a link to that node.

Presentation Layer - Sixth layer of the OSI model, providing services to the application layer.

**Primary Reference Source (PRS) -** Equipment that provides a timing signal whose long-term accuracy is maintained at  $1 \times 10$  -11 or better with verification to universal coordinated time (UTC) and whose timing signal is used as the basis of reference for the control of other clocks within a network.

**Primitive -** an abstract, implementation-independent interaction between a layer service user and a layer service provider.

**Priority -** the parameter of ATM connections that determines the order in which they are reduced from the peak cell rate to the sustained cell rate in times of congestion. Connections with lower priority (4 is low, 1 is high) are reduced first.

**Private Branch Exchange (PBX) -** a private phone system (switch) that connects to the public telephone network and offers in-house connectivity. To reach an outside line, the user must dial a digit like 8 or 9.

**Private Network Node Interface or Private Network-to-Network Interface (PNNI) -** a protocol that defines the interaction of private ATM switches or groups of private ATM switches

**Programmable Read-Only Memory (PROM) -** a chip-based information storage area that can be recorded by an operator but erased only through a physical process.

**Protocol -** a set of rules and formats (semantic and syntactic) that determines the communication behavior of layer entities in the performance of the layer functions.

**Protocol Control Information -** the information exchanged between corresponding entities using a lower layer connection to coordinate their joint operation.

**Protocol Data Unit (PDU) -** a unit of data specified in a layer protocol and consisting of protocol control information and layer user data.

**Proxy** - the process in which one system acts for another system to answer protocol requests.

**Proxy Agent -** an agent that queries on behalf of the manager, used to monitor objects that are not directly manageable.

**Public Data Network (PDN) -** a network designed primarily for data transmission and intended for sharing by many users from many organizations.

**Pulse Code Modulation (PCM) -** a modulation scheme that samples the information signals and transmits a series of coded pulses to represent the data.

**Q.2931 -** Derived from Q.93B, the narrowband ISDN signalling protocol, an ITU standard describing the signalling protocol to be used by switched virtual circuits on ATM LANs.

**Quality of Service (QoS) -** Quality of Service is defined on an end-to-end basis in terms of the following attributes of the end-to-end ATM connection:

Cell Loss Ratio

Cell Transfer Delay

Cell Delay Variation

**Queuing Delay (QD)** - refers to the delay imposed on a cell by its having to be buffered because of unavailability of resources to pass the cell onto the next network function or element. This buffering could be a result of oversubscription of a physical link, or due to a connection of higher priority or tighter service constraints getting the resource of the physical link.

Radio Frequency Interference (RFI) - the unintentional transmission of radio signals. Computer equipment and wiring can both generate and receive RFI.

**Real-Time Clock -** a clock that maintains the time of day, in contrast to a clock that is used to time the electrical pulses on a circuit.

**Red Alarm -** In T1, a red alarm is generated for a locally detected failure such as when a condition like OOF exists for 2.5 seconds, causing a CGA, (Carrier Group Alarm).

**Reduced Instruction Set Computer (RISC) -** a generic name for CPUs that use a simpler instruction set than more traditional designs.

**Redundancy -** In a data transmission, the fragments of characters and bits that can be eliminated with no loss of information.

**Registration -** The address registration function is the mechanism by which Clients provide address information to the LAN Emulation Server.

**Relaying -** a function of a layer by means of which a layer entity receives data from a corresponding entity and transmits it to another corresponding entity.

**Request To Send (RTS) -** an RS-232 modem interface signal (sent from the DTE to the modem on pin 4) which indicates that the DTE has data to transmit.

**Requests For Comment (RFCs) -** IETF documents suggesting protocols and policies of the Internet, inviting comments as to the quality and validity of those policies. These comments are collected and analyzed by the IETF in order to finalize Internet standards.

**RFC1483** - Multiprotocol Encapsulation over ATM Adaptation Layer 5.

RFC1490 - Multiprotocol Interconnect over Frame Relay.

**RFC1577 -** Classical IP and ARP over ATM.

**RFC1755 -** ATM Signaling Support for IP over ATM.

**Robbed-Bit Signaling -** In T1, refers to the use of the least significant bit of every word of frames 6 and 12 (D4), or 6, 12, 18, and 24 (ESF) for signaling purposes.

**Route Server -** A physical device that runs one or more network layer routing protocols, and which uses a route query protocol in order to provide network layer routing forwarding descriptions to clients.

**Router -** a device that forwards traffic between networks or subnetworks based on network layer information.

**Routing Domain (RD) -** A group of topologically contiguous systems which are running one instance of routing.

**Routing Information Protocol (RIP) -** a distance vector-based protocol that provides a measure of distance, or hops, from a transmitting workstation to a receiving workstation.

**Routing Protocol** - A general term indicating a protocol run between routers and/or route servers in order to exchange information used to allow computation of routes. The result of the routing computation will be one or more forwarding descriptions.

**SBus -** hardware interface for add-in boards in later-version Sun 3 workstations.

**Scalable Processor Architecture Reduced instruction set Computer (SPARC) -** a powerful workstation similar to a reduced-instruction-set-computing (RISC) workstation.

**Segment -** a single ATM link or group of interconnected ATM links of an ATM connection.

Segmentation And Reassembly (SAR) - the SAR accepts PDUs from the CS and divides them into very small segments (44 bytes long). If the CS-PDU is less than 44 bytes, it is padded to 44 with zeroes. A two-byte header and trailer are added to this basic segment. The header identifies the message type (beginning, end, continuation, or single) and contains sequence numbering and message identification. The trailer gives the SAR-PDU payload length, exclusive of pad, and contains a CRC check to ensure the SAR-PDU integrity. The result is a 48-byte PDU that fits into the payload field of an ATM cell.

**Selector (SEL) -** A subfield carried in SETUP message part of ATM endpoint address Domain specific Part (DSP) defined by ISO 10589, not used for ATM network routing, used by ATM end systems only.

**Semipermanent Connection -** a connection established via a service order or via network management.

**Serial Line IP (SLIP) -** A protocol used to run IP over serial lines, such as telephone circuits or RS-232 cables, interconnecting two systems.

**Service Access Point (SAP) -** the point at which an entity of a layer provides services to its LM entity or to an entity of the next higher layer.

**Service Data Unit (SDU) -** a unit of interface information whose identity is preserved from one end of a layer connection to the other.

**Service Specific Connection Oriented Protocol (SSCOP) -** an adaptation layer protocol defined in ITU-T Specification: Q.2110.

**Service Specific Convergence Sublayer (SSCS) -** The portion of the convergence sublayer that is dependent upon the type of traffic that is being converted.

**Session Layer -** Layer 5 in the OSI model that is responsible for establishing and managing sessions between the application programs running in different nodes.

**Severely Errored Seconds (SES) -** a second during which more event errors have occurred than the SES threshold (normally 10-3).

**Shaping Descriptor -** *n* ordered pairs of GCRA parameters (I,L) used to define the negotiated traffic shape of an APP connection. The traffic shape refers to the load-balancing of a network, where load-balancing means configuring data flows to maximize network efficiency.

**Shielded Pair -** Two insulated wires in a cable wrapped with metallic braid or foil to prevent interference and provide noise free transmission.

**Shielded Twisted Pair (STP) -** two or more insulated wires, twisted together and then wrapped in a cable with metallic braid or foil to prevent interference and offer noise-free transmissions.

**Signaling System No. 7 (SS7) -** The SS7 protocol has been specified by ITU-T and is a protocol for interexchange signaling.

Simple and Efficient Adaptation Layer (SEAL) - also called AAL 5, this ATM adaptation layer assumes that higher layer processes will provide error recovery, thereby simplifying the SAR portion of the adaptation layer. Using this AAL type packs all 48 bytes of an ATM cell information field with data. It also assumes that only one message is crossing the UNI at a time. That is, multiple end-users at one location cannot interleave messages on the same VC, but must queue them for sequential transmission.

Simple Gateway Management Protocol (SGMP) - the predecessor to SNMP.

Simple Mail Transfer Protocol (SMTP) - the Internet electronic mail protocol used to transfer electronic mail between hosts.

Simple Network Management Protocol (SNMP) - the Internet standard protocol for managing nodes on an IP network.

**Simple Protocol for ATM Network Signalling (SPANS) -** FORE Systems' proprietary signalling protocol used for establishing SVCs between FORE Systems equipment.

**Single Mode Fiber (SMF)** - Fiber optic cable in which the signal or light propagates in a single mode or path. Since all light follows the same path or travels the same distance, a transmitted pulse is not dispersed and does not interfere with adjacent pulses. SMF fibers can support longer distances and are limited mainly by the amount of attenuation. Refer to MMF.

**Small Computer Systems Interface (SCSI) -** a standard for a controller bus that connects hardware devices to their controllers on a computer bus, typically used in small systems.

**Smart PVC (SPVC)** - a generic term for any communications medium which is permanently provisioned at the end points, but switched in the middle. In ATM, there are two kinds of SPVCs: smart permanent virtual path connections (SPVPCs) and smart permanent virtual channel connections (SPVCCs).

**snmpd** - an SMNP agent for a given adapter card.

**Source -** Part of communications system which transmits information.

**Source Address (SA) -** The address from which the message or data originated.

**Source MAC Address (SA) -** A six octet value uniquely identifying an end point and which is sent in an IEEE LAN frame header to indicate source of frame.

**Source Traffic Descriptor -** a set of traffic parameters belonging to the ATM Traffic Descriptor used during the connection set-up to capture the intrinsic traffic characteristics of the connection requested by the source.

**Spanning Tree Protocol -** provides loop-free topology in a network environment where there are redundant paths.

**Static Route -** a route that is entered manually into the routing table.

**Statistical Multiplexing -** a technique for allowing multiple channels and paths to share the same link, typified by the ability to give the bandwidth of a temporarily idle channel to another channel.

**Stick and Click (SC) -** Designation for an Optical Connector featuring a 2.5 mm physically contacting ferrule with a push-pull mating design. Commonly referred to as Structured Cabling, Structured Connectors or Stick and Click

**Stick and Turn (ST) -** A fiber-optic connector designed by AT&T which uses the bayonet style coupling rather than screw-on as the SMA uses. The ST is generally considered the eventual replacement for the SMA type connector.

**Store-and-Forward -** the technique of receiving a message, storing it until the proper outgoing line is available, then retransmitting it, with no direct connection between incoming and outgoing lines.

Straight Tip (ST) - see Stick and Turn.

Structured Cabling (SC) - see Stick and Click.

Structured Connectors (SC) - see Stick and Click.

**Sublayer -** a logical subdivision of a layer.

**SubNetwork Access Protocol (SNAP) -** a specially reserved variant of IEEE 802.2 encoding SNAP indicates to look further into the packet where it will fine a Type field.

**Subscriber Network Interface (SNI) -** the interface between an SMDS end user's CPE and the network directly serving the end user, supported by either a DS1 or DS3 access arrangement.

**Super Frame (SF) -** a term used to describe the repeating 12 D4 frame format that composes a standard (non-ESF) T1 service.

**Super User -** a login ID that allows unlimited access to the full range of a device's functionality, including especially the ability to reconfigure the device and set passwords.

**Sustainable Cell Rate (SCR) -** ATM Forum parameter defined for traffic management. For VBR connections, SCR determines the long-term average cell rate that can be transmitted.

**Sustained Information Rate (SIR) -** In ATM this refers to the long-term average data transmission rate across the User-to-Network Interface. In SMDS this refers to the committed information rate (similar to CIR for Frame Relay Service).

**Switch -** Equipment used to interconnect lines and trunks.

**Switched Connection -** A connection established via signaling.

**Switched Multimegabit Data Service (SMDS) -** a high-speed, datagram-based, public data network service expected to be widely used by telephone companies in their data networks.

**Switched Virtual Channel Connection (SVCC)** - A Switched VCC is one which is established and taken down dynamically through control signaling. A Virtual Channel Connection (VCC) is an ATM connection where switching is performed on the VPI/VCI fields of each cell.

**Switched Virtual Circuit (or Channel) (SVC) -** a channel established on demand by network signalling, used for information transport between two locations and lasting only for the duration of the transfer; the datacom equivalent of a dialed telephone call.

**Switched Virtual Path Connection (SVPC)** - a connection which is established and taken down dynamically through control signaling. A Virtual Path Connection (VPC) is an ATM connection where switching is performed on the VPI field only of each cell.

**Switching System -** A set of one or more systems that act together and appear as a single switch for the purposes of PNNI routing.

 $\textbf{Symmetric Connection -} a \ connection \ with \ the \ same \ bandwidth \ specified \ for \ both \ directions.$ 

**Synchronous -** signals that are sourced from the same timing reference and hence are identical in frequency.

Synchronous Data Link Control (SDLC) - IBM's data link protocol used in SNA networks.

**Synchronous Optical Network (SONET) -** a body of standards that defines all aspects of transporting and managing digital traffic over optical facilities in the public network.

**Synchronous Payload Envelope (SPE) -** the payload field plus a little overhead of a basic SONET signal.

**Synchronous Transfer Mode (STM) -** a transport and switching method that depends on information occurring in regular, fixed patterns with respect to a reference such as a frame pattern.

**Synchronous Transport Signal (STS) -** a SONET electrical signal rate.

**Systeme En Coleur Avec Memoire (SECAM) - Sequential and Memory Color Television -** Started in France in the late 60s, and used by other countries with a political affiliation. This is. The B-Y and R-Y signals are transmitted on alternate lines modulated on an FM subcarrier. The memory is a one line delay line in the receiver to make both color difference signals available at the same time on all lines. Due to FM, the signal is robust in difficult terrain.

**Systems Network Architecture (SNA) -** a proprietary networking architecture used by IBM and IBM-compatible mainframe computers.

**T1 -** a specification for a transmission line. The specification details the input and output characteristics and the bandwidth. T1 lines run at 1.544 Mbps and provide for 24 data channels. In common usage, the term "T1" is used interchangeably with "DS1."

**T1 Link -** A wideband digital carrier facility used for transmission of digitized voice, digital data, and digitized image traffic. This link is composed of two twisted-wire pairs that can carry 24 digital channels, each operating at 64K bps at the aggregate rate of 1.544M bps, full duplex. Also referred to as DS-1.

**T3** - a specification for a transmission line, the equivalent of 28 T1 lines. T3 lines run at 44.736 Mbps. In common usage, the term "T3" is used interchangeably with "DS3."

**Tachometer -** in *ForeView*, the tachometer shows the level of activity on a given port. The number in the tachometer shows the value of a chosen parameter in percentage, with a colored bar providing a semi-logarithmic representation of that percentage.

**Tagged Cell Rate (TCR) -** An ABR service parameter, TCR limits the rate at which a source may send out-of-rate forward RM-cells. TCR is a constant fixed at 10 cells/second.

**Telephony -** The conversion of voices and other sounds into electrical signals which are then transmitted by telecommunications media.

**Telnet** - a TCP/IP protocol that defines a client/server mechanism for emulating directly-connected terminal connections.

**Terminal Equipment (TE) -** Terminal equipment represents the endpoint of ATM connection(s) and termination of the various protocols within the connection(s).

**Throughput -** Measurement of the total useful information processed or communicated by a computer during a specified time period, i.e. packets per second.

**Time Division Multiplexing (TDM) -** a method of traditional digital multiplexing in which a signal occupies a fixed, repetitive time slot within a higher-rate signal.

**Token Ring -** a network access method in which the stations circulate a token. Stations with data to send must have the token to transmit their data.

**topology** - a program that displays the topology of a FORE Systems ATM network. An updated topology can be periodically re-displayed by use of the interval command option.

**Traffic -** the calls being sent and received over a communications network. Also, the packets that are sent on a data network.

**Traffic Management (TM)** - The traffic control and congestion control procedures for ATM. ATM layer traffic control refers to the set of actions taken by the network to avoid congestion conditions. ATM layer congestion control refers to the set of actions taken by the network to minimize the intensity, spread and duration of congestion. The following functions form a framework for managing and controlling traffic and congestion in ATM networks and may be used in appropriate combinations:

Connection Admission Control Feedback Control Usage Parameter Control Priority Control Traffic Shaping Network Resource Management Frame Discard ABR Flow Control

**Traffic Parameter** - A parameter for specifying a particular traffic aspect of a connection.

**Trailer -** the protocol control information located at the end of a PDU.

**Transit Delay -** the time difference between the instant at which the first bit of a PDU crosses one designated boundary, and the instant at which the last bit of the same PDU crosses a second designated boundary.

**Transmission Control Protocol (TCP)** - a specification for software that bundles and unbundles sent and received data into packets, manages the transmission of packets on a network, and checks for errors.

Transmission Control Protocol/Internet Protocol (TCP/IP) - a set of communications protocols that has evolved since the late 1970s, when it was first developed by the Department of Defense. Because programs supporting these protocols are available on so many different computer systems, they have become an excellent way to connect different types of computers over networks.

**Transmission Convergence (TC) -** generates and receives transmission frames and is responsible for all overhead associated with the transmission frame. The TC sublayer packages cells into the transmission frame.

**Transmission Convergence Sublayer (TCS) -** This is part of the ATM physical layer that defines how cells will be transmitted by the actual physical layer.

**Transparent Asynchronous Transmitter/Receiver Interface (TAXI) -** Encoding scheme used for FDDI LANs as well as for ATM; supports speed typical of 100 Mbps over multimode fiber.

**Transport Layer** - Layer Four of the OSI reference model that is responsible for maintaining reliable end-to-end communications across the network.

**trap** - a program interrupt mechanism that automatically updates the state of the network to remote network management hosts. The SNMP agent on the switch supports these SNMP traps.

**Trivial File Transfer Protocol (TFTP) -** Part of IP, a simplified version of FTP that allows files to be transferred from one computer to another over a network.

**Twisted Pair -** Insulated wire in which pairs are twisted together. Commonly used for telephone connections, and LANs because it is inexpensive.

**Unassigned Cells -** a generated cell identified by a standardized virtual path identifier (VPI) and virtual channel identifier (VCI) value, which does not carry information from an application using the ATM Layer service.

**Unavailable Seconds (UAS) -** a measurement of signal quality. Unavailable seconds start accruing when ten consecutive severely errored seconds occur.

**UNI 3.0/3.1 -** the User-to-Network Interface standard set forth by the ATM Forum that defines how private customer premise equipment interacts with private ATM switches.

**Unicasting -** The transmit operation of a single PDU by a source interface where the PDU reaches a single destination.

**Universal Test & Operations Interface for ATM (UTOPIA) -** Refers to an electrical interface between the TC and PMD sublayers of the PHY layer.

**Unshielded Twisted Pair (UTP) -** a cable that consists of two or more insulated conductors in which each pair of conductors are twisted around each other. There is no external protection and noise resistance comes solely from the twists.

**Unspecified Bit Rate (UBR) -** a type of traffic that is not considered time-critical (e.g., ARP messages, pure data), allocated whatever bandwidth is available at any given time. UBR traffic is given a "best effort" priority in an ATM network with no guarantee of successful transmission.

**Uplink -** Represents the connectivity from a border node to an upnode.

**Usage Parameter Control (UPC) -** mechanism that ensures that traffic on a given connection does not exceed the contracted bandwidth of the connection, responsible for policing or enforcement. UPC is sometimes confused with congestion management (see *congestion management*).

**User Datagram Protocol (UDP) -** the TCP/IP transaction protocol used for applications such as remote network management and name-service access; this lets users assign a name, such as "RVAX\*2,S," to a physical or numbered address.

**User-to-Network Interface (UNI) -** the physical and electrical demarcation point between the user and the public network service provider.

**V.35 -** ITU-T standard describing a synchronous, physical layer protocol used for communications between a network access device and a packet network. V.35 is most commonly used in the United States and Europe, and is recommended for speeds up to 48 Kbps.

**Variable Bit Rate (VBR) -** a type of traffic that, when sent over a network, is tolerant of delays and changes in the amount of bandwidth it is allocated (e.g., data applications).

**Virtual Channel (or Circuit) (VC) -** a communications path between two nodes identified by label rather than fixed physical path.

**Virtual Channel Connection (VCC)** - a unidirectional concatenation of VCLs that extends between the points where the ATM service users access the ATM Layer. The points at which the ATM cell payload is passed to, or received from, the users of the ATM Layer (i.e., a higher layer or ATMM-entity) for processing signify the endpoints of a VCC.

**Virtual Channel Identifier (VCI)** - the address or label of a VC; a value stored in a field in the ATM cell header that identifies an individual virtual channel to which the cell belongs. VCI values may be different for each data link hop of an ATM virtual connection.

**Virtual Channel Link (VCL)** - a means of unidirectional transport of ATM cells between the point where a VCI value is assigned and the point where that value is translated or removed.

**Virtual Channel Switch -** a network element that connects VCLs. It terminates VPCs and translates VCI values. The Virtual Channel Switch is directed by Control Plane functions and relays the cells of a VC.

**Virtual Connection -** an endpoint-to-endpoint connection in an ATM network. A virtual connection can be either a virtual path or a virtual channel.

**Virtual Local Area Network (VLAN) -** Work stations connected to an intelligent device which provides the capabilities to define LAN membership.

**Virtual Network Software (VINES) -** Banyan's network operating system based on UNIX and its protocols.

Virtual Path (VP) - a unidirectional logical association or bundle of VCs.

**Virtual Path Connection (VPC) -** a concatenation of VPLs between virtual path terminators (VPTs). VPCs are unidirectional.

**Virtual Path Identifier (VPI)** - the address or label of a particular VP; a value stored in a field in the ATM cell header that identifies an individual virtual path to which the cell belongs. A virtual path may comprise multiple virtual channels.

**Virtual Path Link (VPL) -** a means of unidirectional transport of ATM cells between the point where a VPI value is assigned and the point where that value is translated or removed.

**Virtual Path Switch -** a network element that connects VPLs, it translates VPI (not VCI) values and is directed by Control Plane functions. The Virtual Path Switch relays the cells of a Virtual Path.

**Virtual Path Terminator (VPT) -** a system that unbundles the VCs of a VP for independent processing of each VC.

**Virtual Private Data Network (VPDN) -** a private data communications network built on public switching and transport facilities rather than dedicated leased facilities such as T1s.

**Virtual Private Network (VPN) -** a private voice communications network built on public switching and transport facilities rather than dedicated leased facilities such as T1s.

Virtual Source/Virtual Destination (VS/VD) - An ABR connection may be divided into two or more separately controlled ABR segments. Each ABR control segment, except the first, is sourced by a virtual source. A virtual source implements the behavior of an ABR source endpoint. Backwards RM-cells received by a virtual source are removed from the connection. Each ABR control segment, except the last, is terminated by a virtual destination. A virtual destination assumes the behavior of an ABR destination endpoint. Forward RM-cells received by a virtual destination are turned around and not forwarded to the next segment of the connection.

**Virtual Tributary (VT)** - a structure used to carry payloads such as DS1s that run at significantly lower rates than STS-1s.

Warm Start Trap - an SNMP trap which indicates that SNMP alarm messages or agents have been enabled.

Wide-Area Network (WAN) - a network that covers a large geographic area.

**Wideband Channel -** Communications channel with more capacity (19.2K bps) than the standard capacity of a voice grade line.

**X.21 -** ITU-T standard for serial communications over synchronous digital lines. The X.21 protocol is used primarily in Europe and Japan.

**X.25** - a well-established data switching and transport method that relies on a significant amount of processing to ensure reliable transport over metallic media.

**Yellow Alarm -** An alarm signal sent back toward the source of a failed signal due to the presence of an AIS (may be used by APS equipment to initiate switching).

**Zero Byte Time Slot Interchange (ZBTSI) -** A technique used with the T carrier extended superframe format (ESF) in which an area in the ESF frame carries information about the location of all-zero bytes (eight consecutive "0"s) within the data stream.

**Zero Code Suppression -** The insertion of a "1" bit to prevent the transmission of eight or more consecutive "0" bits. Used primarily with T1 and related digital telephone company facilities, which require a minimum "1's density" in order to keep the individual subchannels of a multiplexed, high speed facility active.

**Zero-Bit Insertion -** A technique used to achieve transparency in bit-oriented protocols. A zero is inserted into sequences of one bits that cause false flag direction.

Glossary

## Index

A	BUS
ABR	С
address registration 2 - 4	CBR
AMI access level	CDVT 1 - 36
all	Cell Delay Variation Tolerance
network	(CDVT) 1 - 36
none	Cell Loss Priority (CLP) bit 1 - 3'
serial 8 - 9	CES (Circuit Emulation Services) B - 1
AMI command privileges 8 - 9	CES connection
admin 8 - 9	creating B - 2
user 8 - 9	displaying
application	Circuit Emulation Services (CES) B - :
configuring on a FramePlus	Classical IP ATM network
network moduleD - 17	<b>configuring 2</b> - 1
area 6 - 6	Classical IP interface 2 - 2
ARP cache 2 - 8	Classical IP over ATM 2 - 1
ARP reply 2 - 8	Classical IP PVC 2 - 10
ARP request	CLP bit
ARP server	configuration
configuring a FORE switch as 2 - 7	emulated LAN, example 4 - 3
ATM ARP (ATM address resolution	connection
protocol) 2 - 6	configuration-direct 3 - 4
ATM Forum LAN Emulation Over ATM	control-direct 3 - 4
Version 1.0	control-distribute
ATM Forum PNNI	data-direct
Available Bit Rate (ABR) 1 - 35	multicast-forward 3 - 4
В	multicast-rol ward 3 - 4
Broadcast and Unknown Server	
(BUS)	connection process, LEC
broadcast packets 4 - 5	Constant Bit Rate (CBR)
BT 1 - 36	containing path
Rurst Tolerance (RT) 1 - 36	crankback 6 - 4

D	F
data encryption 8 - 5	flooding protocol 6 - 2
database exchange protocol 6 - 2	FORE IP
defining an ELAN 3 - 30	ForeThought PNNI 5 - 1
Distributed LAN Emulation	backbone topology 5 - 9
(DLE)3 - 9, 4 - 6	border switches 5 - 8
distributed timing	hello indication messages 5 - 8
ASX-1000 9 - 4	link metrics 5 - 8
ASX-200BX 9 - 4	peer group summary
ASX-200WG 9 - 4	node5 - 8, 5 - 9
ForeRunnerLE 155 9 - 5	physical network 5 - 5
DLE 3 - 9	topology database 5 - 5
DLE peer servers	two-level hierarchy 5 - 6
starting 3 - 40	Frame-based UNI (FUNI) D - 1
domain	FramePlus network module
dynamic path 7 - 8	configuring an applicationD - 17
E	FUNI (Frame-based UNI)
egress rate enforcement	G
configuring D - 10	gateway switch 6 - 5
ELAN	GCRA
components 3 - 2	Generic Cell Rate Algorithm
configuring	(GCRA) 1 - 36
joining	Н
ELAN access control	Hello protocol 6 - 1
disabling 3 - 21	hierarchical routing
enabling 3 - 21	ATMF-PNNI 6 - 3
Emulated LAN (ELAN) 3 - 1	horizontal links
address resolution 3 - 7	HOHZOHtai illiks 0 - 1
data transfer 3 - 8	I
initialization 3 - 6	ILMI (Interim Local Management
operation 3 - 4	Interface) 2 - 4
registration	InARP reply 2 - 8
emulated LAN (ELAN)	InARP request 2 - 8
example configuration 4 - 3	ingress rate enforcement
running multiple 4 - 6	configuring D - 10
ESI (End System Identifier) 2 - 4	initialization process, LEC 4 - 4

Interim Local Management Interface	LES 3 - 3
(ILMI)2 - 4	starting 3 - 40
inverse ARP (InARP)	levels in PNNI 6 - 8
IP filtering 8 - 11	Link Management Interface (LMI)D - 11
authorized IP address table 8 - 11	link-scope UNI
creating an authorized	LIS (Logical IP Subnet) 2 - 2
address 8 - 11	LMI (Link Management Interface)D - 11
deleting an authorized	local authentication 8 - 3
address 8 - 11	Logical IP Subnet (LIS)2 - 2
displaying authorized	logical link (loglink) 5 - 2
addresses	login authentication 8 - 3
information about the last packet dropped 8 - 12	local
statistics	SecurID 8 - 3
IP filtering flags	loglink (logical link) 5 - 2
ir intering riags	M
L	
LAN Emulation Client (LEC) 3 - 3	Maximum Burst Size (MBS) 1 - 35 MBS
LAN emulation client (LEC)	
connection process 4 - 5	MCR
initialization process 4 - 4	Minimum Cell Rate (MCR) 1 - 36
LAN Emulation Configuration Server	multicast packets 4 - 5
(LECS)	N
LAN Emulation Server (LES) 3 - 3	network configuration examples 2 - 11
LAN Emulation services 3 - 2	node 6 - 3
starting 3 - 38	node secret file 8 - 5
leaky bucket algorithm 1 - 36	NSAP addresses 2 - 4
LEC	NSAP prefix 2 - 4
starting 3 - 42	•
LECS	O
well-known address 3 - 42	originating path
LECS configuration file	originating paths
configuring DLE 3 - 30	creating
configuring MPOA 3 - 33	displaying
ELAN access control 3 - 30	displaying advanced information 1 - 10
sample	111101111au011 1 - 10
starting 3 - 38	P
syntax 3 - 24	PCR
LECS control parameters 3 - 32	Peak Cell Rate (PCR) 1 - 35, 1 - 36

peer group leader (PGL) 6 - 3	node secret file 8 - 5
permanent virtual channel	passcode 8 - 4
displaying advanced	personal identification number
information 1 - 16	(PIN) 8 - 3, 8 - 4
permanent virtual circuits (PVCs) 1 - 1	PIN number 8 - 4
PGL	server 8 - 4
peer group leader 6 - 3	tokens 8 - 4
PNNI5 - 1, 6 - 1	SecurID authentication 8 - 3
gateway switch 6 - 5	security 8 - 1
split switch 6 - 5	selector field 2 - 4
PNNI policies 6 - 10	SNMP indexing
PNNI signalling protocol 6 - 4	SNMP trap A - 3
PNNI Topology State Elements 6 - 2	adding
PNNI Topology State Packets 6 - 2	deleting
policy 6 - 10	displaying
advertise 6 - 10	supported on the switch A - 3
summary 6 - 10	source area ID 6 - 10
suppress 6 - 10	SPANS interface 2 - 3
PTSEs 6 - 2	SPANS SPVC
PTSPs	creating 1 - 32
PVC revalidation 2 - 10	SPANS SPVCs
n.	displaying 1 - 33
R	split switch 6 - 5
rate enforcement	SPVC
configuring egress D - 10	configuring on a FramePlus network
configuring ingress D - 10	module D - 15, D - 21
reachability information 6 - 1	SRTS (Synchronous Residual Time
RFC-1577 2 - 1	Stamp) B - 3
S	Sustainable Cell Rate (SCR)1 - 35, 1 - 36
scope 6 - 10	switchclock 9 - 1
SCR	failover
sdconf.rec file 8 - 5	Synchronous Residual Time Stamp
editing 8 - 6	(SRTS)
transferring to the switch 8 - 5	Т
SecurID	tagging cells 1 - 37
data encryption 8 - 5	00 0
installing the server software 8 - 5	Technical Supportii

terminating paths	virtual path
displaying advanced information 1 - 8	virtual path terminators (VPTs) 1 - 6
timing	VP-scope UNI
ASX-1000 9 - 4 ASX-200BX 9 - 4 ASX-200WG 9 - 4 ForeRunnerLE 155 9 - 5 port level 9 - 3 timing on a switch 9 - 1 traffic policing 1 - 36	
•	
UBR	
V	
Variable Bit Rate (VBR)       1 - 35         VBR       1 - 35         VCI allocation range       7 - 1         VC-space       7 - 7         virtual channel       1 - 11	
virtual channel identifier (VCI) 1 - 1	

Index